

Conjugate Direction minimization

Lectures for PHD course on
Non-linear equations and numerical optimization

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Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- 2 Conjugate direction method
- 3 Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



Generic minimization algorithm

In the following we study the convergence rate of the Generic minimization algorithm applied to a quadratic function $q(\mathbf{x})$ with **exact** line search. The function

$$q(\mathbf{x}) = \frac{1}{2} \mathbf{x}^T \mathbf{A} \mathbf{x} - \mathbf{b}^T \mathbf{x} + c$$

can be viewed as a n -dimensional generalization of the 1-dimensional parabolic model.

Generic minimization algorithm

Given an initial guess \mathbf{x}_0 , let $k = 0$;

while not converged do

Find a descent direction \mathbf{p}_k at \mathbf{x}_k ;

Compute a step size α_k using a line-search along \mathbf{p}_k .

Set $\mathbf{x}_{k+1} = \mathbf{x}_k + \alpha_k \mathbf{p}_k$ and increase k by 1.

end while



Assumption (Symmetry)

The matrix \mathbf{A} is assumed to be symmetric, in fact,

$$\mathbf{A} = \mathbf{A}^{\text{Symm}} + \mathbf{A}^{\text{Skew}}$$

where

$$\mathbf{A}^{\text{Symm}} = \frac{1}{2} [\mathbf{A} + \mathbf{A}^T], \quad \mathbf{A}^{\text{Symm}} = (\mathbf{A}^{\text{Symm}})^T$$

$$\mathbf{A}^{\text{Skew}} = \frac{1}{2} [\mathbf{A} - \mathbf{A}^T], \quad \mathbf{A}^{\text{Skew}} = -(\mathbf{A}^{\text{Skew}})^T$$

moreover

$$\mathbf{x}^T \mathbf{A} \mathbf{x} = \mathbf{x}^T \mathbf{A}^{\text{Symm}} \mathbf{x} + \mathbf{x}^T \mathbf{A}^{\text{Skew}} \mathbf{x} = \mathbf{x}^T \mathbf{A}^{\text{Symm}} \mathbf{x}$$

so that only the symmetric part of \mathbf{A} contribute to $q(\mathbf{x})$.



Assumption (SPD)

The matrix A is assumed to be symmetric and positive definite, in fact,

$$\nabla q(x)^T = \frac{1}{2}(A + A^T)x - b = Ax - b$$

and

$$\nabla^2 q(x) = \frac{1}{2}(A + A^T) = A$$

From the **sufficient** condition for a minimum we have that $\nabla q(x_*)^T = 0$, i.e.

$$Ax_* = b$$

and $\nabla^2 q(x_*) = A$ is SPD.



The toy problem

(1/3)

- In the following we study the convergence rate of the Steepest Descent and Conjugate Gradient methods applied to

$$q(x) = \frac{1}{2}x^T Ax - b^T x + c$$

where A is an SPD matrix.

- This assumption simplify the analysis but it is also useful in the non linear case. In fact, by expanding a generic function $f(x)$ near its minimum x_* we have

$$\begin{aligned} f(x) &= f(x_*) + \nabla f(x_*)(x - x_*) \\ &\quad + \frac{1}{2}(x - x_*)^T \nabla^2 f(x_*)(x - x_*) + \mathcal{O}(\|x - x_*\|^3) \end{aligned}$$



The toy problem

(2/3)

- By setting

$$A = \nabla^2 f(x_*),$$

$$b = \nabla^2 f(x_*)x_* - \nabla f(x_*)$$

$$c = f(x_*) - \nabla f(x_*)x_* + \frac{1}{2}x_*^T \nabla^2 f(x_*)x_*$$

we have

$$f(x) = \frac{1}{2}x^T Ax - b^T x + c + \mathcal{O}(\|x - x_*\|^3)$$

- So that we expect that when an iterate x_k is near x_* then we can neglect $\mathcal{O}(\|x - x_*\|^3)$ and the asymptotic behavior is the same of the quadratic problem.



The toy problem

(3/3)

- we can rewrite the quadratic problem in many different way as follows

$$\begin{aligned} q(x) &= \frac{1}{2}(x - x_*)^T A(x - x_*) + c' \\ &= \frac{1}{2}(Ax - b)^T A^{-1}(Ax - b) + c' \end{aligned}$$

where

$$c' = c + \frac{1}{2}x_*^T Ax_*$$

- This last forms are useful in the study of the steepest descent method.



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The steepest descent for quadratic functions

(1/3)

The steepest descent minimization algorithm

Given an initial guess x_0 , let $k = 0$;

while not converged do

Choose as descent direction $p_k = -\nabla q(x_k)^T = b - Ax_k$;

Compute a step size α_k using a line-search along p_k .

Set $x_{k+1} = x_k + \alpha_k p_k$ and increase k by 1.

end while

Definition (Residual)

The expressions

$$r(x) = b - Ax, \quad r_k = b - Ax_k$$

are called the residual. We obviously have $r(x) = -\nabla q(x)^T$ and $r(x_*) = 0$.

The steepest descent for quadratic functions

(2/3)

We can solve exactly the problem

$$\alpha_k = \arg \min_{\alpha \geq 0} q(x_k - \alpha r_k)$$

because $p(\alpha) = q(x_k - \alpha r_k)$ is a parabola. In fact

$$\frac{dp(\alpha)}{d\alpha} = \frac{dq(x_k - \alpha r_k)}{d\alpha} = -\nabla q(x_k - \alpha r_k) r_k = 0$$

but

$$\begin{aligned} 0 &= -\nabla q(x_k - \alpha r_k) r_k = r(x_k - \alpha r_k)^T r_k = (b - A(x_k - \alpha r_k))^T r_k \\ &= (r_k - \alpha A r_k)^T r_k \end{aligned}$$

and the minimum is at α set to $\frac{r_k^T r_k}{r_k^T A r_k}$.

The steepest descent for quadratic functions

(3/3)

The steepest descent minimization algorithm

Given an initial guess x_0 , let $k = 0$;

while not converged do

Compute $r_k = b - Ax_k$;

Compute the step size $\alpha_k = \frac{r_k^T r_k}{r_k^T A r_k}$;

Set $x_{k+1} = x_k + \alpha_k r_k$ and increase k by 1.

end while

Or more compactly

$$x_{k+1} = x_k + \frac{r_k^T r_k}{r_k^T A r_k} r_k$$

The steepest descent reduction step

(1/3)

We want bound $q(x_{k+1})$ by $q(x_k)$:

$$\begin{aligned} q(x_{k+1}) &= q(x_k + \alpha_k r_k) \\ &= \frac{1}{2} (Ax_k + \alpha_k Ar_k - b)^T A^{-1} (Ax_k + \alpha_k Ar_k - b) + c' \\ &= \frac{1}{2} (\alpha_k Ar_k - r_k)^T A^{-1} (\alpha_k Ar_k - r_k) + c' \\ &= \frac{1}{2} r_k^T A^{-1} r_k + \frac{1}{2} \alpha_k^2 r_k^T Ar_k - \alpha_k r_k^T r_k + c' \\ &= q(x_k) + \frac{1}{2} \alpha_k (\alpha_k r_k^T Ar_k - 2r_k^T r_k) \end{aligned}$$



The steepest descent reduction step

(3/3)

or better

$$r_{k+1}^T A^{-1} r_{k+1} = r_k^T A^{-1} r_k \left(1 - \frac{(r_k^T r_k)^2}{(r_k^T A^{-1} r_k)(r_k^T Ar_k)} \right)$$

noticing that $r_k = b - Ax_k = Ax_* - Ax_k = A(x_* - x_k)$ we have

$$\|x_* - x_{k+1}\|_A^2 = \|x_* - x_k\|_A^2 \left(1 - \frac{(r_k^T r_k)^2}{(r_k^T A^{-1} r_k)(r_k^T Ar_k)} \right)$$

where

$$\|x\|_A = \sqrt{x^T A x}$$

is the **energy norm** induced by the SPD matrix A .



The steepest descent reduction step

(2/3)

Substituting $\alpha_k = \frac{r_k^T r_k}{r_k^T Ar_k}$ we obtain

$$q(x_{k+1}) = q(x_k) - \frac{1}{2} \frac{(r_k^T r_k)^2}{r_k^T Ar_k}$$

this shows that the steepest descent method reduce at each step the objective function $q(x)$.

Using the expression $q(x) = \frac{1}{2} r(x)^T A^{-1} r(x) + c'$ we can write:

$$\frac{1}{2} r_{k+1}^T A^{-1} r_{k+1} = \frac{1}{2} r_k^T A^{-1} r_k - \frac{1}{2} \frac{(r_k^T r_k)^2}{r_k^T Ar_k}$$



The estimate of the convergence rate for the **steepest descent** method is linked to the estimate of the term

$$\frac{(r_k^T r_k)^2}{(r_k^T A^{-1} r_k)(r_k^T Ar_k)}$$

in particular we can prove

Lemma (Kantorovic)

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix then the following inequality is valid

$$1 \leq \frac{(x^T A x)(x^T A^{-1} x)}{(x^T x)^2} \leq \frac{(M+m)^2}{4 M m}$$

for all $x \neq 0$. Where $m = \lambda_1$ is the smallest eigenvalue of A and $M = \lambda_n$ is the biggest eigenvalue of A .



Proof.

(1/5).

STEP 1: problem reformulation. First of all notice that

$$\frac{(x^T A x)(x^T A^{-1} x)}{(x^T x)^2} = \frac{(y^T A y)(y^T A^{-1} y)}{(y^T y)^2}$$

for all $y = \alpha x$ with $\alpha \neq 0$. Choosing $\alpha = \|x\|^{-1}$ have:

$$\begin{aligned} \min_{\|z\|=1} (z^T A z)(z^T A^{-1} z) &\leq \\ \frac{(x^T A x)(x^T A^{-1} x)}{(x^T x)^2} &\leq \max_{\|z\|=1} (z^T A z)(z^T A^{-1} z) \end{aligned}$$



Proof.

(2/5).

STEP 2: eigenvector expansions. Matrix $A \in \mathbb{R}^{n \times n}$ is an SPD matrix so that there exists u_1, u_2, \dots, u_n a complete orthonormal eigenvectors set with $0 < \lambda_1 \leq \lambda_2 \leq \dots \leq \lambda_n$ corresponding eigenvalues. Let be $x \in \mathbb{R}^n$ then

$$x = \sum_{k=1}^n \alpha_k u_k, \quad x^T x = \sum_{k=1}^n \alpha_k^2$$

so that $(x^T A x)(x^T A^{-1} x) = h(\alpha_1, \dots, \alpha_n)$ where

$$h(\alpha_1, \dots, \alpha_n) = \left(\sum_{k=1}^n \alpha_k^2 \lambda_k \right) \left(\sum_{k=1}^n \alpha_k^2 \lambda_k^{-1} \right)$$

then the lemma can be reformulated:

- Find maxima and minima of $h(\alpha_1, \dots, \alpha_n)$
- subject to $\sum_{k=1}^n \alpha_k^2 = 1$.



Proof.

(3/5).

STEP 3: problem reduction. By using Lagrange multiplier maxima and minima are the stationary points of:

$$g(\alpha_1, \dots, \alpha_n, \mu) = h(\alpha_1, \dots, \alpha_n) + \mu \left(\sum_{k=1}^n \alpha_k^2 - 1 \right)$$

setting $A = \sum_{k=1}^n \alpha_k^2 \lambda_k$ and $B = \sum_{k=1}^n \alpha_k^2 \lambda_k^{-1}$ we have

$$\frac{\partial g(\alpha_1, \dots, \alpha_n, \mu)}{\partial \alpha_k} = 2\alpha_k (\lambda_k B + \lambda_k^{-1} A + \mu) = 0$$

so that

- Or $\alpha_k = 0$;
 - Or λ_k is a root of the quadratic polynomial $\lambda^2 B + \lambda \mu + A$.
- in any case there are at most 2 coefficients α 's not zero. ^a

^athe argument should be improved in the case of multiple eigenvalues

Proof.

(4/5).

STEP 4: problem reformulation. say α_i and α_j are the only non zero coefficients, then $\alpha_i^2 + \alpha_j^2 = 1$ and we can write

$$\begin{aligned} h(\alpha_1, \dots, \alpha_n) &= (\alpha_i^2 \lambda_i + \alpha_j^2 \lambda_j) (\alpha_i^2 \lambda_i^{-1} + \alpha_j^2 \lambda_j^{-1}) \\ &= \alpha_i^4 + \alpha_j^4 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i} \right) \\ &= \alpha_i^2 (1 - \alpha_j^2) + \alpha_j^2 (1 - \alpha_i^2) + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i} \right) \\ &= 1 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i} - 2 \right) \\ &= 1 + \alpha_i^2 (1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_j} \end{aligned}$$



Proof.

(5/5).

STEP 5: bounding maxima and minima. notice that

$$0 \leq \beta(1 - \beta) \leq \frac{1}{4}, \quad \forall \beta \in [0, 1]$$

$$1 \leq 1 + \alpha_i^2(1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_j} \leq 1 + \frac{(\lambda_i - \lambda_j)^2}{4\lambda_i \lambda_j} = \frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j}$$

to bound $(\lambda_i + \lambda_j)^2 / (4\lambda_i \lambda_j)$ consider the function $f(x) = (1 + x)^2 / x$ which is increasing for $x \geq 1$ so that we have

$$\frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j} \leq \frac{(M + m)^2}{4 M m}$$

and finally

$$1 \leq h(\alpha_1, \dots, \alpha_n) \leq \frac{(M + m)^2}{4 M m}$$



Proof.

Remember from slide N°15

$$\|x_* - x_{k+1}\|_A^2 = \|x_* - x_k\|_A^2 \left(1 - \frac{(r_k^T r_k)^2}{(r_k^T A^{-1} r_k)(r_k^T A r_k)} \right)$$

from Kantorovich inequality

$$1 - \frac{(r_k^T r_k)^2}{(r_k^T A^{-1} r_k)(r_k^T A r_k)} \leq 1 - \frac{4 M m}{(M + m)^2} = \frac{(M - m)^2}{(M + m)^2}$$

so that

$$\|x_* - x_{k+1}\|_A \leq \frac{M - m}{M + m} \|x_* - x_k\|_A$$



Convergence rate of Steepest Descent

The Kantorovich inequality permits to prove:

Theorem (Convergence rate of Steepest Descent)

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix then the **steepest descent** method:

$$x_{k+1} = x_k + \frac{r_k^T r_k}{r_k^T A r_k} r_k$$

converge to the solution $x_* = A^{-1}b$ with at least linear q -rate in the norm $\|\cdot\|_A$. Moreover we have the error estimate

$$\|x_{k+1} - x_*\|_A \leq \frac{\kappa - 1}{\kappa + 1} \|x_k - x_*\|_A$$

 $\kappa = M/m$ is the **condition number** where $m = \lambda_1$ is the smallest eigenvalue of A and $M = \lambda_n$ is the biggest eigenvalue of A .

Remark (One step convergence)

The steepest descent method can converge in one iteration if $\kappa = 1$ or when $r_0 = u_k$ where u_k is an eigenvector of A .

- ④ In the first case ($\kappa = 1$) we have $A = \beta I$ for some $\beta > 0$ so it is not interesting.
- ④ In the second case we have

$$\frac{(u_k^T u_k)^2}{(u_k^T A^{-1} u_k)(u_k^T A u_k)} = \frac{(u_k^T u_k)^2}{\lambda_k^{-1} (u_k^T u_k) \lambda_k (u_k^T u_k)} = 1$$

in both cases we have $r_1 = 0$ i.e. we have found the solution.

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- 1 Convergence rate of Steepest Descent iterative scheme
- 2 **Conjugate direction method**
- 3 Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
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Conjugate direction method

Definition (Conjugate vector)

Given two vectors p and q in \mathbb{R}^n are **conjugate respect to A** if they are orthogonal respect the scalar product induced by A ; i.e.,

$$p^T A q = \sum_{i,j=1}^n A_{ij} p_i q_j = 0.$$

Clearly, n vectors $p_1, p_2, \dots, p_n \in \mathbb{R}^n$ that are pair wise conjugated respect to A form a base of \mathbb{R}^n .

Problem (Linear system)

Find the minimum of $q(x) = \frac{1}{2}x^T A x - b^T x + c$ is equivalent to solve the first order necessary condition, i.e.

$$\text{Find } x_* \in \mathbb{R}^n \text{ such that: } Ax_* = b.$$

Observation

Consider $x_0 \in \mathbb{R}^n$ and decompose the error $e_0 = x_* - x_0$ by the conjugate vectors $p_1, p_2, \dots, p_n \in \mathbb{R}^n$:

$$e_0 = x_* - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \dots + \sigma_n p_n.$$

Evaluating the coefficients $\sigma_1, \sigma_2, \dots, \sigma_n \in \mathbb{R}$ is equivalent to solve the problem $Ax_* = b$, because knowing e_0 we have

$$x_* = x_0 + e_0.$$

Observation

Using conjugacy the coefficients $\sigma_1, \sigma_2, \dots, \sigma_n \in \mathbb{R}$ can be computed as

$$\sigma_i = \frac{p_i^T A e_0}{p_i^T A p_i}, \quad \text{for } i = 1, 2, \dots, n.$$

In fact, for all $1 \leq i \leq n$, we have

$$\begin{aligned} p_i^T A e_0 &= p_i^T A (\sigma_1 p_1 + \sigma_2 p_2 + \dots + \sigma_n p_n), \\ &= \sigma_1 p_i^T A p_1 + \sigma_2 p_i^T A p_2 + \dots + \sigma_n p_i^T A p_n, \\ &= \sigma_i p_i^T A p_i, \end{aligned}$$

because $p_i^T A p_j = 0$ for $i \neq j$.

The conjugate direction method evaluates the coefficients $\sigma_1, \sigma_2, \dots, \sigma_n \in \mathbb{R}$ recursively in n steps, solving for $k \geq 0$ the minimization problem:

Conjugate direction method

Given x_0 ; $k \leftarrow 0$;

repeat

$k \leftarrow k + 1$;

Find $x_k \in x_0 + \mathcal{V}_k$ such that:

$$x_k = \arg \min_{x \in x_0 + \mathcal{V}_k} \|x_* - x\|_A$$

until $k = n$

where \mathcal{V}_k is the subspace of \mathbb{R}^n generated by the first k conjugate direction; i.e.,

$$\mathcal{V}_k = \text{SPAN}\{p_1, p_2, \dots, p_k\}.$$



Solving first step method 1

The minimization problem is the minimum respect to α of the quadratic:

$$\begin{aligned} \Phi(\alpha) &= \|x_* - (x_0 + \alpha p_1)\|_A^2, \\ &= (x_* - (x_0 + \alpha p_1))^T A (x_* - (x_0 + \alpha p_1)), \\ &= (e_0 - \alpha p_1)^T A (e_0 - \alpha p_1), \\ &= e_0^T A e_0 - 2\alpha p_1^T A e_0 + \alpha^2 p_1^T A p_1. \end{aligned}$$

minimum is found by imposing:

$$\frac{d\Phi(\alpha)}{d\alpha} = -2p_1^T A e_0 + 2\alpha p_1^T A p_1 = 0 \quad \Rightarrow$$

$$\alpha_1 = \frac{p_1^T A e_0}{p_1^T A p_1}$$



Step: $x_0 \rightarrow x_1$

At the first step we consider the subspace $x_0 + \text{SPAN}\{p_1\}$ which consists in vectors of the form

$$x(\alpha) = x_0 + \alpha p_1 \quad \alpha \in \mathbb{R}$$

The minimization problem becomes:

Minimization step $x_0 \rightarrow x_1$

Find $x_1 = x_0 + \alpha_1 p_1$ (i.e., find α_1 !) such that:

$$\|x_* - x_1\|_A = \min_{\alpha \in \mathbb{R}} \|x_* - (x_0 + \alpha p_1)\|_A,$$



Solving first step method 2

(1/2)

Remember the error expansion:

$$x_* - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \dots + \sigma_n p_n.$$

Let $x(\alpha) = x_0 + \alpha p_1$, the difference $x_* - x(\alpha)$ becomes:

$$x_* - x(\alpha) = (\sigma_1 - \alpha) p_1 + \sigma_2 p_2 + \dots + \sigma_n p_n$$

due to conjugacy the error $\|x_* - x(\alpha)\|_A$ becomes

$$\begin{aligned} &\|x_* - x(\alpha)\|_A^2 \\ &= \left((\sigma_1 - \alpha) p_1 + \sum_{j=2}^n \sigma_j p_j \right)^T A \left((\sigma_1 - \alpha) p_1 + \sum_{j=2}^n \sigma_j p_j \right) \\ &= (\sigma_1 - \alpha)^2 p_1^T A p_1 + \sum_{j=2}^n \sigma_j^2 p_j^T A p_j \end{aligned}$$



Solving first step method 2

(2/2)

Because

$$\|x_* - x(\alpha)\|_A^2 = (\sigma_1 - \alpha)^2 \|p_1\|_A^2 + \sum_{i=2}^n \sigma_i^2 \|p_i\|_A^2,$$

we have that

$$\|x_* - x(\alpha_1)\|_A^2 = \sum_{i=2}^n \sigma_i^2 \|p_i\|_A^2 \leq \|x_* - x(\alpha)\|_A^2 \quad \text{for all } \alpha \neq \sigma_1$$

so that minimum is found by imposing $\alpha_1 = \sigma_1$:

$$\alpha_1 = \frac{p_1^T A e_0}{p_1^T A p_1}$$

This argument can be generalized for all $k > 1$ (see next slides).

Solving kth Step: $x_{k-1} \rightarrow x_k$

(1/2)

Remember the error expansion:

$$x_* - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \dots + \sigma_n p_n.$$

Consider a vector of the form

$$x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_0 + \alpha^{(1)} p_1 + \alpha^{(2)} p_2 + \dots + \alpha^{(k)} p_k$$

the error $x_* - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})$ can be written as

$$\begin{aligned} x_* - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) &= x_* - x_0 - \sum_{i=1}^k \alpha^{(i)} p_i, \\ &= \sum_{i=1}^k (\sigma_i - \alpha^{(i)}) p_i + \sum_{i=k+1}^n \sigma_i p_i. \end{aligned}$$

Step, $x_{k-1} \rightarrow x_k$

For the step from $k-1$ to k we consider the subspace of \mathbb{R}^n

$$\mathcal{V}_k = \text{SPAN}\{p_1, p_2, \dots, p_k\}$$

which contains vectors of the form:

$$x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_0 + \alpha^{(1)} p_1 + \alpha^{(2)} p_2 + \dots + \alpha^{(k)} p_k$$

The minimization problem becomes:

Minimization step $x_{k-1} \rightarrow x_k$

Find $x_k = x_0 + \alpha_1 p_1 + \alpha_2 p_2 + \dots + \alpha_k p_k$ (i.e. $\alpha_1, \alpha_2, \dots, \alpha_k$) such that:

$$\|x_* - x_k\|_A = \min_{\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)} \in \mathbb{R}} \|x_* - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})\|_A$$

Solving kth Step: $x_{k-1} \rightarrow x_k$

(2/2)

using conjugacy of p_i we obtain the norm of the error:

$$\begin{aligned} \|x_* - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})\|_A^2 \\ = \sum_{i=1}^k (\sigma_i - \alpha^{(i)})^2 \|p_i\|_A^2 + \sum_{i=k+1}^n \sigma_i^2 \|p_i\|_A^2. \end{aligned}$$

So that minimum is found by imposing $\alpha_i = \sigma_i$: for $i = 1, 2, \dots, k$.

$$\alpha_i = \frac{p_i^T A e_0}{p_i^T A p_i} \quad i = 1, 2, \dots, k$$



Successive one dimensional minimization

(1/3)

- notice that $\alpha_i = \sigma_i$ and that

$$\begin{aligned} \mathbf{x}_k &= \mathbf{x}_0 + \alpha_1 \mathbf{p}_1 + \cdots + \alpha_k \mathbf{p}_k \\ &= \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k \end{aligned}$$

- so that \mathbf{x}_{k-1} contains $k-1$ coefficients α_i for the minimization.
- if we consider the one dimensional minimization on the subspace $\mathbf{x}_{k-1} + \text{SPAN}\{\mathbf{p}_k\}$ we find again \mathbf{x}_k !



Successive one dimensional minimization

(3/3)

Using conjugacy of \mathbf{p}_i we obtain the norm of the error:

$$\|\mathbf{x}_* - \mathbf{x}(\alpha)\|_A^2 = (\sigma_k - \alpha)^2 \|\mathbf{p}_k\|_A^2 + \sum_{i=k+1}^n \sigma_i^2 \|\mathbf{p}_i\|_A^2.$$

So that minimum is found by imposing $\alpha = \sigma_k$:

$$\alpha_k = \frac{\mathbf{p}_k^T \mathbf{A} \mathbf{e}_0}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k}$$

Remark

This observation permit to perform the minimization on the k -dimensional space $\mathbf{x}_0 + \mathcal{V}_k$ as successive one dimensional minimizations along the conjugate directions \mathbf{p}_k !



Successive one dimensional minimization

(2/3)

Consider a vector of the form

$$\mathbf{x}(\alpha) = \mathbf{x}_{k-1} + \alpha \mathbf{p}_k$$

remember that $\mathbf{x}_{k-1} = \mathbf{x}_0 + \alpha_1 \mathbf{p}_1 + \cdots + \alpha_{k-1} \mathbf{p}_{k-1}$ so that the error $\mathbf{x}_* - \mathbf{x}(\alpha)$ can be written as

$$\begin{aligned} \mathbf{x}_* - \mathbf{x}(\alpha) &= \mathbf{x}_* - \mathbf{x}_0 - \sum_{i=1}^{k-1} \alpha_i \mathbf{p}_i + \alpha \mathbf{p}_k \\ &= \sum_{i=1}^{k-1} (\sigma_i - \alpha_i) \mathbf{p}_i + (\sigma_k - \alpha) \mathbf{p}_k + \sum_{i=k+1}^n \sigma_i \mathbf{p}_i. \end{aligned}$$

due to the equality $\sigma_i = \alpha_i$ the blue part of the expression is 0.



Problem (one dimensional successive minimization)

Find $\mathbf{x}_k = \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k$ such that:

$$\|\mathbf{x}_* - \mathbf{x}_k\|_A = \min_{\alpha \in \mathbb{R}} \|\mathbf{x}_* - (\mathbf{x}_{k-1} + \alpha \mathbf{p}_k)\|_A,$$

The solution is the minimum respect to α of the quadratic:

$$\begin{aligned} \Phi(\alpha) &= (\mathbf{x}_* - (\mathbf{x}_{k-1} + \alpha \mathbf{p}_k))^T \mathbf{A} (\mathbf{x}_* - (\mathbf{x}_{k-1} + \alpha \mathbf{p}_k)), \\ &= (\mathbf{e}_{k-1} - \alpha \mathbf{p}_k)^T \mathbf{A} (\mathbf{e}_{k-1} - \alpha \mathbf{p}_k), \\ &= \mathbf{e}_{k-1}^T \mathbf{A} \mathbf{e}_{k-1} - 2\alpha \mathbf{p}_k^T \mathbf{A} \mathbf{e}_{k-1} + \alpha^2 \mathbf{p}_k^T \mathbf{A} \mathbf{p}_k. \end{aligned}$$

minimum is found by imposing:

$$\frac{d\Phi(\alpha)}{d\alpha} = -2\mathbf{p}_k^T \mathbf{A} \mathbf{e}_{k-1} + 2\alpha \mathbf{p}_k^T \mathbf{A} \mathbf{p}_k = 0 \quad \Rightarrow \quad \alpha_k = \frac{\mathbf{p}_k^T \mathbf{A} \mathbf{e}_{k-1}}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k}$$



- In the case of minimization on the subspace $x_0 + \mathcal{V}_k$ we have:

$$\alpha_k = p_k^T A e_0 / p_k^T A p_k$$

- In the case of one dimensional minimization on the subspace $x_{k-1} + \text{SPAN}\{p_k\}$ we have:

$$\alpha_k = p_k^T A e_{k-1} / p_k^T A p_k$$

- Apparently they are different results, however by using the conjugacy of the vectors p_i we have

$$\begin{aligned} p_k^T A e_{k-1} &= p_k^T A (x_* - x_{k-1}) \\ &= p_k^T A (x_* - (x_0 + \alpha_1 p_1 + \dots + \alpha_{k-1} p_{k-1})) \\ &= p_k^T A e_0 - \alpha_1 p_k^T A p_1 - \dots - \alpha_{k-1} p_k^T A p_{k-1} \\ &= p_k^T A e_0 \end{aligned}$$



- The **one step minimization** in the space $x_0 + \mathcal{V}_n$ and the **successive minimization** in the space $x_{k-1} + \text{SPAN}\{p_k\}$, $k = 1, 2, \dots, n$ are equivalent if p_i s are conjugate.
- The successive minimization is useful when p_i s are not known in advance but must be computed as the minimization process proceeds.
- The evaluation of α_k is apparently not computable because e_i is not known. However noticing

$$A e_k = A (x_* - x_k) = b - A x_k = r_k$$

we can write

$$\alpha_k = p_k^T A e_{k-1} / p_k^T A p_k = p_k^T r_{k-1} / p_k^T A p_k =$$

- Finally for the residual is valid the recurrence

$$r_k = b - A x_k = b - A (x_{k-1} + \alpha_k p_k) = r_{k-1} - \alpha_k A p_k.$$



Conjugate direction minimization

Algorithm (Conjugate direction minimization)

```

k ← 0; x0 assigned;
r0 ← b - A x0;
while not converged do
    k ← k + 1;
    αk ← (rk-1T pkT) / (pkT A pk);
    xk ← xk-1 + αk pk;
    rk ← rk-1 - αk A pk;
end while

```

Observation (Computational cost)

The conjugate direction minimization requires at each step one matrix-vector product for the evaluation of α_k and two update AXPY for x_k and r_k .



Monotonic behavior of the error

Remark (Monotonic behavior of the error)

The **energy norm** of the error $\|e_k\|_A$ is monotonically decreasing in k . In fact:

$$e_k = x_* - x_k = \alpha_{k+1} p_{k+1} + \dots + \alpha_n p_n,$$

and by conjugacy

$$\|e_k\|_A^2 = \|x_* - x_k\|_A^2 = \sigma_{k+1}^2 \|p_{k+1}\|_A^2 + \dots + \sigma_n^2 \|p_n\|_A^2.$$

Finally from this relation we have $e_n = 0$.



Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- 2 Conjugate direction method
- 3 **Conjugate Gradient method**
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension

Conjugate Gradient method

The Conjugate Gradient method combine the **Conjugate Direction** method with an **orthogonalization process** (like Gram-Schmidt) applied to the residual to construct the conjugate directions. In fact, because A define a scalar product in the next slide we prove:

- each residue is orthogonal to the previous conjugate directions, and consequently linearly independent from the previous conjugate directions.
- if the residual is not null it can be used to construct a new conjugate direction.

Orthogonality of the residue r_k respect \mathcal{V}_k

- The residue r_k is orthogonal to p_1, p_2, \dots, p_k . In fact, from the error expansion

$$e_k = \alpha_{k+1}p_{k+1} + \alpha_{k+2}p_{k+2} + \dots + \alpha_n p_n$$

because $r_k = Ae_k$, for $i = 1, 2, \dots, k$ we have

$$\begin{aligned} p_i^T r_k &= p_i^T A e_k \\ &= p_i^T A \sum_{j=k+1}^n \alpha_j p_j = \sum_{j=k+1}^n \alpha_j p_i^T A p_j \\ &= 0 \end{aligned}$$

Building new conjugate direction

(1/2)

- The conjugate direction method build **one new** direction at each step.
- If $r_k \neq 0$ it can be used to build the new direction p_{k+1} by a Gram-Schmidt orthogonalization process

$$p_{k+1} = r_k + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \dots + \beta_k^{(k+1)} p_k,$$

where the k coefficients $\beta_1^{(k+1)}, \beta_2^{(k+1)}, \dots, \beta_k^{(k+1)}$ must satisfy:

$$p_i^T A p_{k+1} = 0, \quad \text{for } i = 1, 2, \dots, k.$$

Building new conjugate direction

(2/2)

(repeating from previous slide)

$$\mathbf{p}_{k+1} = \mathbf{r}_k + \beta_1^{(k+1)} \mathbf{p}_1 + \beta_2^{(k+1)} \mathbf{p}_2 + \dots + \beta_k^{(k+1)} \mathbf{p}_k,$$

expanding the expression:

$$\begin{aligned} 0 &= \mathbf{p}_i^T \mathbf{A} \mathbf{p}_{k+1}, \\ &= \mathbf{p}_i^T \mathbf{A} (\mathbf{r}_k + \beta_1^{(k+1)} \mathbf{p}_1 + \beta_2^{(k+1)} \mathbf{p}_2 + \dots + \beta_k^{(k+1)} \mathbf{p}_k), \\ &= \mathbf{p}_i^T \mathbf{A} \mathbf{r}_k + \beta_1^{(k+1)} \mathbf{p}_i^T \mathbf{A} \mathbf{p}_1, \end{aligned}$$

$$\Rightarrow \beta_i^{(k+1)} = -\frac{\mathbf{p}_i^T \mathbf{A} \mathbf{r}_k}{\mathbf{p}_i^T \mathbf{A} \mathbf{p}_i} \quad i = 1, 2, \dots, k$$



The choice of the residual $\mathbf{r}_k \neq \mathbf{0}$ for the construction of the new conjugate direction \mathbf{p}_{k+1} has **three** important consequences:

- 1 simplification of the expression for α_k ;
- 2 Orthogonality of the residual \mathbf{r}_k from the previous residue $\mathbf{r}_0, \mathbf{r}_1, \dots, \mathbf{r}_{k-1}$;
- 3 **three point formula** and simplification of the coefficients $\beta_i^{(k+1)}$.

this facts will be examined in the next slides.

Simplification of the expression for α_k

Writing the expression for \mathbf{p}_k from the orthogonalization process

$$\mathbf{p}_k = \mathbf{r}_{k-1} + \beta_1^{(k+1)} \mathbf{p}_1 + \beta_2^{(k+1)} \mathbf{p}_2 + \dots + \beta_{k-1}^{(k+1)} \mathbf{p}_{k-1},$$

using orthogonality of \mathbf{r}_{k-1} and the vectors $\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_{k-1}$, (see slide N.47) we have

$$\begin{aligned} \mathbf{r}_{k-1}^T \mathbf{p}_k &= \mathbf{r}_{k-1}^T (\mathbf{r}_{k-1} + \beta_1^{(k+1)} \mathbf{p}_1 + \beta_2^{(k+1)} \mathbf{p}_2 + \dots + \beta_{k-1}^{(k+1)} \mathbf{p}_{k-1}), \\ &= \mathbf{r}_{k-1}^T \mathbf{r}_{k-1}. \end{aligned}$$

recalling the definition of α_k it follows:

$$\alpha_k = \frac{\mathbf{e}_{k-1}^T \mathbf{A} \mathbf{p}_k}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k} = \frac{\mathbf{r}_{k-1}^T \mathbf{p}_k}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k} = \frac{\mathbf{r}_{k-1}^T \mathbf{r}_{k-1}}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k}$$

Orthogonality of the residue \mathbf{r}_k from $\mathbf{r}_0, \mathbf{r}_1, \dots, \mathbf{r}_{k-1}$

From the definition of \mathbf{p}_{i+1} it follows:

$$\begin{aligned} \mathbf{p}_{i+1} &= \mathbf{r}_i + \beta_1^{(i+1)} \mathbf{p}_1 + \beta_2^{(i+1)} \mathbf{p}_2 + \dots + \beta_i^{(i+1)} \mathbf{p}_i, \\ \Rightarrow \mathbf{r}_i &\in \text{SPAN}\{\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_i, \mathbf{p}_{i+1}\} = \mathcal{V}_{i+1} \quad (\text{obvious}) \end{aligned}$$

using orthogonality of \mathbf{r}_k and the vectors $\mathbf{p}_1, \mathbf{p}_2, \dots, \mathbf{p}_k$, (see slide N.47) for $i < k$ we have

$$\begin{aligned} \mathbf{r}_k^T \mathbf{r}_i &= \mathbf{r}_k^T \left(\mathbf{p}_{i+1} - \sum_{j=1}^i \beta_j^{(i+1)} \mathbf{p}_j \right), \\ &= \mathbf{r}_k^T \mathbf{p}_{i+1} - \sum_{j=1}^i \beta_j^{(i+1)} \mathbf{r}_k^T \mathbf{p}_j = 0. \end{aligned}$$



Three point formula and simplification of $\beta_i^{(k+1)}$

From the relation $\mathbf{r}_k^T \mathbf{r}_i = \mathbf{r}_k^T (\mathbf{r}_{i-1} - \alpha_i \mathbf{A} \mathbf{p}_i)$ we deduce

$$\mathbf{r}_k^T \mathbf{A} \mathbf{p}_i = \frac{\mathbf{r}_k^T \mathbf{r}_{i-1} - \mathbf{r}_k^T \mathbf{r}_i}{\alpha_i} = \begin{cases} -\mathbf{r}_k^T \mathbf{r}_k / \alpha_k & \text{if } i = k; \\ 0 & \text{if } i < k; \end{cases}$$

remembering that $\alpha_k = \mathbf{r}_{k-1}^T \mathbf{r}_{k-1} / \mathbf{p}_k^T \mathbf{A} \mathbf{p}_k$ we obtain

$$\beta_i^{(k+1)} = -\frac{\mathbf{r}_k^T \mathbf{A} \mathbf{p}_i}{\mathbf{p}_i^T \mathbf{A} \mathbf{p}_i} = \begin{cases} \frac{\mathbf{r}_k^T \mathbf{r}_k}{\mathbf{r}_{k-1}^T \mathbf{r}_{k-1}} & i = k; \\ 0 & i < k; \end{cases}$$

i.e. there is only one non zero coefficient $\beta_k^{(k+1)}$, so we write $\beta_k = \beta_k^{(k+1)}$ and obtain the **three point formula**:

$$\mathbf{p}_{k+1} = \mathbf{r}_k + \beta_k \mathbf{p}_k$$

Conjugate gradient algorithm

initial step:

$k \leftarrow 0$; \mathbf{x}_0 assigned;

$\mathbf{r}_0 \leftarrow \mathbf{b} - \mathbf{A} \mathbf{x}_0$;

$\mathbf{p}_1 \leftarrow \mathbf{r}_0$;

while $\|\mathbf{r}_k\| > \epsilon$ **do**

$k \leftarrow k + 1$;

Conjugate direction method

$\alpha_k \leftarrow \frac{\mathbf{r}_{k-1}^T \mathbf{r}_{k-1}}{\mathbf{p}_k^T \mathbf{A} \mathbf{p}_k}$;

$\mathbf{x}_k \leftarrow \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_k$;

$\mathbf{r}_k \leftarrow \mathbf{r}_{k-1} - \alpha_k \mathbf{A} \mathbf{p}_k$;

Residual orthogonalization

$\beta_k \leftarrow \frac{\mathbf{r}_k^T \mathbf{r}_k}{\mathbf{r}_{k-1}^T \mathbf{r}_{k-1}}$;

$\mathbf{p}_{k+1} \leftarrow \mathbf{r}_k + \beta_k \mathbf{p}_k$;

end while

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Polynomial residual expansions

(1/5)

From the Conjugate Gradient iterative scheme on slide 54 we have

Lemma

There exists k -degree polynomial $P_k(x)$ and $Q_k(x)$ such that

$$\mathbf{r}_k = P_k(\mathbf{A}) \mathbf{r}_0 \quad k = 0, 1, \dots, n$$

$$\mathbf{p}_k = Q_{k-1}(\mathbf{A}) \mathbf{r}_0 \quad k = 1, 2, \dots, n$$

Moreover $P_k(0) = 1$ for all k .

Proof.

(1/2).

The proof is by induction.

Base $k = 0$

$$\mathbf{p}_1 = \mathbf{r}_0$$

so that $P_0(x) = 1$ and $Q_0(x) = 1$.

Polynomial residual expansions

(2/5)

Proof.

(2/2).

let the expansion valid for $k-1$ Consider the recursion for the residual:

$$\begin{aligned} r_k &= r_{k-1} - \alpha_k A p_k \\ &= P_{k-1}(A) r_0 + \alpha_k A Q_{k-1}(A) r_0 \\ &= (P_{k-1}(A) + \alpha_k A Q_{k-1}(A)) r_0 \end{aligned}$$

then $P_k(x) = P_{k-1}(x) + \alpha_k x Q_{k-1}(x)$ and $P_k(0) = P_{k-1}(0) = 1$. Consider the recursion for the conjugate direction

$$\begin{aligned} p_{k+1} &= P_k(A) r_0 + \beta_k Q_{k-1}(A) r_0 \\ &= (P_k(A) + \beta_k Q_{k-1}(A)) r_0 \end{aligned}$$

then $Q_k(x) = P_k(x) + \beta_k Q_{k-1}(x)$.



Polynomial residual expansions

(3/5)

We have the following trivial equality

$$\begin{aligned} \mathcal{V}_k &= \text{SPAN}\{p_1, p_2, \dots, p_k\} \\ &= \text{SPAN}\{r_0, r_1, \dots, r_{k-1}\} \\ &= \{q(A)r_0 \mid q \in \mathbb{P}^{k-1}\} \\ &= \{p(A)e_0 \mid p \in \mathbb{P}^k, p(0) = 0\} \end{aligned}$$

In this way the optimality of CG step can be written as

$$\begin{aligned} \|x_* - x_k\|_A &\leq \|x_* - x\|_A, & \forall x \in x_0 + \mathcal{V}_k \\ \|x_* - x_k\|_A &\leq \|x_* - (x_0 + p(A)e_0)\|_A, & \forall p \in \mathbb{P}^k, p(0) = 0 \\ \|x_* - x_k\|_A &\leq \|P(A)e_0\|_A, & \forall P \in \mathbb{P}^k, P(0) = 1 \end{aligned}$$

Polynomial residual expansions

(4/5)

Recalling that

$$A^{-1}r_k = A^{-1}(b - Ax_k) = x_* - x_k = e_k$$

we can write

$$\begin{aligned} e_k &= x_* - x_k = A^{-1}r_k \\ &= A^{-1}P_k(A)r_0 \\ &= P_k(A)A^{-1}r_0 \\ &= P_k(A)(x_* - x_0) \\ &= P_k(A)e_0. \end{aligned}$$

due to the optimality of the conjugate gradient we have:



Polynomial residual expansions

(5/5)

Using the results of slide 58 and 59 we can write

$$e_k = P_k(A)e_0,$$

$$\|e_k\|_A = \|P_k(A)e_0\|_A \leq \|P(A)e_0\|_A \quad \forall P \in \mathbb{P}^k, P(0) = 1$$

and from this equation we have the estimate

$$\|e_k\|_A \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \|P(A)e_0\|_A$$

So an estimate of the form

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \|P(A)e_0\|_A \leq C_k \|e_0\|_A$$

can be used to proof a convergence rate theorem, as for the steepest descent algorithm.



Convergence rate calculation

Lemma

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix, and $p \in \mathbb{P}^k$ a polynomial, then

$$\|p(A)x\|_A \leq \|p(A)\|_2 \|x\|_A$$

Proof.

(1/2).

The matrix A is SPD so that we can write

$$A = U^T \Lambda U, \quad \Lambda = \text{DIAG}\{\lambda_1, \lambda_2, \dots, \lambda_n\}$$

where U is an orthogonal matrix (i.e. $U^T U = I$) and $\Lambda \geq 0$ is diagonal. We can define the SPD matrix $A^{1/2}$ as follows

$$A^{1/2} = U^T \Lambda^{1/2} U, \quad \Lambda^{1/2} = \text{DIAG}\{\lambda_1^{1/2}, \lambda_2^{1/2}, \dots, \lambda_n^{1/2}\}$$

and obviously $A^{1/2} A^{1/2} = A$.



Proof.

(2/2).

Notice that

$$\|x\|_A^2 = x^T A x = x^T A^{1/2} A^{1/2} x = \|A^{1/2} x\|_2^2$$

so that

$$\begin{aligned} \|p(A)x\|_A &= \|A^{1/2} p(A)x\|_2 \\ &= \|p(A) A^{1/2} x\|_2 \\ &\leq \|p(A)\|_2 \|A^{1/2} x\|_2 \\ &= \|p(A)\|_2 \|x\|_A \end{aligned}$$



Lemma

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix, and $p \in \mathbb{P}^k$ a polynomial, then

$$\|p(A)\|_2 = \max_{\lambda \in \sigma(A)} |p(\lambda)|$$

Proof.

The matrix $p(A)$ is symmetric, and for a generic symmetric matrix B we have

$$\|B\|_2 = \max_{\lambda \in \sigma(B)} |\lambda|$$

observing that if λ is an eigenvalue of A then $p(\lambda)$ is an eigenvalue of $p(A)$ the thesis easily follows. \square



- Starting the error estimate

$$\|e_k\|_A \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \|P(A)e_0\|_A$$

- Combining the last two lemma we easily obtain the estimate

$$\|e_k\|_A \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \left[\max_{\lambda \in \sigma(A)} |P(\lambda)| \right] \|e_0\|_A$$

- The convergence rate is estimated by bounding the constant

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \left[\max_{\lambda \in \sigma(A)} |P(\lambda)| \right]$$



Finite termination of Conjugate Gradient

Theorem (Finite termination of Conjugate Gradient)

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix, the the **Conjugate Gradient** applied to the linear system $Ax = b$ terminate finding the exact solution in at most n -step.

Proof.

From the estimate

$$\|e_k\|_A \leq \inf_{P \in \mathcal{P}^k, P(0)=1} \left[\max_{\lambda \in \sigma(A)} |P(\lambda)| \right] \|e_0\|_A$$

choosing
$$P(x) = \prod_{\lambda \in \sigma(A)} (x - \lambda) / \prod_{\lambda \in \sigma(A)} (0 - \lambda)$$

we have $\max_{\lambda \in \sigma(A)} |P(\lambda)| = 0$ and $\|e_n\|_A = 0$. □

Chebyshev Polynomials

(1/4)

- The **Chebyshev Polynomials of the First Kind** are the right polynomial for this estimate. This polynomial have the following definition in the interval $[-1, 1]$:

$$T_k(x) = \cos(k \arccos(x))$$

- Another equivalent definition valid in the interval $(-\infty, \infty)$ is the following

$$T_k(x) = \frac{1}{2} \left[\left(x + \sqrt{x^2 - 1} \right)^k + \left(x - \sqrt{x^2 - 1} \right)^k \right]$$

- In spite of these definition, $T_k(x)$ is effectively a polynomial.

Convergence rate of Conjugate Gradient

- The constant

$$\inf_{P \in \mathcal{P}^k, P(0)=1} \left[\max_{\lambda \in \sigma(A)} |P(\lambda)| \right]$$

is not easy to evaluate,

- The following bound, is useful

$$\max_{\lambda \in \sigma(A)} |P(\lambda)| \leq \max_{\lambda \in [\lambda_1, \lambda_n]} |P(\lambda)|$$

- in particular the final estimate will be obtained by

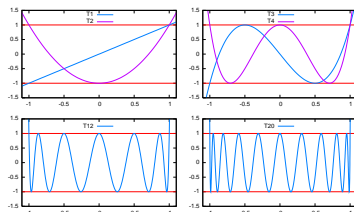
$$\inf_{P \in \mathcal{P}^k, P(0)=1} \left[\max_{\lambda \in \sigma(A)} |P(\lambda)| \right] \leq \max_{\lambda \in [\lambda_1, \lambda_n]} |\tilde{P}_k(\lambda)|$$

where $\tilde{P}_k(x)$ is an opportune k -degree polynomial for which $\tilde{P}_k(0) = 1$ and it is easy to evaluate $\max_{\lambda \in [\lambda_1, \lambda_n]} |\tilde{P}_k(\lambda)|$.

Chebyshev Polynomials

(2/4)

Some example of Chebyshev Polynomials.



Chebyshev Polynomials

(3/4)

- It is easy to show that $T_k(x)$ is a polynomial by the use of

$$\cos(\alpha + \beta) = \cos \alpha \cos \beta - \sin \alpha \sin \beta$$

$$\cos(\alpha + \beta) + \cos(\alpha - \beta) = 2 \cos \alpha \cos \beta$$

let $\theta = \arccos(x)$:

- $T_0(x) = \cos(0\theta) = 1$;
 - $T_1(x) = \cos(1\theta) = x$;
 - $T_2(x) = \cos(2\theta) = \cos(\theta)^2 - \sin(\theta)^2 = 2\cos(\theta)^2 - 1 = 2x^2 - 1$;
 - $T_{k+1}(x) + T_{k-1}(x) = \cos((k+1)\theta) + \cos((k-1)\theta)$
 $= 2 \cos(k\theta) \cos(\theta) = 2x T_k(x)$
- In general we have the following recurrence:
- $T_0(x) = 1$;
 - $T_1(x) = x$;
 - $T_{k+1}(x) = 2x T_k(x) - T_{k-1}(x)$.



Chebyshev Polynomials

(4/4)

- Solving the recurrence:
 - $T_0(x) = 1$;
 - $T_1(x) = x$;
 - $T_{k+1}(x) = 2x T_k(x) - T_{k-1}(x)$.
- We obtain the explicit form of the Chebyshev Polynomials

$$T_k(x) = \frac{1}{2} \left[\left(x + \sqrt{x^2 - 1} \right)^k + \left(x - \sqrt{x^2 - 1} \right)^k \right]$$

- The translated and scaled polynomial is useful in the study of the conjugate gradient method:

$$T_k(x; a, b) = T_k\left(\frac{a+b-2x}{b-a}\right)$$

where we have $|T_k(x; a, b)| \leq 1$ for all $x \in [a, b]$.



Convergence rate of Conjugate Gradient method

Theorem (Convergence rate of Conjugate Gradient method)

Let $A \in \mathbb{R}^{n \times n}$ an SPD matrix then the **Conjugate Gradient** method converge to the solution $x_* = A^{-1}b$ with at least linear r -rate in the norm $\|\cdot\|_A$. Moreover we have the error estimate

$$\|e_k\|_A \lesssim 2 \left(\frac{\sqrt{\kappa} - 1}{\sqrt{\kappa} + 1} \right)^k \|e_0\|_A$$

$\kappa = M/m$ is the **condition number** where $m = \lambda_1$ is the smallest eigenvalue of A and $M = \lambda_n$ is the biggest eigenvalue of A .

The expression $a_k \lesssim b_k$ means that for all $\epsilon > 0$ there exists $k_0 > 0$ such that:

$$a_k \leq (1 - \epsilon)b_k, \quad \forall k > k_0$$



Proof.

From the estimate

$$\|e_k\|_A \leq \max_{\lambda \in [m, M]} |P(\lambda)| \|e_0\|_A, \quad P \in \mathbb{P}^k, P(0) = 1$$

choosing $P(x) = T_k(x; m, M)/T_k(0; m, M)$ from the fact that $|T_k(x; m, M)| \leq 1$ for $x \in [m, M]$ we have

$$\|e_k\|_A \leq T_k(0; m, M)^{-1} \|e_0\|_A = T_k\left(\frac{M+m}{M-m}\right)^{-1} \|e_0\|_A$$

observe that $\frac{M+m}{M-m} = \frac{\kappa+1}{\kappa-1}$ and

$$T_k\left(\frac{\kappa+1}{\kappa-1}\right)^{-1} = 2 \left[\left(\frac{\sqrt{\kappa}+1}{\sqrt{\kappa}-1} \right)^k + \left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1} \right)^k \right]^{-1}$$

finally notice that $\left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1} \right)^k \rightarrow 0$ as $k \rightarrow \infty$.



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- 3 Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 **Preconditioning the Conjugate Gradient method**
- 6 Nonlinear Conjugate Gradient extension



Now we reformulate the preconditioned system:

Problem (Preconditioned linear system)

Given $A, P \in \mathbb{R}^{n \times n}$, with A an SPD matrix and P non singular matrix and $b \in \mathbb{R}^n$ the preconditioned problem is the following:

$$\text{Find } \tilde{x}_* \in \mathbb{R}^n \text{ such that: } \tilde{A} \tilde{x}_* = \tilde{b}$$

where

$$\tilde{A} = P^{-T} A P^{-1} \quad \tilde{b} = P^{-T} b$$

notice that if x_* is the solution of the linear system $Ax = b$ then $\tilde{x}_* = Px_*$ is the solution of the linear system $\tilde{A}x = \tilde{b}$.



Preconditioning

Problem (Preconditioned linear system)

Given $A, P \in \mathbb{R}^{n \times n}$, with A an SPD matrix and P non singular matrix and $b \in \mathbb{R}^n$.

$$\text{Find } x_* \in \mathbb{R}^n \text{ such that: } P^{-T} A x_* = P^{-T} b.$$

A **good** choice for P should be such that $M = P^T P \approx A$, where \approx denotes that M is an approximation of A in **some sense to precise later**.

Notice that:

- ♦ P non singular imply:

$$P^{-T}(b - Ax) = 0 \iff b - Ax = 0;$$

- ♦ A SPD imply $\tilde{A} = P^{-T} A P^{-1}$ is also SPD (obvious proof).



PCG: preliminary version

initial step:

$k \leftarrow 0$; x_0 assigned;

$\tilde{x}_0 \leftarrow Px_0$; $\tilde{r}_0 \leftarrow \tilde{b} - \tilde{A}\tilde{x}_0$; $\tilde{p}_1 \leftarrow \tilde{r}_0$;

while $\|\tilde{r}_k\| > \epsilon$ **do**

$k \leftarrow k + 1$;

Conjugate direction method

$$\tilde{\alpha}_k \leftarrow \frac{\tilde{r}_{k-1}^T \tilde{r}_{k-1}}{\tilde{p}_k^T A \tilde{p}_k};$$

$$\tilde{x}_k \leftarrow \tilde{x}_{k-1} + \tilde{\alpha}_k \tilde{p}_k;$$

$$\tilde{r}_k \leftarrow \tilde{r}_{k-1} - \tilde{\alpha}_k A \tilde{p}_k;$$

Residual orthogonalization

$$\tilde{\beta}_k \leftarrow \frac{\tilde{r}_k^T \tilde{r}_k}{\tilde{r}_{k-1}^T \tilde{r}_{k-1}};$$

$$\tilde{p}_{k+1} \leftarrow \tilde{r}_k + \tilde{\beta}_k \tilde{p}_k;$$

end while

final step

$$P^{-1} \tilde{x}_k;$$



Conjugate gradient algorithm applied to $\tilde{A}\tilde{x} = \tilde{b}$ require the evaluation of thing like:

$$\tilde{A}\tilde{p}_k = P^{-T}AP^{-1}\tilde{p}_k.$$

this can be done **without evaluate directly the matrix \tilde{A}** , by the following operations:

- 1 solve $P s'_k = \tilde{p}_k$ for $s'_k = P^{-1}\tilde{p}_k$;
- 2 evaluate $s''_k = A s'_k$;
- 3 solve $P^T s'''_k = s''_k$ for $s'''_k = P^{-T}s''_k$.

Step 1 and 3 require the solution of two auxiliary linear system. This is not a big problem if P and P^T are triangular matrices (see e.g. **incomplete Cholesky**).

Definition

For all $k \geq 1$, we introduce the vectors

$$\mathbf{x}_k = \mathbf{x}_{k-1} + \tilde{\alpha}_k \mathbf{q}_k.$$

Observation

If we assume, **by construction**, $\tilde{\mathbf{x}}_0 = P\mathbf{x}_0$, then we have

$$\tilde{\mathbf{x}}_k = P\mathbf{x}_k, \quad \text{for all } k \text{ with } 1 \leq k \leq n.$$

In fact, if $\tilde{\mathbf{x}}_{k-1} = P\mathbf{x}_{k-1}$ (inductive hypothesis), then

$$\begin{aligned} \tilde{\mathbf{x}}_k &= \tilde{\mathbf{x}}_{k-1} + \tilde{\alpha}_k \tilde{\mathbf{p}}_k && [\text{preconditioned CG}] \\ &= P\mathbf{x}_{k-1} + \tilde{\alpha}_k P\mathbf{q}_k && [\text{inductive Hyp. defs of } \mathbf{q}_k] \\ &= P(\mathbf{x}_{k-1} + \tilde{\alpha}_k \mathbf{q}_k) && [\text{obvious}] \\ &= P\mathbf{x}_k && [\text{defs. of } \mathbf{x}_k] \end{aligned}$$

However... we can reformulate the algorithm using only the matrices A and P !

Definition

For all $k \geq 1$, we introduce the vector $\mathbf{q}_k = P^{-1}\tilde{\mathbf{p}}_k$.

Observation

If the vectors $\tilde{\mathbf{p}}_1, \tilde{\mathbf{p}}_2, \dots, \tilde{\mathbf{p}}_n$ for all $1 \leq k \leq n$ are \tilde{A} -conjugate, then the corresponding vectors $\mathbf{q}_1, \mathbf{q}_2, \dots, \mathbf{q}_n$ are A -conjugate. In fact:

$$\mathbf{q}_j^T A \mathbf{q}_i = \underbrace{\tilde{\mathbf{p}}_j^T P^{-T}}_{=\mathbf{q}_j^T} \underbrace{A P^{-1}}_{=\mathbf{q}_i^T} \tilde{\mathbf{p}}_i = \tilde{\mathbf{p}}_j^T \underbrace{\tilde{A}}_{=P^{-T}AP^{-1}} \tilde{\mathbf{p}}_i = 0, \quad \text{if } i \neq j,$$

that is a consequence of \tilde{A} -conjugation of vectors $\tilde{\mathbf{p}}_i$.

Observation

Because $\tilde{\mathbf{x}}_k = P\mathbf{x}_k$ for all $k \geq 0$, we have the recurrence between the corresponding residue $\tilde{\mathbf{r}}_k = \tilde{\mathbf{b}} - \tilde{A}\tilde{\mathbf{x}}_k$ and $\mathbf{r}_k = \mathbf{b} - A\mathbf{x}_k$:

$$\tilde{\mathbf{r}}_k = P^{-T} \mathbf{r}_k.$$

In fact,

$$\begin{aligned} \tilde{\mathbf{r}}_k &= \tilde{\mathbf{b}} - \tilde{A}\tilde{\mathbf{x}}_k, && [\text{defs. of } \tilde{\mathbf{r}}_k] \\ &= P^{-T}\mathbf{b} - P^{-T}AP^{-1}P\mathbf{x}_k, && [\text{defs. of } \tilde{\mathbf{b}}, \tilde{A}, \tilde{\mathbf{x}}_k] \\ &= P^{-T}(\mathbf{b} - A\mathbf{x}_k), && [\text{obvious}] \\ &= P^{-T}\mathbf{r}_k. && [\text{defs. of } \mathbf{r}_k] \end{aligned}$$

Definition

For all k , with $1 \leq k \leq n$, the vector z_k is the solution of the linear system

$$M z_k = r_k.$$

where $M = P^T P$. Formally,

$$z_k = M^{-1} r_k = P^{-1} P^{-T} r_k.$$

Using the vectors $\{z_k\}$,

- we can express $\tilde{\alpha}_k$ and $\tilde{\beta}_k$ in terms of A , the residual r_k , and conjugate direction q_k ;
- we can build a recurrence relation for the A -conjugate directions q_k .



Observation

$$\begin{aligned} \tilde{\alpha}_k &= \frac{\tilde{r}_{k-1}^T \tilde{r}_{k-1}}{\tilde{p}_k^T A \tilde{p}_k} = \frac{r_{k-1}^T P^{-1} P^{-T} r_{k-1}}{q_k^T P^T P^{-1} A P^{-1} P q_k} = \frac{r_{k-1}^T M^{-1} r_{k-1}}{q_k^T A q_k}, \\ &= \frac{r_{k-1}^T z_{k-1}}{q_k^T A q_k}. \end{aligned}$$

Observation

$$\begin{aligned} \tilde{\beta}_k &= \frac{\tilde{r}_k^T \tilde{r}_k}{\tilde{r}_{k-1}^T \tilde{r}_{k-1}} = \frac{r_k^T P^{-1} P^{-T} r_k}{r_{k-1}^T P^{-1} P^{-T} r_{k-1}} = \frac{r_k^T M^{-1} r_k}{r_{k-1}^T M^{-1} r_{k-1}}, \\ &= \frac{r_k^T z_k}{r_{k-1}^T z_{k-1}}. \end{aligned}$$



Observation

Using the vector $z_k = M^{-1} r_k$, the following recurrence is true

$$q_{k+1} = z_k + \tilde{\beta}_k q_k$$

In fact:

$$\begin{aligned} \tilde{p}_{k+1} &= \tilde{r}_k + \tilde{\beta}_k \tilde{p}_k && [\text{preconditioned CG}] \\ P^{-1} \tilde{p}_{k+1} &= P^{-1} \tilde{r}_k + \tilde{\beta}_k P^{-1} \tilde{p}_k && [\text{left mult } P^{-1}] \\ P^{-1} \tilde{p}_{k+1} &= P^{-1} P^{-T} r_k + \tilde{\beta}_k P^{-1} \tilde{p}_k && [r_{k+1} = P^{-T} r_{k+1}] \\ P^{-1} \tilde{p}_{k+1} &= M^{-1} r_k + \tilde{\beta}_k P^{-1} \tilde{p}_k && [M^{-1} = P^{-1} P^{-T}] \\ q_{k+1} &= z_k + \tilde{\beta}_k q_k && [q_k = P^{-1} \tilde{p}_k] \end{aligned}$$



PCG: final version

initial step:

$k \leftarrow 0$; x_0 assigned;

$r_0 \leftarrow b - A x_0$; $q_1 \leftarrow r_0$;

while $\|z_k\| > \epsilon$ do

$k \leftarrow k + 1$;

Conjugate direction method

$$\tilde{\alpha}_k \leftarrow \frac{r_{k-1}^T z_{k-1}}{q_k^T A q_k};$$

$$x_k \leftarrow x_{k-1} + \tilde{\alpha}_k q_k;$$

$$r_k \leftarrow r_{k-1} - \tilde{\alpha}_k A q_k;$$

Preconditioning

$$z_k = M^{-1} r_k;$$

Residual orthogonalization

$$\tilde{\beta}_k \leftarrow \frac{r_k^T z_k}{r_{k-1}^T z_{k-1}};$$

$$q_{k+1} \leftarrow z_k + \tilde{\beta}_k q_k;$$

end while



Outline

- 1 Convergence rate of Steepest Descent iterative scheme
- 2 Conjugate direction method
- 3 Conjugate Gradient method
- 4 Conjugate Gradient convergence rate
- 5 Preconditioning the Conjugate Gradient method
- 6 Nonlinear Conjugate Gradient extension



Nonlinear Conjugate Gradient extension

- 4 The conjugate gradient algorithm can be extended for nonlinear minimization.
- 5 Fletcher and Reeves extend CG for the minimization of a general non linear function $f(\mathbf{x})$ as follows:
 - 1 Substitute the evaluation of α_k by a line search
 - 2 Substitute the residual \mathbf{r}_k with the gradient $\nabla f(\mathbf{x}_k)$
- 6 We also translate the index for the search direction \mathbf{p}_k to be more consistent with the gradients. The resulting algorithm is in the next slide



Fletcher and Reeves Nonlinear Conjugate Gradient

initial step:

$k \leftarrow 0$; \mathbf{x}_0 assigned;
 $\mathbf{f}_0 \leftarrow f(\mathbf{x}_0)$; $\mathbf{g}_0 \leftarrow \nabla f(\mathbf{x}_0)^T$;
 $\mathbf{p}_0 \leftarrow -\mathbf{g}_0$;
while $\|\mathbf{g}_k\| > \epsilon$ **do**
 $k \leftarrow k + 1$;
 Conjugate direction method
 Compute α_k by line-search;
 $\mathbf{x}_k \leftarrow \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_{k-1}$;
 $\mathbf{g}_k \leftarrow \nabla f(\mathbf{x}_k)^T$;
 Residual orthogonalization
 $\beta_k^{FR} \leftarrow \frac{\mathbf{g}_k^T \mathbf{g}_k}{\mathbf{g}_{k-1}^T \mathbf{g}_{k-1}}$;
 $\mathbf{p}_k \leftarrow -\mathbf{g}_k + \beta_k^{FR} \mathbf{p}_{k-1}$;
end while



- 5 To ensure convergence and apply Zoutendijk global convergence theorem we need to ensure that \mathbf{p}_k is a descent direction.
- 6 \mathbf{p}_0 is a descent direction by construction, for \mathbf{p}_k we have

$$\mathbf{g}_k^T \mathbf{p}_k = -\|\mathbf{g}_k\|^2 + \beta_k^{FR} \mathbf{g}_k^T \mathbf{p}_{k-1}$$

if the line-search is **exact** then $\mathbf{g}_k^T \mathbf{p}_{k-1} = 0$ because \mathbf{p}_{k-1} is the direction of the line-search. So by induction \mathbf{p}_k is a descent direction.

- 7 Exact line-search is expensive, however if we use inexact line-search with **strong Wolfe** conditions
 - 1 **sufficient decrease**: $f(\mathbf{x}_k + \alpha_k \mathbf{p}_k) \leq f(\mathbf{x}_k) + c_1 \alpha_k \nabla f(\mathbf{x}_k)^T \mathbf{p}_k$;
 - 2 **curvature condition**: $|\nabla f(\mathbf{x}_k + \alpha_k \mathbf{p}_k)^T \mathbf{p}_k| \leq c_2 |\nabla f(\mathbf{x}_k)^T \mathbf{p}_k|$.
 with $0 < c_1 < c_2 < 1/2$ then we can prove that \mathbf{p}_k is a descent direction.



The previous consideration permits to say that Fletcher and Reeves nonlinear conjugate gradient method with strong Wolfe line-search is globally convergent¹

To prove global convergence we need the following lemma:

Lemma (descent direction bound)

Suppose we apply Fletcher and Reeves nonlinear conjugate gradient method to $f(x)$ with strong Wolfe line-search with $0 < c_2 < 1/2$. The method generates descent direction p_k that satisfy the following inequality

$$-\frac{1}{1-c_2} \leq \frac{g_k^T p_k}{\|g_k\|^2} \leq -\frac{1-2c_2}{1-c_2}, \quad k = 0, 1, 2, \dots$$

¹globally here means that Zoutendijk like theorem apply

Proof. (2/3).

Using update direction formula's of the algorithm:

$$\beta_k^{FR} = \frac{g_k^T g_k}{g_{k-1}^T g_{k-1}} \quad p_k = -g_k + \beta_k^{FR} p_{k-1}$$

we can write

$$\frac{g_k^T p_k}{\|g_k\|^2} = -1 + \beta_k^{FR} \frac{g_k^T p_{k-1}}{\|g_k\|^2} = -1 + \frac{g_k^T p_{k-1}}{\|g_{k-1}\|^2}$$

and by using second strong Wolfe condition:

$$-1 + c_2 \frac{g_{k-1}^T p_{k-1}}{\|g_{k-1}\|^2} \leq \frac{g_k^T p_k}{\|g_k\|^2} \leq -1 - c_2 \frac{g_{k-1}^T p_{k-1}}{\|g_{k-1}\|^2}$$

Proof. (1/3).

The proof is by induction. First notice that the function

$$t(\xi) = \frac{2\xi - 1}{1 - \xi}$$

is monotonically increasing on the interval $[0, 1/2]$ and that $t(0) = -1$ and $t(1/2) = 0$. Hence, because of $c_2 \in (0, 1/2)$ we have:

$$-1 < \frac{2c_2 - 1}{1 - c_2} < 0. \quad (*)$$

base of induction $k = 0$: For $k = 0$ we have $p_0 = -g_0$ so that $\frac{g_0^T p_0}{\|g_0\|^2} = -1$. From $(*)$ the lemma inequality is trivially satisfied.

Proof. (3/3).

by induction we have

$$\frac{1}{1-c_2} \geq -\frac{g_{k-1}^T p_{k-1}}{\|g_{k-1}\|^2} > 0$$

so that

$$\frac{g_k^T p_k}{\|g_k\|^2} \leq -1 - c_2 \frac{g_{k-1}^T p_{k-1}}{\|g_{k-1}\|^2} \leq -1 + c_2 \frac{1}{1-c_2} = \frac{2c_2 - 1}{1 - c_2}$$

and

$$\frac{g_k^T p_k}{\|g_k\|^2} \geq -1 + c_2 \frac{g_{k-1}^T p_{k-1}}{\|g_{k-1}\|^2} \geq -1 - c_2 \frac{1}{1-c_2} = -\frac{1}{1-c_2}$$

□

- The inequality of the the previous lemma can be written as:

$$\frac{1}{1-c_2} \frac{\|g_k\|}{\|p_k\|} \geq -\frac{g_k^T p_k}{\|g_k\| \|p_k\|} \geq \frac{1-2c_2}{1-c_2} \frac{\|g_k\|}{\|p_k\|} > 0$$

- Remembering the Zoutendijk theorem we have

$$\sum_{k=1}^{\infty} (\cos \theta_k)^2 \|g_k\|^2 < \infty, \quad \text{where} \quad \cos \theta_k = -\frac{g_k^T p_k}{\|g_k\| \|p_k\|}$$

- so that if $\|g_k\| / \|p_k\|$ is bounded from below we have that $\cos \theta_k \geq \delta$ for all k and then from Zoutendijk theorem the scheme converge.
- Unfortunately this bound cant be proved so that Zoutendijk theorem cant be applied directly. However it is possible to prove a weaker results, i.e. that $\liminf_{k \rightarrow \infty} \|g_k\| = 0!$



Convergence of Fletcher and Reeves method

Assumption (Regularity assumption)

We assume $f \in C^1(\mathbb{R}^n)$ with Lipschitz continuous gradient, i.e. there exists $\gamma > 0$ such that

$$\|\nabla f(x)^T - \nabla f(y)^T\| \leq \gamma \|x - y\|, \quad \forall x, y \in \mathbb{R}^n$$



Theorem (Convergence of Fletcher and Reeves method)

Suppose the method of **Fletcher and Reeves** is implemented with strong Wolfe line-search with $0 < c_1 < c_2 < 1/2$. If $f(x)$ and x_0 satisfy the previous regularity assumptions, then

$$\liminf_{k \rightarrow \infty} \|g_k\| = 0$$

Proof.

(1/4).

From previous Lemma we have

$$\cos \theta_k \geq \frac{1}{1-c_2} \frac{\|g_k\|}{\|p_k\|} \quad k = 1, 2, \dots$$

substituting in Zoutendijk condition we have $\sum_{k=1}^{\infty} \frac{\|g_k\|^4}{\|p_k\|^2} < \infty$.

The proof is by contradiction. in fact if theorem is not true than the series diverge. Next we want to bound $\|p_k\|$.



Proof. (bounding $\|p_k\|$)

(2/4).

Using second Wolfe condition and previous Lemma

$$|g_k^T p_{k-1}| \leq -c_2 g_k^T p_{k-1} \leq \frac{c_2}{1-c_2} \|g_{k-1}\|^2$$

using $p_k \leftarrow -g_k + \beta_k^{FR} p_{k-1}$ we have

$$\begin{aligned} \|p_k\|^2 &\leq \|g_k\|^2 + 2\beta_k^{FR} |g_k^T p_{k-1}| + (\beta_k^{FR})^2 \|p_{k-1}\|^2 \\ &\leq \|g_k\|^2 + \frac{2c_2}{1-c_2} \beta_k^{FR} \|g_{k-1}\|^2 + (\beta_k^{FR})^2 \|p_{k-1}\|^2 \end{aligned}$$

recall that $\beta_k^{FR} \leftarrow \|g_k\|^2 / \|g_{k-1}\|^2$ then

$$\|p_k\|^2 \leq \frac{1+c_2}{1-c_2} \|g_k\|^2 + (\beta_k^{FR})^2 \|p_{k-1}\|^2$$



Proof. (bounding $\|p_k\|$)

(3/4).

setting $c_3 = \frac{1+c_2}{1-c_2}$ and using repeatedly the last inequality we obtain:

$$\begin{aligned}\|p_k\|^2 &\leq c_3 \|g_k\|^2 + (\beta_k^{FR})^2 (c_3 \|g_{k-1}\|^2 + (\beta_{k-1}^{FR})^2 \|p_{k-2}\|^2) \\ &= c_3 \|g_k\|^4 \left(\|g_k\|^{-2} + \|g_{k-1}\|^{-2} \right) + \frac{\|g_k\|^4}{\|g_{k-2}\|^4} \|p_{k-2}\|^2 \\ &\leq c_3 \|g_k\|^4 \left(\|g_k\|^{-2} + \|g_{k-1}\|^{-2} + \|g_{k-2}\|^{-2} \right) \\ &\quad + \frac{\|g_k\|^4}{\|g_{k-3}\|^4} \|p_{k-3}\|^2 \\ &\leq c_3 \|g_k\|^4 \sum_{j=1}^k \|g_j\|^{-2}\end{aligned}$$



Weakness of Fletcher and Reeves method

- Suppose that p_k is a **bad** search direction, i.e. $\cos \theta_k \approx 0$.
- From the **descent direction** Lemma (see slide 89) we have

$$\frac{1}{1-c_2} \frac{\|g_k\|}{\|p_k\|} \geq \cos \theta_k \geq \frac{1-2c_2}{1-c_2} \frac{\|g_k\|}{\|p_k\|} > 0$$

- so that to have $\cos \theta_k \approx 0$ we need $\|p_k\| \gg \|g_k\|$.
- since p_k is a bad direction near orthogonal to g_k it is likely that the step is small and $x_{k+1} \approx x_k$. If so we have also $g_{k+1} \approx g_k$ and $\beta_{k+1}^{FR} \approx 1$.
- but remember that $p_{k+1} \leftarrow -g_{k+1} + \beta_{k+1}^{FR} p_k$, so that $p_{k+1} \approx p_k$.
- This means that a **long sequence of unproductive iterates** will follow.



Proof.

(4/4).

Suppose now **by contradiction** there exists $\delta > 0$ such that $\|g_k\| \geq \delta$ by using the regularity assumptions we have

$$\|p_k\|^2 \leq c_3 \|g_k\|^4 \sum_{j=1}^k \|g_j\|^{-2} \leq c_3 \|g_k\|^4 \delta^{-2} k$$

Substituting in Zoutendijk condition we have

$$\infty > \sum_{k=1}^{\infty} \frac{\|g_k\|^4}{\|p_k\|^2} \geq \frac{\delta^2}{c_4} \sum_{k=1}^{\infty} \frac{1}{k} = \infty$$

this contradict assumption. □

^athe correct assumption is that there exists k_0 such that $\|g_k\| \geq \delta$ for $k \geq k_0$ but this complicate a little bit the following inequality without introducing new idea.



Polack and Ribière Nonlinear Conjugate Gradient

- The previous problem can be elided if we restart anew when the iterate stagnate.
- Restarting is obtained by simply set $\beta_k^{FR} = 0$.
- A more elegant solution can be obtained with a new definition of β_k due to Polack and Ribière is the following:

$$\beta_k^{PR} = \frac{g_k^T (g_k - g_{k-1})}{g_{k-1}^T g_{k-1}}$$

- This definition of β_k^{PR} is identical of β_k^{FR} in the case of quadratic function because $g_k^T g_{k-1} = 0$. The definition **differs** in non linear case and in particular when there is stagnation i.e. $g_k \approx g_{k-1}$ we have $\beta_k^{PR} \approx 0$, i.e. we have an **automatic restart**.



Polack and Ribière Nonlinear Conjugate Gradient

initial step:

$k \leftarrow 0$; \mathbf{x}_0 assigned;
 $f_0 \leftarrow f(\mathbf{x}_0)$; $\mathbf{g}_0 \leftarrow \nabla f(\mathbf{x}_0)^T$;
 $\mathbf{p}_0 \leftarrow -\mathbf{g}_0$;

while $\|\mathbf{g}_k\| > \epsilon$ **do**

$k \leftarrow k + 1$;

Conjugate direction method

Compute α_k by line-search;

$\mathbf{x}_k \leftarrow \mathbf{x}_{k-1} + \alpha_k \mathbf{p}_{k-1}$;

$\mathbf{g}_k \leftarrow \nabla f(\mathbf{x}_k)^T$;

Residual orthogonalization

$$\beta_k^{PR} \leftarrow \frac{\mathbf{g}_k^T (\mathbf{g}_k - \mathbf{g}_{k-1})}{\mathbf{g}_{k-1}^T \mathbf{g}_{k-1}};$$

$\mathbf{p}_k \leftarrow -\mathbf{g}_k + \beta_k^{PR} \mathbf{p}_{k-1}$;

end while



Weakness of Polack and Ribière method

(2/2)

- Polack and Ribière choice on the average perform better than Fletcher and Reeves but there is **not** convergence results!
- Although there is not convergence results there is a negative results due to Powell:

Theorem

Consider the Polack and Ribière method with exact line-search. There exists a twice continuously differentiable function $f: \mathbb{R}^3 \mapsto \mathbb{R}$ and a starting point \mathbf{x}_0 such that the sequence of gradients $\{\|\mathbf{g}_k\|\}$ is bounded away from zero.

- However in spite of this results Polack and Ribière is the first choice among conjugate direction methods.



Weakness of Polack and Ribière method

(1/2)

- Although the modification is minimal, for the Polack and Ribière method with strong Wolfe line-search it can happen that \mathbf{p}_k is not a descent direction.
- If \mathbf{p}_k is not a descent direction we can restart i.e. set $\beta_k^{PR} = 0$ or modify β_k^{PR} as follows

$$\beta_k^{PR+} = \max\{\beta_k^{PR}, 0\}$$

this new coefficient with a modified Wolfe line-search ensure that \mathbf{p}_k is a descent direction.



Other choices




- There are many other modification of the coefficient β_k that collapse to the same coefficient in the case of quadratic function. One important choice is the Hestenes and Stiefel choice

$$\beta_k^{HS} = \frac{\mathbf{g}_k^T (\mathbf{g}_k - \mathbf{g}_{k-1})}{(\mathbf{g}_k^T - \mathbf{g}_{k-1}^T) \mathbf{p}_{k-1}}$$

- For this choice there is similar convergence results of Fletcher and Reeves and similar performance.



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