# Conjugate Direction minimization Lectures for PHD course on Non-linear equations and numerical optimization

Enrico Bertolazzi

DIMS - Università di Trento

March 2005



## Outline

- Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



# Generic minimization algorithm

In the following we study the convergence rate of the Generic minimization algorithm applied to a quadratic function  $\mathbf{q}(x)$  with exact line search. The function

$$q(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x} - \boldsymbol{b}^T \boldsymbol{x} + c$$

can be viewed as a n-dimensional generalization of the 1-dimensional parabolic model.

#### Generic minimization algorithm

Given an initial guess  $x_0$ , let k = 0;

while not converged do

Find a descent direction  $p_k$  at  $x_k$ ;

Compute a step size  $\alpha_k$  using a line-search along  $p_k$ .

Set  $x_{k+1} = x_k + \alpha_k p_k$  and increase k by 1.

end while



#### Assumption (Symmetry)

The matrix A is assumed to be symmetric, in fact,

$$\boldsymbol{A} = \boldsymbol{A}^{Symm} + \boldsymbol{A}^{Skew}$$

where

$$egin{align} oldsymbol{A}^{Symm} &= rac{1}{2}ig[oldsymbol{A} + oldsymbol{A}^Tig], & oldsymbol{A}^{Symm} &= (oldsymbol{A}^{Symm})^T \ oldsymbol{A}^{Skew} &= rac{1}{2}ig[oldsymbol{A} - oldsymbol{A}^Tig], & oldsymbol{A}^{Skew} &= -(oldsymbol{A}^{Skew})^T \ oldsy$$

moreover

$$x^T A x = x^T A^{Symm} x + x^T A^{Skew} x = x^T A^{Symm} x$$

so that only the symmetric part of A contribute to q(x).





## Assumption (SPD)

The matrix A is assumed to be symmetric and positive definite, in fact,

$$abla \mathsf{q}(oldsymbol{x})^T = rac{1}{2}ig(oldsymbol{A} + oldsymbol{A}^Tig)oldsymbol{x} - oldsymbol{b} = oldsymbol{A}oldsymbol{x} - oldsymbol{b}$$

and

$$abla^2 \mathtt{q}(oldsymbol{x}) = rac{1}{2}ig(oldsymbol{A} + oldsymbol{A}^Tig) = oldsymbol{A}$$

From the sufficient condition for a minimum we have that  $\nabla \mathbf{q}(x_{\star})^T = \mathbf{0}$ , i.e.

$$Ax_{\star} = b$$

and  $abla^2 \mathtt{q}(x_\star) = A$  is SPD.





 In the following we study the convergence rate of the Steepest Descent and Conjugate Gradient methods applied to

$$q(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x} - \boldsymbol{b}^T \boldsymbol{x} + c$$

where A is an SPD matrix.

• This assumption simplify the analysis but it is also useful in the non linear case. In fact, by expanding a generic function f(x) near its minimum  $x_{\star}$  we have

$$egin{aligned} \mathsf{f}(oldsymbol{x}) &= \mathsf{f}(oldsymbol{x}_\star) + 
abla \mathsf{f}(oldsymbol{x}_\star) + \nabla \mathsf{f}(oldsymbol{x}_\star) (oldsymbol{x} - oldsymbol{x}_\star) \\ &+ rac{1}{2} (oldsymbol{x} - oldsymbol{x}_\star)^T 
abla^2 \mathsf{f}(oldsymbol{x}_\star) (oldsymbol{x} - oldsymbol{x}_\star) + \mathcal{O}(\|oldsymbol{x} - oldsymbol{x}_\star\|^3) \end{aligned}$$



By setting

$$egin{aligned} oldsymbol{A} &= 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}), \ oldsymbol{b} &= 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} - 
abla \mathsf{f}(oldsymbol{x}_{\star}) \ c &= \mathsf{f}(oldsymbol{x}_{\star}) - 
abla \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} + rac{1}{2} oldsymbol{x}_{\star}^T 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} \end{aligned}$$

we have

$$\mathsf{f}(oldsymbol{x}) = rac{1}{2} oldsymbol{x}^T oldsymbol{A} oldsymbol{x} - oldsymbol{b}^T oldsymbol{x} + c + \mathcal{O}(\|oldsymbol{x} - oldsymbol{x}_\star\|^3)$$

• So that we expect that when an iterate  $x_k$  is near  $x_\star$  then we can neglect  $\mathcal{O}(\|x-x_\star\|^3)$  and the asymptotic behavior is the same of the quadratic problem.



 we can rewrite the quadratic problem in many different way as follows

$$\mathbf{q}(oldsymbol{x}) = rac{1}{2}(oldsymbol{x} - oldsymbol{x}_{\star})^T oldsymbol{A}(oldsymbol{x} - oldsymbol{x}_{\star}) + c'$$

$$= rac{1}{2}(oldsymbol{A}oldsymbol{x} - oldsymbol{b})^T oldsymbol{A}^{-1}(oldsymbol{A}oldsymbol{x} - oldsymbol{b}) + c'$$

where

$$c' = c + \frac{1}{2} \boldsymbol{x}_{\star}^T \boldsymbol{A} \boldsymbol{x}_{\star}$$

 This last forms are useful in the study of the steepest descent method.



#### Outline

- Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



# The steepest descent for quadratic functions

(1/3)

#### The steepest descent minimization algorithm

Given an initial guess  $x_0$ , let k = 0;

#### while not converged do

Choose as descent direction  $p_k = -\nabla q(x_k)^T = b - Ax_k$ ;

Compute a step size  $\alpha_k$  using a line-search along  $p_k$ .

Set  $x_{k+1} = x_k + \alpha_k p_k$  and increase k by 1.

end while

#### Definition (Residual)

The expressions

$$r(x) = b - Ax$$
,  $r_k = b - Ax_k$ 

are called the residual. We obviously have  $r(x) = -\nabla q(x)^T$  and  $r(x_\star) = \mathbf{0}$ .



# The steepest descent for quadratic functions

We can solve exactly the problem

$$lpha_k = rg \min_{lpha \geq 0} \ \mathsf{q}(oldsymbol{x}_k - lpha oldsymbol{r}_k)$$

because  $p(\alpha) = q(x_k - \alpha r_k)$  is a parabola. In fact

$$\frac{\mathrm{d}p(\alpha)}{\mathrm{d}\alpha} = \frac{\mathrm{d}\mathsf{q}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)}{\mathrm{d}\alpha} = -\nabla \mathsf{q}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)\boldsymbol{r}_k = 0$$

but

$$egin{aligned} \mathbf{0} &= -
abla \mathsf{q}(oldsymbol{x}_k - lpha oldsymbol{r}_k) oldsymbol{r}_k = oldsymbol{r}(oldsymbol{x}_k - lpha oldsymbol{r}_k)^T oldsymbol{r}_k \ &= oldsymbol{(r_k - lpha oldsymbol{A} oldsymbol{r}_k)}^T oldsymbol{r}_k \end{aligned}$$

and the minimum is at lpha set to  $\dfrac{m{r}_k^Tm{r}_k}{m{r}_k^Tm{A}m{r}_k}.$ 



#### The steepest descent minimization algorithm

Given an initial guess  $x_0$ , let k = 0;

while not converged do

Compute  $r_k = b - Ax_k$ ;

Compute the step size  $\alpha_k = \frac{\boldsymbol{r}_k^T \boldsymbol{r}_k}{\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k}$ ;

Set  $x_{k+1} = x_k + \alpha_k r_k$  and increase k by 1.

end while

Or more compactly

$$oldsymbol{x}_{k+1} = oldsymbol{x}_k + rac{oldsymbol{r}_k^Toldsymbol{r}_k}{oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k}oldsymbol{r}_k$$



# The steepest descent reduction step

We want bound  $q(x_{k+1})$  by  $q(x_k)$ :

$$\begin{aligned} \mathsf{q}(\boldsymbol{x}_{k+1}) &= \mathsf{q}\left(\boldsymbol{x}_k + \alpha_k \boldsymbol{r}_k\right) \\ &= \frac{1}{2}\left(\boldsymbol{A}\boldsymbol{x}_k + \alpha_k \boldsymbol{A}\boldsymbol{r}_k - \boldsymbol{b}\right)^T \boldsymbol{A}^{-1}\left(\boldsymbol{A}\boldsymbol{x}_k + \alpha_k \boldsymbol{A}\boldsymbol{r}_k - \boldsymbol{b}\right) + c' \\ &= \frac{1}{2}\left(\alpha_k \boldsymbol{A}\boldsymbol{r}_k - \boldsymbol{r}_k\right)^T \boldsymbol{A}^{-1}\left(\alpha_k \boldsymbol{A}\boldsymbol{r}_k - \boldsymbol{r}_k\right) + c' \\ &= \frac{1}{2}\boldsymbol{r}_k^T \boldsymbol{A}^{-1}\boldsymbol{r}_k + \frac{1}{2}\alpha_k^2 \boldsymbol{r}_k^T \boldsymbol{A}\boldsymbol{r}_k - \alpha_k \boldsymbol{r}_k^T \boldsymbol{r}_k + c' \\ &= \mathsf{q}(\boldsymbol{x}_k) + \frac{1}{2}\alpha_k\left(\alpha_k \boldsymbol{r}_k^T \boldsymbol{A}\boldsymbol{r}_k - 2\boldsymbol{r}_k^T \boldsymbol{r}_k\right) \end{aligned}$$



# The steepest descent reduction step

Substituting  $lpha_k = rac{m{r}_k^Tm{r}_k}{m{r}_k^Tm{A}m{r}_k}$  we obtain

$$\mathsf{q}(oldsymbol{x}_{k+1}) = \mathsf{q}(oldsymbol{x}_k) - rac{1}{2} rac{(oldsymbol{r}_k^T oldsymbol{r}_k)^2}{oldsymbol{r}_k^T oldsymbol{A} oldsymbol{r}_k}$$

this shows that the steepest descent method reduce at each step the objective function q(x).

Using the expression  $q(x) = \frac{1}{2}r(x)^TA^{-1}r(x) + c'$  we can write:

$$\frac{1}{2} \boldsymbol{r}_{k+1}^T \boldsymbol{A}^{-1} \boldsymbol{r}_{k+1} = \frac{1}{2} \boldsymbol{r}_k^T \boldsymbol{A}^{-1} \boldsymbol{r}_k - \frac{1}{2} \frac{(\boldsymbol{r}_k^T \boldsymbol{r}_k)^2}{\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k}$$



# The steepest descent reduction step

or better

$$oldsymbol{r}_{k+1}^Toldsymbol{A}^{-1}oldsymbol{r}_{k+1} = oldsymbol{r}_k^Toldsymbol{A}^{-1}oldsymbol{r}_k \left(1-rac{(oldsymbol{r}_k^Toldsymbol{A}^T)^2}{(oldsymbol{r}_k^Toldsymbol{A}^{-1}oldsymbol{r}_k)(oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k)}
ight)$$

noticing that  $m{r}_k = m{b} - m{A}m{x}_k = m{A}m{x}_\star - m{A}m{x}_k = m{A}(m{x}_\star - m{x}_k)$  we have

$$\left\|oldsymbol{x}_{\star}-oldsymbol{x}_{k+1}
ight\|_{oldsymbol{A}}^{2}=\left\|oldsymbol{x}_{\star}-oldsymbol{x}_{k}
ight\|_{oldsymbol{A}}^{2}\left(1-rac{(oldsymbol{r}_{k}^{T}oldsymbol{r}_{k})^{2}}{(oldsymbol{r}_{k}^{T}oldsymbol{A}^{-1}oldsymbol{r}_{k})(oldsymbol{r}_{k}^{T}oldsymbol{A}oldsymbol{r}_{k})}
ight)$$

where

$$\|\boldsymbol{x}\|_{\boldsymbol{A}} = \sqrt{\boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x}}$$

is the energy norm induced by the SPD matrix A.



The estimate of the convergence rate for the steepest descent method is linked to the estimate of the term

$$rac{(oldsymbol{r}_k^Toldsymbol{r}_k)^2}{(oldsymbol{r}_k^Toldsymbol{A}^{-1}oldsymbol{r}_k)(oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k)}$$

in particular we can prove

#### Lemma (Kantorovic)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the following inequality is valid

$$1 \leq rac{(m{x}^Tm{A}m{x})(m{x}^Tm{A}^{-1}m{x})}{(m{x}^Tm{x})^2} \leq rac{(M+m)^2}{4\,M\,m}$$

for all  $x \neq 0$ . Where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.





Proof. (1/5).

STEP 1: problem reformulation. First of all notice that

$$rac{(m{x}^Tm{A}m{x})(m{x}^Tm{A}^{-1}m{x})}{(m{x}^Tm{x})^2} = rac{(m{y}^Tm{A}m{y})(m{y}^Tm{A}^{-1}m{y})}{(m{y}^Tm{y})^2}$$

for all  $y = \alpha x$  with  $\alpha \neq 0$ . Choosing  $\alpha = ||x||^{-1}$  have:





Proof. (2/5).

STEP 2: eigenvector expansions. Matrix  $A \in \mathbb{R}^{n \times n}$  is an SPD matrix so that there exists  $u_1, u_2, \ldots, u_n$  a complete orthonormal eigenvectors set with  $0 < \lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n$  corresponding eigenvalues. Let be  $x \in \mathbb{R}^n$  then

$$oldsymbol{x} = \sum_{k=1}^n lpha_k oldsymbol{u}_k, \qquad oldsymbol{x}^T oldsymbol{x} = \sum_{k=1}^n lpha_k^2$$

so that  $(\boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x}) (\boldsymbol{x}^T \boldsymbol{A}^{-1} \boldsymbol{x}) = h(\alpha_1, \dots, \alpha_n)$  where

$$h(\alpha_1,\ldots,\alpha_n) = \left(\sum_{k=1}^n \alpha_k^2 \lambda_k\right) \left(\sum_{k=1}^n \alpha_k^2 \lambda_k^{-1}\right)$$

then the lemma can be reformulated:

- Find maxima and minima of  $h(\alpha_1, \ldots, \alpha_n)$
- subject to  $\sum_{k=1}^{n} \alpha_k^2 = 1$ .



# Proof. (3/5).

STEP 3: problem reduction. By using Lagrange multiplier maxima and minima are the stationary points of:

$$g(\alpha_1,\ldots,\alpha_n,\mu)=h(\alpha_1,\ldots,\alpha_n)+\mu\left(\sum_{k=1}^n\alpha_k^2-1\right)$$

setting  $A=\sum_{k=1}^n \alpha_k^2 \lambda_k$  and  $B=\sum_{k=1}^n \alpha_k^2 \lambda_k^{-1}$  we have

$$\frac{\partial g(\alpha_1,\ldots,\alpha_n,\mu)}{\partial \alpha_k} = 2\alpha_k (\lambda_k B + \lambda_k^{-1} A + \mu) = 0$$

so that

- ② Or  $\lambda_k$  is a root of the quadratic polynomial  $\lambda^2 B + \lambda \mu + A$ . in any case there are at most 2 coefficients  $\alpha$ 's not zero. <sup>a</sup>



athe argument should be improved in the case of multiple eigenvalues

Proof. (4/5).

STEP 4: problem reformulation. say  $\alpha_i$  and  $\alpha_j$  are the only non zero coefficients, then  $\alpha_i^2 + \alpha_i^2 = 1$  and we can write

$$h(\alpha_1, \dots, \alpha_n) = (\alpha_i^2 \lambda_i + \alpha_j^2 \lambda_j) (\alpha_i^2 \lambda_i^{-1} + \alpha_j^2 \lambda_j^{-1})$$

$$= \alpha_i^4 + \alpha_j^4 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i}\right)$$

$$= \alpha_i^2 (1 - \alpha_j^2) + \alpha_j^2 (1 - \alpha_i^2) + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i}\right)$$

$$= 1 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i} - 2\right)$$

$$= 1 + \alpha_i^2 (1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_i}$$





#### Proof. (5/5).

STEP 5: bounding maxima and minima. notice that

$$0 \le \beta(1-\beta) \le \frac{1}{4}, \quad \forall \beta \in [0,1]$$

$$1 \le 1 + \alpha_i^2 (1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_j} \le 1 + \frac{(\lambda_i - \lambda_j)^2}{4\lambda_i \lambda_j} = \frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j}$$

to bound  $(\lambda_i + \lambda_j)^2/(4\lambda_i\lambda_j)$  consider the function  $f(x) = (1+x)^2/x$  which is increasing for  $x \ge 1$  so that we have

$$\frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j} \le \frac{(M+m)^2}{4Mm}$$

and finally

$$1 \leq h(\alpha_1, \ldots, \alpha_n) \leq \frac{(M+m)^2}{4 M m}$$



# Convergence rate of Steepest Descent

The Kantorovich inequality permits to prove:

#### Theorem (Convergence rate of Steepest Descent)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the steepest descent method:

$$oldsymbol{x}_{k+1} = oldsymbol{x}_k + rac{oldsymbol{r}_k^Toldsymbol{r}_k}{oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k}oldsymbol{r}_k$$

converge to the solution  $x_\star = A^{-1}b$  with at least linear q-rate in the norm  $\|\cdot\|_A$ . Moreover we have the error estimate

$$\left\|oldsymbol{x}_{k+1} - oldsymbol{x}_{\star}
ight\|_{oldsymbol{A}} \leq rac{\kappa-1}{\kappa+1} \left\|oldsymbol{x}_{k} - oldsymbol{x}_{\star}
ight\|_{oldsymbol{A}}$$

 $\kappa = M/m$  is the condition number where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.



#### Proof.

Remember from slide  $N^{\circ}15$ 

$$\left\|oldsymbol{x}_{\star}-oldsymbol{x}_{k+1}
ight\|_{oldsymbol{A}}^{2}=\left\|oldsymbol{x}_{\star}-oldsymbol{x}_{k}
ight\|_{oldsymbol{A}}^{2}\left(1-rac{(oldsymbol{r}_{k}^{T}oldsymbol{r}_{k})^{2}}{(oldsymbol{r}_{k}^{T}oldsymbol{A}^{-1}oldsymbol{r}_{k})(oldsymbol{r}_{k}^{T}oldsymbol{A}oldsymbol{r}_{k})}
ight)$$

from Kantorovich inequality

$$1 - \frac{(\boldsymbol{r}_k^T \boldsymbol{r}_k)^2}{(\boldsymbol{r}_k^T \boldsymbol{A}^{-1} \boldsymbol{r}_k)(\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k)} \le 1 - \frac{4 M m}{(M+m)^2} = \frac{(M-m)^2}{(M+m)^2}$$

so that

$$\left\| \boldsymbol{x}_{\star} - \boldsymbol{x}_{k+1} \right\|_{\boldsymbol{A}} \leq \frac{M-m}{M+m} \left\| \boldsymbol{x}_{\star} - \boldsymbol{x}_{k} \right\|_{\boldsymbol{A}}$$





#### Remark (One step convergence)

The steepest descent method can converge in one iteration if  $\kappa = 1$  or when  $r_0 = u_k$  where  $u_k$  is an eigenvector of A.

- In the first case  $(\kappa = 1)$  we have  $A = \beta I$  for some  $\beta > 0$  so it is not interesting.
- 2 In the second case we have

$$\frac{(\boldsymbol{u}_k^T\boldsymbol{u}_k)^2}{(\boldsymbol{u}_k^T\boldsymbol{A}^{-1}\boldsymbol{u}_k)(\boldsymbol{u}_k^T\boldsymbol{A}\boldsymbol{u}_k)} = \frac{(\boldsymbol{u}_k^T\boldsymbol{u}_k)^2}{\lambda_k^{-1}(\boldsymbol{u}_k^T\boldsymbol{u}_k)\lambda_k(\boldsymbol{u}_k^T\boldsymbol{u}_k)} = 1$$

in both cases we have  $r_1 = \mathbf{0}$  i.e. we have found the solution.



## Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



# Conjugate direction method

## Definition (Conjugate vector)

Given two vectors p and q in  $\mathbb{R}^n$  are conjugate respect to A if they are orthogonal respect the scalar product induced by A; i.e.,

$$\boldsymbol{p}^T \boldsymbol{A} \boldsymbol{q} = \sum_{i,j=1}^n A_{ij} p_i q_j = 0.$$

Clearly, n vectors  $p_1, p_2, \dots p_n \in \mathbb{R}^n$  that are pair wise conjugated respect to A form a base of  $\mathbb{R}^n$ .



## Problem (Linear system)

Find the minimum of  $q(x) = \frac{1}{2}x^TAx - b^Tx + c$  is equivalent to solve the first order necessary condition, i.e.

Find  $\mathbf{x}_{\star} \in \mathbb{R}^n$  such that:  $A\mathbf{x}_{\star} = \mathbf{b}$ .

#### Observation

Consider  $x_0 \in \mathbb{R}^n$  and decompose the error  $e_0 = x_\star - x_0$  by the conjugate vectors  $p_1$ ,  $p_2, \ldots, p_n \in \mathbb{R}^n$ :

$$e_0 = x_{\star} - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \cdots + \sigma_n p_n.$$

Evaluating the coefficients  $\sigma_1$ ,  $\sigma_2$ , ...,  $\sigma_n \in \mathbb{R}$  is equivalent to solve the problem  $Ax_{\star} = b$ , because knowing  $e_0$  we have

$$x_{\star} = x_0 + e_0.$$





#### Observation

Using conjugacy the coefficients  $\sigma_1$ ,  $\sigma_2$ , ...,  $\sigma_n \in \mathbb{R}$  can be computed as

$$\sigma_i = rac{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{e}_0}{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i}, \qquad for \ i=1,2,\ldots,n.$$

In fact, for all  $1 \le i \le n$ , we have

$$\begin{aligned} \boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{e}_0 &= \boldsymbol{p}_i^T \boldsymbol{A} \left( \sigma_1 \boldsymbol{p}_1 + \sigma_2 \boldsymbol{p}_2 + \ldots + \sigma_n \boldsymbol{p}_n \right), \\ &= \sigma_1 \boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_1 + \sigma_2 \boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_2 + \ldots + \sigma_n \boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_n, \\ &= \sigma_i \boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_i, \end{aligned}$$

because  $\mathbf{p}_i^T \mathbf{A} \mathbf{p}_j = 0$  for  $i \neq j$ .





The conjugate direction method evaluate the coefficients  $\sigma_1$ ,  $\sigma_2, \ldots, \sigma_n \in \mathbb{R}$  recursively in n steps, solving for  $k \geq 0$  the minimization problem:

#### Conjugate direction method

 $\begin{aligned} &\text{Given } \boldsymbol{x}_0; \ k \leftarrow 0; \\ &\textbf{repeat} \\ & k \leftarrow k+1; \\ &\text{Find } \boldsymbol{x}_k \in \boldsymbol{x}_0 + \mathcal{V}_k \text{ such that:} \end{aligned}$ 

$$oldsymbol{x}_k = \mathop{\mathsf{arg\,min}}\limits_{oldsymbol{x} \in oldsymbol{x}_0 + \mathcal{V}_k} \|oldsymbol{x}_\star - oldsymbol{x}\|_{oldsymbol{A}}$$

until k=n

where  $\mathcal{V}_k$  is the subspace of  $\mathbb{R}^n$  generated by the first k conjugate direction; i.e.,

$$\mathcal{V}_k = \operatorname{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \dots, \boldsymbol{p}_k\}.$$





# Step: $x_0 \rightarrow x_1$

At the first step we consider the subspace  $x_0 + \mathrm{SPAN}\{p_1\}$  which consists in vectors of the form

$$x(\alpha) = x_0 + \alpha p_1 \qquad \alpha \in \mathbb{R}$$

The minimization problem becomes:

#### Minimization step $x_0 o x_1$

Find  $x_1 = x_0 + \alpha_1 p_1$  (i.e., find  $\alpha_1$ !) such that:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{1}\|_{\boldsymbol{A}} = \min_{\alpha \in \mathbb{R}} \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})\|_{\boldsymbol{A}},$$



# Solving first step method 1

The minimization problem is the minimum respect to  $\alpha$  of the quadratic:

$$\begin{aligned} \Phi(\alpha) &= \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})\|_{\boldsymbol{A}}^{2}, \\ &= (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1}))^{T} \boldsymbol{A} (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})), \\ &= (\boldsymbol{e}_{0} - \alpha \boldsymbol{p}_{1})^{T} \boldsymbol{A} (\boldsymbol{e}_{0} - \alpha \boldsymbol{p}_{1}), \\ &= \boldsymbol{e}_{0}^{T} \boldsymbol{A} \boldsymbol{e}_{0} - 2\alpha \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{e}_{0} + \alpha^{2} \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{p}_{1}. \end{aligned}$$

minimum is found by imposing:

$$\frac{\mathsf{d}\Phi(\alpha)}{\mathsf{d}\alpha} = -2\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{e}_0 + 2\alpha\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{p}_1 = 0 \quad \Rightarrow \quad \alpha_1 = \frac{\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{e}_0}{\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{p}_1}$$





Remember the error expansion:

$$\boldsymbol{x}_{\star} - \boldsymbol{x}_0 = \sigma_1 \boldsymbol{p}_1 + \sigma_2 \boldsymbol{p}_2 + \cdots + \sigma_n \boldsymbol{p}_n.$$

Let  $x(\alpha) = x_0 + \alpha p_1$ , the difference  $x_{\star} - x(\alpha)$  becomes:

$$\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha) = (\sigma_1 - \alpha)\boldsymbol{p}_1 + \sigma_2\boldsymbol{p}_2 + \ldots + \sigma_n\boldsymbol{p}_n$$

due to conjugacy the error  $\| \boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha) \|_{\boldsymbol{A}}$  becomes

$$\begin{aligned} \|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2} \\ &= \left( (\sigma_{1} - \alpha)\boldsymbol{p}_{1} + \sum_{i=2}^{n} \sigma_{i}\boldsymbol{p}_{i} \right)^{T} \boldsymbol{A} \left( (\sigma_{1} - \alpha)\boldsymbol{p}_{1} + \sum_{j=2}^{n} \sigma_{j}\boldsymbol{p}_{i} \right) \\ &= (\sigma_{1} - \alpha)^{2} \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{p}_{1} + \sum_{j=2}^{n} \sigma_{j}^{2} \boldsymbol{p}_{j}^{T} \boldsymbol{A} \boldsymbol{p}_{j} \end{aligned}$$





# Solving first step method 2

(2/2)

**Because** 

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2} = (\sigma_{1} - \alpha)^{2} \|\boldsymbol{p}_{1}\|_{\boldsymbol{A}}^{2} + \sum_{i=2}^{n} \sigma_{2}^{2} \|\boldsymbol{p}_{i}\|_{\boldsymbol{A}}^{2},$$

we have that

$$\|oldsymbol{x}_{\star} - oldsymbol{x}(lpha_1)\|_{oldsymbol{A}}^2 = \sum_{i=2}^n \sigma_i^2 \, \|oldsymbol{p}_i\|_{oldsymbol{A}}^2 \leq \|oldsymbol{x}_{\star} - oldsymbol{x}(lpha)\|_{oldsymbol{A}}^2 \qquad ext{for all } lpha 
eq \sigma_1$$

so that minimum is found by imposing  $\alpha_1 = \sigma_1$ :

$$\alpha_1 = \frac{\boldsymbol{p}_1^T \boldsymbol{A} \boldsymbol{e}_0}{\boldsymbol{p}_1^T \boldsymbol{A} \boldsymbol{p}_1}$$

This argument can be generalized for all k > 1 (see next slides).



## Step, $oldsymbol{x}_{k-1} ightarrow oldsymbol{x}_k$

For the step from k-1 to k we consider the subspace of  $\mathbb{R}^n$ 

$$\mathcal{V}_k = \operatorname{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \dots, \boldsymbol{p}_k\}$$

which contains vectors of the form:

$$x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_0 + \alpha^{(1)}p_1 + \alpha^{(2)}p_2 + \dots + \alpha^{(k)}p_k$$

The minimization problem becomes:

#### Minimization step $oldsymbol{x}_{k-1} ightarrow oldsymbol{x}_k$

Find  $x_k = x_0 + \alpha_1 p_1 + \alpha_2 p_2 + \ldots + \alpha_k p_k$  (i.e.  $\alpha_1, \alpha_2, \ldots, \alpha_k$ ) such that:

$$\left\|oldsymbol{x}_{\star} - oldsymbol{x}_{k}
ight\|_{oldsymbol{A}} = \min_{lpha^{(1)},lpha^{(2)},...,lpha^{(k)} \in \mathbb{R}} \left\|oldsymbol{x}_{\star} - oldsymbol{x}(lpha^{(1)},lpha^{(2)},\ldots,lpha^{(k)})
ight\|_{oldsymbol{A}}$$





# Solving kth Step: $\boldsymbol{x}_{k-1} \rightarrow \boldsymbol{x}_k$

(1/2)

Remember the error expansion:

$$x_{\star} - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \cdots + \sigma_n p_n.$$

Consider a vector of the form

$$x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_0 + \alpha^{(1)}p_1 + \alpha^{(2)}p_2 + \dots + \alpha^{(k)}p_k$$

the error  $x_{\star} - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})$  can be written as

$$x_{\star} - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_{\star} - x_{0} - \sum_{i=1}^{k} \alpha^{(i)} p_{i},$$

$$= \sum_{i=1}^{k} (\sigma_i - \alpha^{(i)}) \boldsymbol{p}_i + \sum_{i=k+1}^{n} \sigma_i \boldsymbol{p}_i.$$





# Solving kth Step: $\boldsymbol{x}_{k-1} \rightarrow \boldsymbol{x}_k$

using conjugacy of  $p_i$  we obtain the norm of the error:

$$\left\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})\right\|_{\boldsymbol{A}}^{2}$$

$$= \sum_{i=1}^{k} \left(\sigma_{i} - \alpha^{(i)}\right)^{2} \left\|\boldsymbol{p}_{i}\right\|_{\boldsymbol{A}}^{2} + \sum_{i=k+1}^{n} \sigma_{i}^{2} \left\|\boldsymbol{p}_{i}\right\|_{\boldsymbol{A}}^{2}.$$

So that minimum is found by imposing  $\alpha_i = \sigma_i$ : for  $i = 1, 2, \dots, k$ .

$$egin{aligned} lpha_i = rac{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{e}_0}{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i} \end{aligned} \qquad i = 1, 2, \ldots, k$$





• notice that  $\alpha_i = \sigma_i$  and that

$$x_k = x_0 + \alpha_1 p_1 + \dots + \alpha_k p_k$$
  
=  $x_{k-1} + \alpha_k p_k$ 

- so that  $x_{k-1}$  contains k-1 coefficients  $\alpha_i$  for the minimization.
- if we consider the one dimensional minimization on the subspace  $x_{k-1} + \operatorname{SPAN}\{p_k\}$  we find again  $x_k$ !





(2/3)

### Successive one dimensional minimization

Consider a vector of the form

$$\boldsymbol{x}(\alpha) = \boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_k$$

remember that  $x_{k-1} = x_0 + \alpha_1 p_1 + \cdots + \alpha_{k-1} p_{k-1}$  so that the error  $x_{\star} - x(\alpha)$  can be written as

$$egin{aligned} oldsymbol{x}_{\star} - oldsymbol{x}(lpha) &= oldsymbol{x}_{\star} - oldsymbol{x}_0 - \sum_{i=1}^{k-1} lpha_i oldsymbol{p}_i + lpha oldsymbol{p}_k + lpha oldsymbol{p}_k + \sum_{i=k+1}^n \sigma_i oldsymbol{p}_i. \end{aligned}$$

due to the equality  $\sigma_i = \alpha_i$  the blue part of the expression is 0.





Using conjugacy of  $p_i$  we obtain the norm of the error:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2} = \left(\sigma_{k} - \alpha\right)^{2} \|\boldsymbol{p}_{k}\|_{\boldsymbol{A}}^{2} + \sum_{i=k+1}^{n} \sigma_{i}^{2} \|\boldsymbol{p}_{i}\|_{\boldsymbol{A}}^{2}.$$

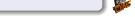
So that minimum is found by imposing  $\alpha = \sigma_k$ :

$$lpha_k = rac{oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_0}{oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_k}$$

#### Remark

This observation permit to perform the minimization on the k-dimensional space  $x_0 + \mathcal{V}_k$  as successive one dimensional minimizations along the conjugate directions  $p_k!$ .





#### Problem (one dimensional successive minimization)

Find  $x_k = x_{k-1} + \alpha_k p_k$  such that:

$$\left\|oldsymbol{x}_{\star} - oldsymbol{x}_{k}
ight\|_{oldsymbol{A}} = \min_{lpha \in \mathbb{R}} \left\|oldsymbol{x}_{\star} - (oldsymbol{x}_{k-1} + lpha oldsymbol{p}_{k})
ight\|_{oldsymbol{A}},$$

The solution is the minimum respect to  $\alpha$  of the quadratic:

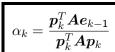
$$\Phi(\alpha) = (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_{k}))^{T} \boldsymbol{A} (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_{k})),$$

$$= (\boldsymbol{e}_{k-1} - \alpha \boldsymbol{p}_{k})^{T} \boldsymbol{A} (\boldsymbol{e}_{k-1} - \alpha \boldsymbol{p}_{k}),$$

$$= \boldsymbol{e}_{k-1}^{T} \boldsymbol{A} \boldsymbol{e}_{k-1} - 2\alpha \boldsymbol{p}_{k}^{T} \boldsymbol{A} \boldsymbol{e}_{k-1} + \alpha^{2} \boldsymbol{p}_{k}^{T} \boldsymbol{A} \boldsymbol{p}_{k}.$$

minimum is found by imposing:

$$\frac{\mathsf{d}\Phi(\alpha)}{\mathsf{d}\alpha} = -2\boldsymbol{p}_k^T\boldsymbol{A}\boldsymbol{e}_{k-1} + 2\alpha\boldsymbol{p}_k^T\boldsymbol{A}\boldsymbol{p}_k = 0 \quad \Rightarrow \quad \left[\alpha_k = \frac{\boldsymbol{p}_k^T\boldsymbol{A}\boldsymbol{e}_{k-1}}{\boldsymbol{p}_k^T\boldsymbol{A}\boldsymbol{p}_k}\right]$$







• In the case of minimization on the subspace  $x_0 + \mathcal{V}_k$  we have:

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_0 / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k$$

• In the case of one dimensional minimization on the subspace  $x_{k-1} + \operatorname{SPAN}\{p_k\}$  we have:

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k$$

ullet Apparently they are different results, however by using the conjugacy of the vectors  $oldsymbol{p}_i$  we have

$$egin{aligned} oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_{k-1} &= oldsymbol{p}_k^T oldsymbol{A} (oldsymbol{x}_{\star} - oldsymbol{x}_{k-1}) \ &= oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_{0} - lpha_1 oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_{1} + \dots + lpha_{k-1} oldsymbol{p}_{k-1}) ig) \ &= oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_{0} - lpha_1 oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_{1} - \dots - lpha_{k-1} oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_{k-1} \\ &= oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_{0} \end{aligned}$$



- The one step minimization in the space  $x_0 + \mathcal{V}_n$  and the successive minimization in the space  $x_{k-1} + \operatorname{SPAN}\{p_k\}$ ,  $k = 1, 2, \ldots, n$  are equivalent if  $p_i$ s are conjugate.
- The successive minimization is useful when  $p_i$ s are not known in advance but must be computed as the minimization process proceeds.
- ullet The evaluation of  $lpha_k$  is apparently not computable because  $e_i$  is not known. However noticing

$$oldsymbol{A}oldsymbol{e}_k = oldsymbol{A}(oldsymbol{x}_\star - oldsymbol{x}_k) = oldsymbol{b} - oldsymbol{A}oldsymbol{x}_k = oldsymbol{r}_k$$

we can write

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k = \boldsymbol{p}_k^T \boldsymbol{r}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k =$$

• Finally for the residual is valid the recurrence

$$r_k = b - Ax_k = b - A(x_{k-1} + \alpha_k p_k) = r_{k-1} - \alpha_k Ap_k.$$



### Conjugate direction minimization

#### Algorithm (Conjugate direction minimization)

$$\begin{aligned} k &\leftarrow 0; \ x_0 \ assigned; \\ r_0 &\leftarrow b - Ax_0; \\ \text{while } \textit{not converged do} \\ k &\leftarrow k+1; \\ \alpha_k &\leftarrow \frac{r_{k-1}^T p_k^T}{p_k A p_k}; \\ x_k &\leftarrow x_{k-1} + \alpha_k p_k; \\ r_k &\leftarrow r_{k-1} - \alpha_k A p_k; \\ \textit{end while} \end{aligned}$$

### Observation (Computazional cost)

The conjugate direction minimization requires at each step one matrix–vector product for the evaluation of  $\alpha_k$  and two update AXPY for  $x_k$  and  $r_k$ .





#### Monotonic behavior of the error

#### Remark (Monotonic behavior of the error)

The energy norm of the error  $\|e_k\|_A$  is monotonically decreasing in k. In fact:

$$e_k = x_{\star} - x_k = \alpha_{k+1} p_{k+1} + \ldots + \alpha_n p_n$$

and by conjugacy

$$\|\boldsymbol{e}_{k}\|_{\boldsymbol{A}}^{2} = \|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}}^{2} = \sigma_{k+1}^{2} \|\boldsymbol{p}_{k+1}\|_{\boldsymbol{A}}^{2} + \ldots + \sigma_{n}^{2} \|\boldsymbol{p}_{n}\|_{\boldsymbol{A}}^{2}.$$

Finally from this relation we have  $e_n = \mathbf{0}$ .



#### Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



## Conjugate Gradient method

The Conjugate Gradient method combine the Conjugate Direction method with an orthogonalization process (like Gram-Schmidt) applied to the residual to construct the conjugate directions. In fact, because  $\boldsymbol{A}$  define a scalar product in the next slide we prove:

- each residue is orthogonal to the previous conjugate directions, and consequently linearly independent from the previous conjugate directions.
- if the residual is not null is can be used to construct a new conjugate direction.





## Orthogonality of the residue $r_k$ respect $\mathcal{V}_k$

• The residue  $r_k$  is orthogonal to  $p_1, p_2, \ldots, p_k$ . In fact, from the error expansion

$$e_k = \alpha_{k+1} p_{k+1} + \alpha_{k+2} p_{k+2} + \cdots + \alpha_n p_n$$

because  $r_k = Ae_k$ , for  $i = 1, 2, \dots, k$  we have

$$egin{aligned} m{p}_i^Tm{r}_k &= m{p}_i^Tm{A}m{e}_k \ &= m{p}_i^Tm{A}\sum_{j=k+1}^n lpha_jm{p}_j = \sum_{j=k+1}^n lpha_jm{p}_i^Tm{A}m{p}_j \ &= 0 \end{aligned}$$





## Building new conjugate direction

- The conjugate direction method build one new direction at each step.
- ullet If  $m{r}_k 
  eq m{0}$  it can be used to build the new direction  $m{p}_{k+1}$  by a Gram-Schmidt orthogonalization process

$$p_{k+1} = r_k + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \ldots + \beta_k^{(k+1)} p_k,$$

where the k coefficients  $\beta_1^{(k+1)}$ ,  $\beta_2^{(k+1)}$ , ...,  $\beta_k^{(k+1)}$  must satisfy:

$$\boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_{k+1} = 0, \quad \text{for } i = 1, 2, \dots, k.$$





## Building new conjugate direction

(repeating from previous slide)

$$p_{k+1} = r_k + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \dots + \beta_k^{(k+1)} p_k,$$

expanding the expression:

$$0 = \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{k+1},$$

$$= \boldsymbol{p}_{i}^{T} \boldsymbol{A} (\boldsymbol{r}_{k} + \beta_{1}^{(k+1)} \boldsymbol{p}_{1} + \beta_{2}^{(k+1)} \boldsymbol{p}_{2} + \dots + \beta_{k}^{(k+1)} \boldsymbol{p}_{k}),$$

$$= \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{r}_{k} + \beta_{i}^{(k+1)} \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{i},$$

$$\Rightarrow \beta_{i}^{(k+1)} = -\frac{\boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{r}_{k}}{\boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{i}} \qquad i = 1, 2, \dots, k$$





The choice of the residual  $r_k \neq \mathbf{0}$  for the construction of the new conjugate direction  $p_{k+1}$  has three important consequences:

- simplification of the expression for  $\alpha_k$ ;
- ② Orthogonality of the residual  $r_k$  from the previous residue  $r_0$ ,  $r_1, \ldots, r_{k-1}$ ;
- § three point formula and simplification of the coefficients  $\beta_i^{(k+1)}$ .

this facts will be examined in the next slides.



## Simplification of the expression for $\alpha_k$

Writing the expression for  $p_k$  from the orthogonalization process

$$p_k = r_{k-1} + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \ldots + \beta_{k-1}^{(k+1)} p_{k-1},$$

using orthogonality of  $r_{k-1}$  and the vectors  $p_1$ ,  $p_2$ ,  $\ldots$ ,  $p_{k-1}$ , (see slide N.47) we have

$$egin{aligned} m{r}_{k-1}^T m{p}_k &= m{r}_{k-1}^T ig( m{r}_{k-1} + eta_1^{(k+1)} m{p}_1 + eta_3^{(k+1)} m{p}_2 + \ldots + eta_{k-1}^{(k+1)} m{p}_{k-1} ig), \ &= m{r}_{k-1}^T m{r}_{k-1}. \end{aligned}$$

recalling the definition of  $\alpha_k$  it follows:

$$\alpha_k = \frac{\boldsymbol{e}_{k-1}^T \boldsymbol{A} \boldsymbol{p}_k}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k} = \frac{\boldsymbol{r}_{k-1}^T \boldsymbol{p}_k}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k} = \boxed{\frac{\boldsymbol{r}_{k-1}^T \boldsymbol{r}_{k-1}}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k}}$$





### Orthogonally of the residue $r_k$ from $r_0, r_1, \ldots, r_{k-1}$

From the definition of  $p_{i+1}$  it follows:

$$\begin{aligned} \boldsymbol{p}_{i+1} &= \boldsymbol{r}_i + \beta_1^{(i+1)} \boldsymbol{p}_1 + \beta_2^{(i+1)} \boldsymbol{p}_2 + \ldots + \beta_i^{(i+1)} \boldsymbol{p}_i, \\ &\Rightarrow \quad \boldsymbol{r}_i \in \text{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \ldots, \boldsymbol{p}_i, \boldsymbol{p}_{i+1}\} = \mathcal{V}_{i+1} \end{aligned} \quad \text{(obvious)}$$

using orthogonality of  $r_k$  and the vectors  $p_1$ ,  $p_2$ , ...,  $p_k$ , (see slide N.47) for i < k we have

$$egin{aligned} m{r}_k^Tm{r}_i &= m{r}_k^Tigg(m{p}_{i+1} - \sum_{j=1}^ieta_j^{(i+1)}m{p}_jigg), \ &= m{r}_k^Tm{p}_{i+1} - \sum_{j=1}^ieta_j^{(i+1)}m{r}_k^Tm{p}_j = 0. \end{aligned}$$





# Three point formula and simplification of $eta_i^{(k+1)}$

From the relation  $m{r}_k^Tm{r}_i = m{r}_k^T(m{r}_{i-1} - lpha_im{A}m{p}_i)$  we deduce

$$m{r}_k^T m{A} m{p}_i = rac{m{r}_k^T m{r}_{i-1} - m{r}_k^T m{r}_i}{lpha_i} = egin{cases} -m{r}_k^T m{r}_k / lpha_k & ext{if } i = k; \\ 0 & ext{if } i < k; \end{cases}$$

remembering that  $lpha_k = m{r}_{k-1}^T m{r}_{k-1} \ / \ m{p}_k^T m{A} m{p}_k$  we obtain

$$eta_i^{(k+1)} = -rac{oldsymbol{r}_k^Toldsymbol{A}oldsymbol{p}_i}{oldsymbol{p}_i^Toldsymbol{A}oldsymbol{p}_i} = \left\{ egin{array}{c} rac{oldsymbol{r}_k^Toldsymbol{r}_k}{oldsymbol{r}_{k-1}^Toldsymbol{r}_{k-1}} & i = k; \ 0 & i < k; \end{array} 
ight.$$

i.e. there is only one non zero coefficient  $\beta_k^{(k+1)}$ , so we write  $\beta_k = \beta_k^{(k+1)}$  and obtain the three point formula:

$$\boldsymbol{p}_{k+1} = \boldsymbol{r}_k + \beta_k \boldsymbol{p}_k$$





## Conjugate gradient algorithm

#### initial step:

$$k \leftarrow 0$$
;  $x_0$  assigned;  $r_0 \leftarrow b - Ax_0$ ;  $p_1 \leftarrow r_0$ ; while  $||r_k|| > \epsilon$  do  $k \leftarrow k + 1$ :

#### Conjugate direction method

$$\alpha_k \leftarrow \frac{\mathbf{r}_{k-1}^T \mathbf{r}_{k-1}}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k};$$
 $\mathbf{x}_k \leftarrow \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k;$ 
 $\mathbf{r}_k \leftarrow \mathbf{r}_{k-1} - \alpha_k \mathbf{A} \mathbf{p}_k;$ 

#### Residual orthogonalization

$$eta_k \leftarrow rac{oldsymbol{r}_k^T oldsymbol{r}_k}{oldsymbol{r}_{k-1}^T oldsymbol{r}_{k-1}}; \ oldsymbol{p}_{k+1} \leftarrow oldsymbol{r}_k + eta_k oldsymbol{p}_k;$$

end while





#### Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



## Polynomial residual expansions

(1/5)

From the Conjugate Gradient iterative scheme on slide 54 we have

#### Lemma

There exists k-degree polynomial  $P_k(x)$  and  $Q_k(x)$  such that

$$r_k = P_k(\mathbf{A})r_0$$

$$k = 0, 1, \dots, n$$

$$p_k = Q_{k-1}(A)r_0$$
  $k = 1, 2, ..., n$ 

Moreover  $P_k(0) = 1$  for all k.

Proof.

(1/2).

The proof is by induction.

Base k = 0

$$p_1 = r_0$$

so that  $P_0(x) = 1$  and  $Q_0(x) = 1$ .



Proof. (2/2).

let the expansion valid for k-1 Consider the recursion for the residual:

$$r_k = r_{k-1} - \alpha_k \mathbf{A} \mathbf{p}_k$$

$$= P_{k-1}(\mathbf{A}) r_0 + \alpha_k \mathbf{A} Q_{k-1}(\mathbf{A}) r_0$$

$$= (P_{k-1}(\mathbf{A}) + \alpha_k \mathbf{A} Q_{k-1}(\mathbf{A})) r_0$$

then  $P_k(x) = P_{k-1}(x) + \alpha_k x Q_{k-1}(x)$  and  $P_k(0) = P_{k-1}(0) = 1$ . Consider the recursion for the conjugate direction

$$p_{k+1} = P_k(\mathbf{A})\mathbf{r}_0 + \beta_k Q_{k-1}(\mathbf{A})\mathbf{r}_0$$
$$= (P_k(\mathbf{A}) + \beta_k Q_{k-1}(\mathbf{A}))\mathbf{r}_0$$

then 
$$Q_k(x) = P_k(x) + \beta_k Q_{k-1}(x)$$
.



(3/5)

## Polynomial residual expansions

We have the following trivial equality

$$\mathcal{V}_k = \text{SPAN}\{p_1, p_2, \dots p_k\}$$

$$= \text{SPAN}\{r_0, r_1, \dots r_{k-1}\}$$

$$= \{q(A)r_0 \mid q \in \mathbb{P}^{k-1}, \}$$

$$= \{p(A)e_0 \mid p \in \mathbb{P}^k, p(0) = 0\}$$

In this way the optimality of CG step can be written as

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} \leq \|\boldsymbol{x}_{\star} - \boldsymbol{x}\|_{\boldsymbol{A}}, \qquad \forall \boldsymbol{x} \in \boldsymbol{x}_{0} + \mathcal{V}_{k}$$

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} \leq \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + p(\boldsymbol{A})\boldsymbol{e}_{0})\|_{\boldsymbol{A}}, \qquad \forall p \in \mathbb{P}^{k}, \ p(0) = 0$$

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} \leq \|P(\boldsymbol{A})\boldsymbol{e}_{0}\|_{\boldsymbol{A}}, \qquad \forall P \in \mathbb{P}^{k}, P(0) = 1$$





## Polynomial residual expansions

Recalling that

$$oldsymbol{A}^{-1}oldsymbol{r}_k = oldsymbol{A}^{-1}(oldsymbol{b} - oldsymbol{A}oldsymbol{x}_k) = oldsymbol{x}_\star - oldsymbol{x}_k = oldsymbol{e}_k$$

we can write

$$e_k = x_* - x_k = A^{-1}r_k$$
  
=  $A^{-1}P_k(A)r_0$   
=  $P_k(A)A^{-1}r_0$   
=  $P_k(A)(x_* - x_0)$   
=  $P_k(A)e_0$ .

due to the optimality of the conjugate gradient we have:





Using the results of slide 58 and 59 we can write

$$e_k = P_k(A)e_0,$$

$$\|e_k\|_{A} = \|P_k(A)e_0\|_{A} \le \|P(A)e_0\|_{A} \qquad \forall P \in \mathbb{P}^k, P(0) = 1$$

and from this equation we have the estimate

$$\|e_k\|_{\boldsymbol{A}} \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \|P(\boldsymbol{A})e_0\|_{\boldsymbol{A}}$$

So an estimate of the form

$$\inf_{P \in \mathbb{P}^k, \ P(0)=1} \left\| P(\boldsymbol{A}) \boldsymbol{e}_0 \right\|_{\boldsymbol{A}} \leq C_k \left\| \boldsymbol{e}_0 \right\|_{\boldsymbol{A}}$$

can be used to proof a convergence rate theorem, as for the steepest descent algorithm.



## Convergence rate calculation

#### Lemma

Let  $\pmb{A} \in \mathbb{R}^{n \times n}$  an SPD matrix, and  $p \in \mathbb{P}^k$  a polynomial, then

$$\|p(\boldsymbol{A})\boldsymbol{x}\|_{\boldsymbol{A}} \le \|p(\boldsymbol{A})\|_2 \|\boldsymbol{x}\|_{\boldsymbol{A}}$$

Proof. (1/2).

The matrix A is SPD so that we can write

$$A = U^T \Lambda U$$
,  $\Lambda = \text{DIAG}\{\lambda_1, \lambda_2, \dots, \lambda_n\}$ 

where U is an orthogonal matrix (i.e.  $U^TU=I$ ) and  $\mathbf{\Lambda} \geq \mathbf{0}$  is diagonal. We can define the SPD matrix  $\mathbf{A}^{1/2}$  as follows

$$A^{1/2} = U^T \mathbf{\Lambda}^{1/2} U, \qquad \mathbf{\Lambda}^{1/2} = \text{DIAG}\{\lambda_1^{1/2}, \lambda_2^{1/2}, \dots, \lambda_n^{1/2}\}$$

and obviously  $oldsymbol{A}^{1/2}oldsymbol{A}^{1/2}=oldsymbol{A}$  .



Proof. (2/2).

Notice that

$$\left\| oldsymbol{x} 
ight\|_{oldsymbol{A}}^2 = oldsymbol{x}^T oldsymbol{A} oldsymbol{x} = oldsymbol{x}^T oldsymbol{A}^{1/2} oldsymbol{x} = \left\| oldsymbol{A}^{1/2} oldsymbol{x} 
ight\|_2^2$$

so that

$$\begin{aligned} \|p(\mathbf{A})\mathbf{x}\|_{\mathbf{A}} &= \left\|\mathbf{A}^{1/2}p(\mathbf{A})\mathbf{x}\right\|_{2} \\ &= \left\|p(\mathbf{A})\mathbf{A}^{1/2}\mathbf{x}\right\|_{2} \\ &\leq \|p(\mathbf{A})\|_{2} \left\|\mathbf{A}^{1/2}\mathbf{x}\right\|_{2} \\ &= \|p(\mathbf{A})\|_{2} \left\|\mathbf{x}\right\|_{\mathbf{A}} \end{aligned}$$





#### Lemma

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix, and  $p \in \mathbb{P}^k$  a polynomial, then

$$||p(\mathbf{A})||_2 = \max_{\lambda \in \sigma(\mathbf{A})} |p(\lambda)|$$

#### Proof.

The matrix p(A) is symmetric, and for a generic symmetric matrix B we have

$$\|\boldsymbol{B}\|_2 = \max_{\lambda \in \sigma(\boldsymbol{B})} |\lambda|$$

observing that if  $\lambda$  is an eigenvalue of A then  $p(\lambda)$  is an eigenvalue of p(A) the thesis easily follows.



Starting the error estimate

$$\|\boldsymbol{e}_{k}\|_{\boldsymbol{A}} \leq \inf_{P \in \mathbb{P}^{k}, P(\mathbf{0})=1} \|P(\boldsymbol{A})\boldsymbol{e}_{\mathbf{0}}\|_{\boldsymbol{A}}$$

Combining the last two lemma we easily obtain the estimate

$$\|e_k\|_{oldsymbol{A}} \leq \inf_{P \in \mathbb{P}^k, \, P(\mathbf{0}) = 1} \left[ \max_{\lambda \in \sigma(oldsymbol{A})} |P(\lambda)| \, \right] \|e_0\|_{oldsymbol{A}}$$

• The convergence rate is estimated by bounding the constant

$$\inf_{P \in \mathbb{P}^k,\, P(\mathbf{0}) = 1} \left[ \max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \, \right]$$





### Finite termination of Conjugate Gradient

#### Theorem (Finite termination of Conjugate Gradient)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix, the the Conjugate Gradient applied to the linear system Ax = b terminate finding the exact solution in at most n-step.

#### Proof.

From the estimate

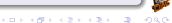
$$\|\boldsymbol{e}_{k}\|_{\boldsymbol{A}} \leq \inf_{P \in \mathbb{P}^{k}, \, P(0)=1} \left[ \max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \, \right] \|\boldsymbol{e}_{0}\|_{\boldsymbol{A}}$$

choosing

$$P(x) = \prod_{\lambda \in \sigma(A)} (x - \lambda) / \prod_{\lambda \in \sigma(A)} (0 - \lambda)$$

we have  $\max_{\lambda \in \sigma(A)} |P(\lambda)| = 0$  and  $||e_n||_A = 0$ .





## Convergence rate of Conjugate Gradient

The constant

$$\inf_{P \in \mathbb{P}^k, P(\mathbf{0}) = 1} \left[ \max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \right]$$

is not easy to evaluate,

2 The following bound, is useful

$$\max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \leq \max_{\lambda \in [\lambda_1, \lambda_n]} |P(\lambda)|$$

in particular the final estimate will be obtained by

$$\inf_{P \in \mathbb{P}^k, \, P(\mathbf{0}) = 1} \left[ \max_{\lambda \in \sigma(\mathbf{A})} |P(\lambda)| \right] \leq \max_{\lambda \in [\lambda_1, \lambda_n]} \left| \bar{P}_k(\lambda) \right|$$

where  $\bar{P}_k(x)$  is an opportune k-degree polynomial for which  $\bar{P}_k(0)=1$  and it is easy to evaluate  $\max_{\lambda\in[\lambda_1,\lambda_2]}|\bar{P}_k(\lambda)|$ .





(1/4)

$$T_k(x) = \cos(k \arccos(x))$$

② Another equivalent definition valid in the interval  $(-\infty,\infty)$  is the following

$$T_k(x) = \frac{1}{2} \left[ \left( x + \sqrt{x^2 - 1} \right)^k + \left( x - \sqrt{x^2 - 1} \right)^k \right]$$

lacksquare In spite of these definition,  $T_k(x)$  is effectively a polynomial.

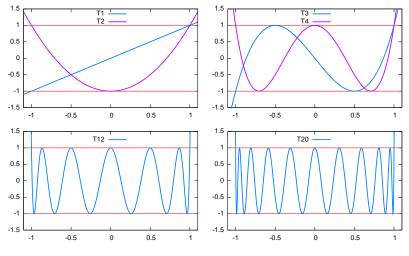




## Chebyshev Polynomials

(2/4)

#### Some example of Chebyshev Polynomials.







## Chebyshev Polynomials

lacksquare It is easy to show that  $T_k(x)$  is a polynomial by the use of

$$\cos(\alpha + \beta) = \cos \alpha \cos \beta - \sin \alpha \sin \beta$$
$$\cos(\alpha + \beta) + \cos(\alpha - \beta) = 2\cos \alpha \cos \beta$$

let  $\theta = \arccos(x)$ :

- **1**  $T_0(x) = \cos(0\,\theta) = 1;$
- **2**  $T_1(x) = \cos(1 \theta) = x;$
- $T_2(x) = \cos(2\theta) = \cos(\theta)^2 \sin(\theta)^2 = 2\cos(\theta)^2 1 = 2x^2 1;$
- $T_{k+1}(x) + T_{k-1}(x) = \cos((k+1)\theta) + \cos((k-1)\theta)$   $= 2\cos(k\theta)\cos(\theta) = 2xT_k(x)$
- 2 In general we have the following recurrence:
  - $\mathbf{0} \ T_0(x) = 1;$
  - $T_1(x) = x;$
  - $T_{k+1}(x) = 2x T_k(x) T_{k-1}(x).$





## Chebyshev Polynomials

- Solving the recurrence:
  - **1**  $T_0(x) = 1$ ;
  - **2**  $T_1(x) = x$ ;
  - 3  $T_{k+1}(x) = 2xT_k(x) T_{k-1}(x)$ .
- We obtain the explicit form of the Chebyshev Polynomials

$$T_k(x) = \frac{1}{2} \left[ \left( x + \sqrt{x^2 - 1} \right)^k + \left( x - \sqrt{x^2 - 1} \right)^k \right]$$

 The translated and scaled polynomial is useful in the study of the conjugate gradient method:

$$T_k(x; a, b) = T_k\left(\frac{a+b-2x}{b-a}\right)$$

where we have  $|T_k(x; a, b)| \le 1$  for all  $x \in [a, b]$ .





## Convergence rate of Conjugate Gradient method

### Theorem (Convergence rate of Conjugate Gradient method)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the Conjugate Gradient method converge to the solution  $x_\star = A^{-1}b$  with at least linear r-rate in the norm  $\|\cdot\|_A$ . Moreover we have the error estimate

$$\|\boldsymbol{e}_k\|_{\boldsymbol{A}} \lesssim 2\left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k \|\boldsymbol{e}_0\|_{\boldsymbol{A}}$$

 $\kappa = M/m$  is the condition number where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.

The expression  $a_k \lesssim b_k$  means that for all  $\epsilon > 0$  there exists  $k_0 > 0$  such that:

$$a_k \le (1 - \epsilon)b_k, \quad \forall k > k_0$$





#### Proof.

From the estimate

$$\|e_k\|_{\boldsymbol{A}} \leq \max_{\lambda \in [m,M]} |P(\lambda)| \|e_0\|_{\boldsymbol{A}}, \qquad P \in \mathbb{P}^k, P(0) = 1$$

choosing  $P(x)=T_k(x;m,M)/T_k(0;m,M)$  from the fact that  $|T_k(x;m,M)|\leq 1$  for  $x\in[m,M]$  we have

$$\|e_k\|_{\mathbf{A}} \le T_k(0; m, M)^{-1} \|e_0\|_{\mathbf{A}} = T_k \left(\frac{M+m}{M-m}\right)^{-1} \|e_0\|_{\mathbf{A}}$$

observe that  $\frac{M+m}{M-m} = \frac{\kappa+1}{\kappa-1}$  and

$$T_k \left(\frac{\kappa+1}{\kappa-1}\right)^{-1} = 2\left[\left(\frac{\sqrt{\kappa}+1}{\sqrt{\kappa}-1}\right)^k + \left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k\right]^{-1}$$

finally notice that  $\left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k \to 0$  as  $k \to \infty$ .



## Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



# Preconditioning

### Problem (Preconditioned linear system)

Given  $A, P \in \mathbb{R}^{n \times n}$ , with A an SPD matrix and P non singular matrix and  $b \in \mathbb{R}^n$ .

Find 
$$x_{\star} \in \mathbb{R}^n$$
 such that:  $P^{-T}Ax_{\star} = P^{-T}b$ .

A good choice for P should be such that  $M = P^T P \approx A$ , where  $\approx$  denotes that M is an approximation of A in some sense to precise later.

Notice that:

P non singular imply:

$$P^{-T}(b-Ax)=\mathbf{0}$$
  $\iff$   $b-Ax=0$ ;

 $oldsymbol{A}$  SPD imply  $\widetilde{A} = oldsymbol{P}^{-T} oldsymbol{A} oldsymbol{P}^{-1}$  is also SPD (obvious proof).



Now we reformulate the preconditioned system:

### Problem (Preconditioned linear system)

Given  $A, P \in \mathbb{R}^{n \times n}$ , with A an SPD matrix and P non singular matrix and  $b \in \mathbb{R}^n$  the preconditioned problem is the following:

Find 
$$\widetilde{m{x}_{\star}} \in \mathbb{R}^n$$
 such that:  $\widetilde{m{A}}\widetilde{m{x}_{\star}} = \widetilde{m{b}}$ 

where

$$\widetilde{A} = P^{-T}AP^{-1}$$
  $\widetilde{b} = P^{-T}b$ 

notice that if  $x_\star$  is the solution of the linear system Ax = b then  $\widetilde{x_\star} = Px_\star$  is the solution of the linear system  $\widetilde{A}x = \widetilde{b}$ .



# PCG: preliminary version

### initial step:

$$\begin{array}{l} k \leftarrow 0; \ \boldsymbol{x}_0 \ \text{assigned}; \\ \widetilde{\boldsymbol{x}}_0 \leftarrow \boldsymbol{P} \boldsymbol{x}_0; \ \widetilde{\boldsymbol{r}}_0 \leftarrow \widetilde{\boldsymbol{b}} - \widetilde{\boldsymbol{A}} \widetilde{\boldsymbol{x}}_0; \ \widetilde{\boldsymbol{p}}_1 \leftarrow \widetilde{\boldsymbol{r}}_0; \\ \text{while } \|\widetilde{\boldsymbol{r}}_k\| > \epsilon \ \text{do} \\ k \leftarrow k+1; \end{array}$$

### Conjugate direction method

$$\widetilde{\alpha}_{k} \leftarrow \frac{\widetilde{r}_{k-1}^{T} \widetilde{r}_{k-1}}{\widetilde{p}_{k}^{T} \widetilde{A} \widetilde{p}_{k}}; 
\widetilde{x}_{k} \leftarrow \widetilde{x}_{k-1} + \widetilde{\alpha}_{k} \widetilde{p}_{k}; 
\widetilde{r}_{k} \leftarrow \widetilde{r}_{k-1} - \widetilde{\alpha}_{k} \widetilde{A} \widetilde{p}_{k};$$

### Residual orthogonalization

$$\widetilde{\beta}_{k} \leftarrow \frac{\widetilde{r}_{k}^{T} \widetilde{r}_{k}}{\widetilde{r}_{k-1}^{T} \widetilde{r}_{k-1}}; 
\widetilde{p}_{k+1} \leftarrow \widetilde{r}_{k} + \widetilde{\beta}_{k} \widetilde{p}_{k};$$

# end while

### final step

$$oldsymbol{P}^{-1}\widetilde{oldsymbol{x}}_{k}$$
;





Conjugate gradient algorithm applied to  $\widetilde{A}\widetilde{x}=\widetilde{b}$  require the evaluation of thing like:

$$\widetilde{A}\widetilde{p}_k = P^{-T}AP^{-1}\widetilde{p}_k.$$

this can be done without evaluate directly the matrix  $\widetilde{A}$ , by the following operations:

- $lacksquare{1}{3}$  solve  $m{P}m{s}_k'=\widetilde{m{p}}_k$  for  $m{s}_k'=m{P}^{-1}\widetilde{m{p}}_k$ ;
- $oldsymbol{2}$  evaluate  $s_k'' = A s_k';$
- $\bullet$  solve  $P^Ts_k''' = s_k''$  for  $s_k''' = P^{-T}s''$ .

Step 1 and 3 require the solution of two auxiliary linear system. This is not a big problem if P and  $P^T$  are triangular matrices (see e.g. incomplete Cholesky).



However... we can reformulate the algorithm using only the matrices A and P!

#### Definition

For all  $k \geq 1$ , we introduce the vector  $q_k = P^{-1}\widetilde{p}$ .

### Observation

If the vectors  $\widetilde{p}_1$ ,  $\widetilde{p}_2$ , ...  $\widetilde{p}_k$  for all  $1 \leq k \leq n$  are  $\widetilde{A}$ -conjugate, then the corresponding vectors  $q_1$ ,  $q_2$ , ...  $q_k$  are A-conjugate. In fact:

$$\mathbf{q}_{j}^{T} \mathbf{A} \mathbf{q}_{i} = \underbrace{\widetilde{\mathbf{p}}_{j}^{T} \mathbf{P}^{-T}}_{=\mathbf{q}_{i}^{T}} \mathbf{A} \underbrace{\mathbf{P}^{-1} \widetilde{\mathbf{p}}_{i}}_{=\mathbf{q}_{j}^{T}} = \widetilde{\mathbf{p}}_{j}^{T} \underbrace{\widetilde{\mathbf{A}}}_{=\mathbf{P}^{-T} \mathbf{A} \mathbf{P}^{-1}}_{=\mathbf{P}^{-T} \mathbf{A} \mathbf{P}^{-1}} \text{ if } i \neq j,$$

that is a consequence of  $\widetilde{A}$ -conjugation of vectors  $\widetilde{p}_i$ .





#### **Definition**

For all  $k \geq 1$ , we introduce the vectors

$$\boldsymbol{x}_k = \boldsymbol{x}_{k-1} + \widetilde{\alpha}_k \boldsymbol{q}_k.$$

### Observation

If we assume, by construction,  $\widetilde{x}_0 = Px_0$ , then we have

$$\widetilde{\boldsymbol{x}}_k = \boldsymbol{P} \boldsymbol{x}_k,$$
 for all  $k$  with  $1 \le k \le n$ .

In fact, if  $\widetilde{x}_{k-1} = Px_{k-1}$  (inductive hypothesis), then

$$egin{aligned} \widetilde{m{x}}_k &= \widetilde{m{x}}_{k-1} + \widetilde{lpha}_k \widetilde{m{p}}_k & [ extit{preconditioned CG}] \ &= m{P}m{x}_{k-1} + \widetilde{lpha}_k m{P}m{q}_k & [ extit{inductive Hyp. defs of }m{q}_k] \ &= m{P}m{(x}_{k-1} + \widetilde{lpha}_k m{q}_k) & [ extit{obvious}] \ &= m{P}m{x}_k & [ extit{defs. of }m{x}_k] \end{aligned}$$



#### Observation

Because  $\widetilde{x}_k = Px_k$  for all  $k \geq 0$ , we have the recurrence between the corresponding residue  $\widetilde{r}_k = \widetilde{b} - \widetilde{A}\widetilde{x}$  and  $r_k = b - Ax_k$ :

$$\widetilde{\boldsymbol{r}}_k = \boldsymbol{P}^{-T} \boldsymbol{r}_k.$$

In fact,

$$egin{aligned} \widetilde{m{r}}_k &= \widetilde{m{b}} - \widetilde{m{A}}\widetilde{m{x}}_k, & [ ext{defs. of } \widetilde{m{r}}_k] \ &= m{P}^{-T}m{b} - m{P}^{-T}m{A}m{P}^{-1}m{P}m{x}_k, & [ ext{defs. of } \widetilde{m{b}}, \ \widetilde{m{A}}, \ \widetilde{m{x}}_k] \ &= m{P}^{-T}m{(b-Am{x}_k)}, & [ ext{obvious}] \ &= m{P}^{-T}m{r}_k. & [ ext{defs. of } m{r}_k] \end{aligned}$$





#### **Definition**

For all k, with  $1 \le k \le n$ , the vector  $z_k$  is the solution of the linear system

$$Mz_k = r_k$$
.

where  $M = P^T P$ . Formally,

$$z_k = M^{-1}r_k = P^{-1}P^{-T}r_k.$$

Using the vectors  $\{z_k\}$ ,

- we can express  $\widetilde{\alpha}_k$  and  $\widetilde{\beta}_k$  in terms of A, the residual  $r_k$ , and conjugate direction  $q_k$ ;
- we can build a recurrence relation for the  $m{A}$ -conjugate directions  $m{q}_k$ .



#### Observation

$$egin{aligned} \widetilde{lpha}_k &= rac{\widetilde{m{r}}_{k-1}^T\widetilde{m{r}}_{k-1}}{\widetilde{m{p}}_k^T\widetilde{m{A}}\widetilde{m{p}}_k} = rac{m{r}_{k-1}m{P}^{-1}m{P}^{-1}m{P}^{-1}m{r}_{k-1}}{m{q}_k^Tm{P}^Tm{P}^{-T}m{A}m{P}^{-1}m{P}m{q}_k} = rac{m{r}_{k-1}m{M}^{-1}m{r}_{k-1}}{m{q}_km{A}m{q}_k}, \ &= \boxed{rac{m{r}_{k-1}m{z}_{k-1}}{m{q}_km{A}m{q}_k}.} \end{aligned}$$

### Observation

$$egin{aligned} \widetilde{eta}_k &= rac{\widetilde{m{r}}_k^T \widetilde{m{r}}_k}{\widetilde{m{r}}_{k-1}^T \widetilde{m{r}}_{k-1}} = rac{m{r}_k^T m{P}^{-1} m{P}^{-1} m{r}_k}{m{r}_{k-1}^T m{P}^{-1} m{P}^{-T} m{r}_{k-1}} = rac{m{r}_k^T m{M}^{-1} m{r}_k}{m{r}_{k-1}^T m{M}^{-1} m{r}_{k-1}}, \ &= \boxed{ rac{m{r}_k^T m{z}_k}{m{r}_{k-1}^T m{z}_{k-1}}. \end{aligned} }$$



#### Observation

Using the vector  $z_k = M^{-1}r_k$ , the following recurrence is true

$$\boldsymbol{q}_{k+1} = \boldsymbol{z}_k + \widetilde{\beta}_k \boldsymbol{q}_k$$

In fact:

$$egin{aligned} \widetilde{p}_{k+1} &= \widetilde{r}_k + \widetilde{eta}_k \widetilde{p}_k & [ ext{preconditioned CG}] \ P^{-1} \widetilde{p}_{k+1} &= P^{-1} \widetilde{r}_k + \widetilde{eta}_k P^{-1} \widetilde{p}_k & [ ext{left mult } P^{-1}] \ P^{-1} \widetilde{p}_{k+1} &= P^{-1} P^{-T} r_k + \widetilde{eta}_k P^{-1} \widetilde{p}_k & [ ext{r}_{k+1} &= P^{-T} r_{k+1}] \ P^{-1} \widetilde{p}_{k+1} &= M^{-1} r_k + \widetilde{eta}_k P^{-1} \widetilde{p}_k & [M^{-1} &= P^{-1} P^{-T}] \ q_{k+1} &= z_k + \widetilde{eta}_k q_k & [q_k &= P^{-1} \widetilde{p}_k] \end{aligned}$$





# PCG: final version

#### initial step:

$$k \leftarrow 0$$
;  $x_0$  assigned;  $r_0 \leftarrow b - Ax_0$ ;  $q_1 \leftarrow r_0$ ; while  $\|z_k\| > \epsilon$  do  $k \leftarrow k + 1$ :

### Conjugate direction method

$$\widetilde{lpha}_k \leftarrow rac{oldsymbol{r}_{k-1}^T oldsymbol{z}_{k-1}}{oldsymbol{q}_k^T \widetilde{oldsymbol{A}} oldsymbol{q}_k}; \ oldsymbol{x}_k \leftarrow oldsymbol{x}_{k-1} + \widetilde{lpha}_k oldsymbol{q}_k;$$

$$r_k \leftarrow r_{k-1} - \widetilde{\alpha}_k A q_k;$$

#### Preconditioning

$$oldsymbol{z}_k = oldsymbol{M}^{-1} oldsymbol{r}_k;$$

#### Residual orthogonalization

$$egin{aligned} \widetilde{eta}_k &\leftarrow rac{oldsymbol{r}_k^Toldsymbol{z}_k}{oldsymbol{r}_{k-1}^Toldsymbol{z}_{k-1}}; \ oldsymbol{q}_{k+1} &\leftarrow oldsymbol{z}_k + \widetilde{eta}_k oldsymbol{q}_k; \end{aligned}$$





## Outline

- Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



# Nonlinear Conjugate Gradient extension

- The conjugate gradient algorithm can be extended for nonlinear minimization.
- ② Fletcher and Reeves extend CG for the minimization of a general non linear function f(x) as follows:
  - **1** Substitute the evaluation of  $\alpha_k$  by an line search
  - **2** Substitute the residual  $r_k$  with the gradient  $\nabla \mathsf{f}(x_k)$
- $oldsymbol{\circ}$  We also translate the index for the search direction  $oldsymbol{p}_k$  to be more consistent with the gradients. The resulting algorithm is in the next slide



# Fletcher and Reeves Nonlinear Conjugate Gradient

### initial step:

$$\begin{array}{l} k \leftarrow 0; \ x_0 \ \text{assigned}; \\ f_0 \leftarrow \mathrm{f}(x_0); \ g_0 \leftarrow \nabla \mathrm{f}(x_0)^T; \\ p_0 \leftarrow -g_0; \\ \text{while} \ \|g_k\| > \epsilon \ \text{do} \\ k \leftarrow k+1; \\ \text{Conjugate direction method} \\ \text{Compute} \ \alpha_k \ \text{by line-search}; \\ x_k \leftarrow x_{k-1} + \alpha_k p_{k-1}; \\ g_k \leftarrow \nabla \mathrm{f}(x_k)^T; \\ \text{Residual orthogonalization} \\ \beta_k^{FR} \leftarrow \frac{g_k^T g_k}{g_{k-1}^T g_{k-1}}; \\ p_k \leftarrow -g_k + \beta_k^{FR} p_{k-1}; \\ \text{end while} \end{array}$$



- lacktriangledown To ensure convergence and apply Zoutendijk global convergence theorem we need to ensure that  $p_k$  is a descent direction.
- $oldsymbol{arphi}$   $p_0$  is a descent direction by construction, for  $p_k$  we have

$$\left\|oldsymbol{g}_k^Toldsymbol{p}_k = -\left\|oldsymbol{g}_k
ight\|^2 + eta_k^{FR}oldsymbol{g}_k^Toldsymbol{p}_{k-1}$$

if the line-search is exact than  $g_k^T p_{k-1} = 0$  because  $p_{k-1}$  is the direction of the line-search. So by induction  $p_k$  is a descent direction.

- Exact line-search is expensive, however if we use inexact line-search with strong Wolfe conditions
  - sufficient decrease:  $f(x_k + \alpha_k p_k) \le f(x_k) + c_1 \alpha_k \nabla f(x_k) p_k$ ;

with  $0 < c_1 < c_2 < 1/2$  then we can prove that  $p_k$  is a descent direction.



The previous consideration permits to say that Fletcher and Reeves nonlinear conjugate gradient method with strong Wolfe line-search is globally convergent<sup>1</sup>

To prove globally convergence we need the following lemma:

## Lemma (descent direction bound)

Suppose we apply Fletcher and Reeves nonlinear conjugate gradient method to f(x) with strong Wolfe line-search with  $0 < c_2 < 1/2$ . The the method generates descent direction  $p_k$  that satisfy the following inequality

$$-\frac{1}{1-c_2} \le \frac{m{g}_k^T m{p}_k}{\|m{g}_k\|^2} \le -\frac{1-2c_2}{1-c_2}, \qquad k = 0, 1, 2, \dots$$



<sup>&</sup>lt;sup>1</sup>globally here means that Zoutendijk like theorem apply (7) + (3) + (3) + (3)

Proof. (1/3).

The proof is by induction. First notice that the function

$$t(\xi) = \frac{2\xi - 1}{1 - \xi}$$

is monotonically increasing on the interval [0,1/2] and that t(0)=-1 and t(1/2)=0. Hence, because of  $c_2\in(0,1/2)$  we have:

$$-1 < \frac{2c_2 - 1}{1 - c_2} < 0. \tag{*}$$

base of induction k = 0: For k = 0 we have  $p_0 = -g_0$  so that  $g_0^T p_0 / ||g_0||^2 = -1$ . From  $(\star)$  the lemma inequality is trivially satisfied.





Proof. (2/3).

Using update direction formula's of the algorithm:

$$eta_k^{FR} = rac{oldsymbol{g}_k^Toldsymbol{g}_k}{oldsymbol{g}_{k-1}^Toldsymbol{g}_{k-1}} \qquad oldsymbol{p}_k = -oldsymbol{g}_k + eta_k^{FR}oldsymbol{p}_{k-1}$$

we can write

$$\frac{{{\boldsymbol{g}}_k^T{\boldsymbol{p}_k}}}{{{{\left\| {{\boldsymbol{g}}_k} \right\|}^2}}} = - 1 + \beta _k^{FR}\frac{{{\boldsymbol{g}}_k^T{\boldsymbol{p}_{k - 1}}}}{{{{\left\| {{\boldsymbol{g}}_k} \right\|}^2}}} = - 1 + \frac{{{\boldsymbol{g}}_k^T{\boldsymbol{p}_{k - 1}}}}{{{{\left\| {{\boldsymbol{g}}_{k - 1}} \right\|}^2}}}$$

and by using second strong Wolfe condition:

$$-1 + c_2 \frac{\boldsymbol{g}_{k-1}^T \boldsymbol{p}_{k-1}}{\left\| \boldsymbol{g}_{k-1} \right\|^2} \le \frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\left\| \boldsymbol{g}_k \right\|^2} \le -1 - c_2 \frac{\boldsymbol{g}_{k-1}^T \boldsymbol{p}_{k-1}}{\left\| \boldsymbol{g}_{k-1} \right\|^2}$$





Proof. (3/3).

by induction we have

$$rac{1}{1-c_2} \geq -rac{oldsymbol{g}_{k-1}^Toldsymbol{p}_{k-1}}{\left\|oldsymbol{g}_{k-1}
ight\|^2} > 0$$

so that

$$\frac{\boldsymbol{g}_{k}^{T}\boldsymbol{p}_{k}}{\left\|\boldsymbol{g}_{k}\right\|^{2}} \leq -1 - c_{2} \frac{\boldsymbol{g}_{k-1}^{T}\boldsymbol{p}_{k-1}}{\left\|\boldsymbol{g}_{k-1}\right\|^{2}} \leq -1 + c_{2} \frac{1}{1 - c_{2}} = \frac{2c_{2} - 1}{1 - c_{2}}$$

and

$$\left\| rac{oldsymbol{g}_{k}^{T}oldsymbol{p}_{k}}{\left\| oldsymbol{g}_{k} 
ight\|^{2}} \geq -1 + c_{2}rac{oldsymbol{g}_{k-1}^{T}oldsymbol{p}_{k-1}}{\left\| oldsymbol{g}_{k-1} 
ight\|^{2}} \geq -1 - c_{2}rac{1}{1-c_{2}} = -rac{1}{1-c_{2}}$$





The inequality of the the previous lemma can be written as:

$$\frac{1}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} \ge - \frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{g}_k\| \|\boldsymbol{p}_k\|} \ge \frac{1 - 2c_2}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} > 0$$

Remembering the Zoutendijk theorem we have

$$\sum_{k=1}^{\infty} (\cos \theta_k)^2 \|\boldsymbol{g}_k\|^2 < \infty, \quad \text{where} \quad \cos \theta_k = -\frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{g}_k\| \|\boldsymbol{p}_k\|}$$

- **3** so that if  $\|g_k\|/\|p_k\|$  is bounded from below we have that  $\cos \theta_k \geq \delta$  for all k and then from Zoutendijk theorem the scheme converge.
- ① Unfortunately this bound cant be proved so that Zoutendijk theorem cant be applied directly. However it is possible to prove a weaker results, i.e. that  $\liminf_{k\to\infty}\|g_k\|=0!$





# Convergence of Fletcher and Reeves method

## Assumption (Regularity assumption)

We assume  $f \in C^1(\mathbb{R}^n)$  with Lipschitz continuous gradient, i.e. there exists  $\gamma > 0$  such that

$$\|\nabla f(\boldsymbol{x})^T - \nabla f(\boldsymbol{y})^T\| \le \gamma \|\boldsymbol{x} - \boldsymbol{y}\|, \quad \forall \boldsymbol{x}, \boldsymbol{y} \in \mathbb{R}^n$$



## Theorem (Convergence of Fletcher and Reeves method)

Suppose the method of Fletcher and Reeves is implemented with strong Wolfe line-search with  $0 < c_1 < c_2 < 1/2$ . If f(x) and  $x_0$  satisfy the previous regularity assumptions, then

$$\liminf_{k \to \infty} \|\boldsymbol{g}_k\| = 0$$

# Proof. (1/4).

From previous Lemma we have

$$\cos \theta_k \ge \frac{1}{1 - c_2} \frac{\|g_k\|}{\|p_k\|}$$
  $k = 1, 2, ...$ 

substituting in Zoutendijk condition we have  $\sum_{k=1}^{\infty} \frac{\|g_k\|^4}{\|p_k\|^2} < \infty.$ 

The proof is by contradiction. in fact if theorem is not true than the series diverge. Next we want to bound  $||p_k||$ .



## Proof. (bounding $\|p_k\|$ )

(2/4).

Using second Wolfe condition and previous Lemma

$$\left| \boldsymbol{g}_{k}^{T} \boldsymbol{p}_{k-1} \right| \leq -c_{2} \boldsymbol{g}_{k}^{T} \boldsymbol{p}_{k-1} \leq \frac{c_{2}}{1-c_{2}} \left\| \boldsymbol{g}_{k-1} \right\|^{2}$$

using  $oldsymbol{p}_k \leftarrow -oldsymbol{g}_k + eta_k^{FR} oldsymbol{p}_{k-1}$  we have

$$\|\mathbf{p}_{k}\|^{2} \leq \|\mathbf{g}_{k}\|^{2} + 2\beta_{k}^{FR} |\mathbf{g}_{k}^{T}\mathbf{p}_{k-1}| + (\beta_{k}^{FR})^{2} \|\mathbf{p}_{k-1}\|^{2}$$

$$\leq \|\mathbf{g}_{k}\|^{2} + \frac{2c_{2}}{1 - c_{2}} \beta_{k}^{FR} \|\mathbf{g}_{k-1}\|^{2} + (\beta_{k}^{FR})^{2} \|\mathbf{p}_{k-1}\|^{2}$$

recall that  $eta_k^{FR} \leftarrow \left\|oldsymbol{g}_k 
ight\|^2 / \left\|oldsymbol{g}_{k-1} 
ight\|^2$  then

$$\left\lVert oldsymbol{p}_k 
ight
Vert^2 \leq rac{1+c_2}{1-c_2} \left\lVert oldsymbol{g}_k 
ight
Vert^2 + (eta_k^{FR})^2 \left\lVert oldsymbol{p}_{k-1} 
ight
Vert^2$$





# Proof. (bounding $\|\boldsymbol{p}_k\|$ )

(3/4).

setting  $c_3 = \frac{1+c_2}{1-c_2}$  and using repeatedly the last inequality we obtain:

$$\begin{split} \left\| \boldsymbol{p}_{k} \right\|^{2} & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{2} + (\beta_{k}^{FR})^{2} (c_{3} \left\| \boldsymbol{g}_{k-1} \right\|^{2} + (\beta_{k-1}^{FR})^{2} \left\| \boldsymbol{p}_{k-2} \right\|^{2}) \\ & = c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \left( \left\| \boldsymbol{g}_{k} \right\|^{-2} + \left\| \boldsymbol{g}_{k-1} \right\|^{-2} \right) + \frac{\left\| \boldsymbol{g}_{k} \right\|^{4}}{\left\| \boldsymbol{g}_{k-2} \right\|^{4}} \left\| \boldsymbol{p}_{k-2} \right\|^{2} \\ & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \left( \left\| \boldsymbol{g}_{k} \right\|^{-2} + \left\| \boldsymbol{g}_{k-1} \right\|^{-2} + \left\| \boldsymbol{g}_{k-2} \right\|^{-2} \right) \\ & + \frac{\left\| \boldsymbol{g}_{k} \right\|^{4}}{\left\| \boldsymbol{g}_{k-3} \right\|^{4}} \left\| \boldsymbol{p}_{k-3} \right\|^{2} \\ & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \sum_{i=1}^{k} \left\| \boldsymbol{g}_{i} \right\|^{-2} \end{split}$$



### Proof.

(4/4).

Suppose now by contradiction there exists  $\delta>0$  such that  $\|g_k\|\geq \delta$  a by using the regularity assumptions we have

$$\|\boldsymbol{p}_k\|^2 \le c_3 \|\boldsymbol{g}_k\|^4 \sum_{j=1}^k \|\boldsymbol{g}_j\|^{-2} \le c_3 \|\boldsymbol{g}_k\|^4 \delta^{-2} k$$

Substituting in Zoutendijk condition we have

$$\infty > \sum_{k=1}^{\infty} \frac{\|g_k\|^4}{\|p_k\|^2} \ge \frac{\delta^2}{c_4} \sum_{k=1}^{\infty} \frac{1}{k} = \infty$$

this contradict assumption.





<sup>&</sup>lt;sup>a</sup>the correct assumption is that there exists  $k_0$  such that  $\|g_k\| \geq \delta$  for  $k \geq k_0$  but this complicate a little bit the following inequality without introducing new idea.

### Weakness of Fletcher and Reeves method

- Suppose that  $p_k$  is a bad search direction, i.e.  $\cos \theta_k \approx 0$ .
- From the descent direction bound Lemma (see slide 89) we have

$$\frac{1}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} \ge \cos \theta_k \ge \frac{1 - 2c_2}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} > 0$$

- so that to have  $\cos \theta_k \approx 0$  we needs  $\|\boldsymbol{p}_k\| \gg \|\boldsymbol{g}_k\|$ .
- since  $p_k$  is a bad direction near orthogonal to  $g_k$  it is likely that the step is small and  $x_{k+1} \approx x_k$ . If so we have also  $g_{k+1} \approx g_k$  and  $\beta_{k+1}^{FR} \approx 1$ .
- but remember that  $m{p}_{k+1} \leftarrow -m{g}_{k+1} + eta_{k+1}^{FR} m{p}_k$ , so that  $m{p}_{k+1} pprox m{p}_k$ .
- This means that a long sequence of unproductive iterates will follows.



# Polack and Ribiére Nonlinear Conjugate Gradient

- The previous problem can be elided if we restart anew when the iterate stagnate.
- 2 Restarting is obtained by simply set  $\beta_k^{FR} = 0$ .
- **3** A more elegant solution can be obtained with a new definition of  $\beta_k$  due to Polack and Ribiére is the following:

$$eta_k^{PR} = rac{oldsymbol{g}_k^T(oldsymbol{g}_k - oldsymbol{g}_{k-1})}{oldsymbol{g}_{k-1}^Toldsymbol{g}_{k-1}}$$

• This definition of  $\beta_k^{PR}$  is identical of  $\beta_k^{FR}$  in the case of quadratic function because  $\boldsymbol{g}_k^T\boldsymbol{g}_{k-1}=0$ . The definition differs in non linear case and in particular when there is stagnation i.e.  $\boldsymbol{g}_k \approx \boldsymbol{g}_{k-1}$  we have  $\beta_k^{PR} \approx 0$ , i.e. we have an automatic restart.



# Polack and Ribiére Nonlinear Conjugate Gradient

#### initial step:

$$\begin{array}{l} k \leftarrow 0; \ x_0 \ \text{assigned}; \\ f_0 \leftarrow \mathsf{f}(x_0); \ g_0 \leftarrow \nabla \mathsf{f}(x_0)^T; \\ p_0 \leftarrow -g_0; \\ \textbf{while} \ \|g_k\| > \epsilon \ \textbf{do} \\ k \leftarrow k+1; \\ \textbf{Conjugate direction method} \\ \textbf{Compute} \ \alpha_k \ \text{by line-search}; \\ x_k \leftarrow x_{k-1} + \alpha_k p_{k-1}; \\ g_k \leftarrow \nabla \mathsf{f}(x_k)^T; \\ \textbf{Residual orthogonalization} \\ \beta_k^{PR} \leftarrow \frac{g_k^T(g_k - g_{k-1})}{g_{k-1}^T g_{k-1}}; \\ p_k \leftarrow -g_k + \beta_k^{PR} p_{k-1}; \\ \textbf{end while} \end{array}$$

## Weakness of Polack and Ribiére method

- Although the modification is minimal, for the Polack and Ribiére method with strong Wolfe line-search it can happen that  $p_k$  is not a descent direction.
- If  $p_k$  is not a descent direction we can restart i.e. set  $\beta_k^{PR}=0$  or modify  $\beta_k^{PR}$  as follows

$$\beta_k^{PR+} = \max\{\beta_k^{PR}, 0\}$$

this new coefficient with a modified Wolfe line-search ensure that  $p_k$  is a descent direction.



(2/2)

- Polack and Ribiére choice on the average perform better than Fletcher and Reeves but there is not convergence results!
- Although there is not convergence results there is a negative results due to Powell:

#### Theorem

Consider the Polack and Ribiére method with exact line-search. There exists a twice continuously differentiable function  $f: \mathbb{R}^3 \mapsto \mathbb{R}$  and a starting point  $x_0$  such that the sequence of gradients  $\{\|g_k\|\}$  is bounded away from zero.

• However is spite of this results Polack and Ribiére is the first choice among conjugate direction methods.



### Other choices

• There are many other modification of the coefficient  $\beta_k$  that collapse to the same coefficient in the case o quadratic function. One important choice is the Hestenes and Stiefel choice

$$eta_k^{HS} = rac{oldsymbol{g}_k^T(oldsymbol{g}_k - oldsymbol{g}_{k-1})}{(oldsymbol{g}_k^T - oldsymbol{g}_{k-1}^T)oldsymbol{p}_{k-1}}$$

• For this choice there is similar convergence results of Fletcher and Reeves and similar performance.





### References



J. E. Dennis, Jr. and Robert B. Schnabel Numerical Methods for Unconstrained Optimization and Nonlinear Equations SIAM, Classics in Applied Mathematics, 16, 1996.



J. Nocedal and S. J. Wrigth Numerical Optimization Springer Series in Operation Research, 1999.



J. Stoer and R. Bulirsch Introduction to numerical analysis Springer-Verlag, Texts in Applied Mathematics, 12, 2002.

