Unconstrained minimization Lectures for PHD course on

Non-linear equations and numerical optimization

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Outline

- General iterative scheme
 - Descent direction failure
- Backtracking Armijo line-search
 - · Global convergence of backtracking Armijo line-search
 - · Global convergence of steepest descent
- Wolfe-Zoutendiik global convergence
 - The Wolfe conditions The Armijo-Goldstein conditions
- Algorithms for line-search
 - · Armijo Parabolic-Cubic search
 - Wolfe linesearch



The problem

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Given $f : \mathbb{R}^n \mapsto \mathbb{R}$:

 $\text{minimize} \quad f(x)$

the following regularity about f(x) is assumed in the following:

Assumption (Regularity assumption)

We assume $f \in C^1(\mathbb{R}^n)$ with Lipschitz continuous gradient, i.e. there exists $\gamma > 0$ such that

$$\|\nabla f(x)^T - \nabla f(y)^T\| \le \gamma \|x - y\|, \quad \forall x, y \in \mathbb{R}^n$$

The problem

(2/3)

Definition (Global minimum)

Given $f: \mathbb{R}^n \mapsto \mathbb{R}$ a point $x_* \in \mathbb{R}^n$ is a global minimum if

$$f(x_+) \le f(x)$$
, $\forall x \in \mathbb{R}^n$.

Definition (Local minimum)

Given $\mathsf{f}:\mathbb{R}^n\mapsto\mathbb{R}$ a point $x_\star\in\mathbb{R}^n$ is a local minimum if

$$f(x_+) \le f(x)$$
, $\forall x \in B(x_+; \delta)$.

Obviously a global minimum is a local minimum. Find a global minimum in general is not an easy task. The algorithms presented in the seguel will approximate local minima's.





Definition (Strict global minimum)

Given $f: \mathbb{R}^n \mapsto \mathbb{R}$ a point $x_* \in \mathbb{R}^n$ is a strict global minimum if

$$f(x_*) < f(x), \quad \forall x \in \mathbb{R}^n \setminus \{x_*\}.$$

Definition (Strict local minimum)

Given $f:\mathbb{R}^n\mapsto\mathbb{R}$ a point $x_\star\in\mathbb{R}^n$ is a strict local minimum if

$$f(x_{\star}) < f(x), \quad \forall x \in B(x_{\star}; \delta) \setminus \{x_{\star}\}.$$

Obviously a strict global minimum is a strict local minimum.

- The first order necessary condition do not discriminate maximum, minimum, or saddle points.
- To discriminate maximum and minimum we need more information, e.g. second order derivative of f(x).
- With second order derivative we can build necessary and sufficient condition for a minima
- In general using only first and second order derivative at the point x_{*} it is not possible to deduce a necessary and sufficient condition for a minima.

First order Necessary condition

Lemma (First order Necessary condition for local minimum)

Given $f: \mathbb{R}^n \mapsto \mathbb{R}$ satisfying the regularity assumption. If a point $x_\star \in \mathbb{R}^n$ is a local minimum then

$$\nabla f(x_{\star})^T = 0$$

Proof.

Consider a generic direction d, then for δ small enough we have

$$\lambda^{-1}(f(x_{+} + \lambda d) - f(x_{+})) \le 0, \quad 0 \le \lambda \le \delta$$

so that

$$\lim_{\lambda} \lambda^{-1} (f(x_{\star} + \lambda d) - f(x_{\star})) = \nabla f(x_{\star})d \leq 0,$$

because $oldsymbol{d}$ is a generic direction we have $abla f(x_\star)^T = oldsymbol{0}$.



Second order Necessary condition

Lemma (Second order Necessary condition for local minimum)

Given $f \in C^2(\mathbb{R}^n)$ if a point $x_* \in \mathbb{R}^n$ is a local minimum then $\nabla f(x_*)^T = \mathbf{0}$ and $\nabla^2 f(x_*)$ is semi-definite positive, i.e.

$$d^T \nabla^2 f(x_+) d \ge 0$$
, $\forall d \in \mathbb{R}^n$

Example

This condition is only, necessary, in fact consider $f(x) = x_1^2 - x_2^3$,

$$\nabla f(\mathbf{x}) = (2x_1, -3x_2^2), \quad \nabla^2 f(\mathbf{x}) = \begin{pmatrix} 2 & 0 \\ 0 & -6x_2 \end{pmatrix}$$

for the point $x_\star=0$ we have $\nabla f(\mathbf{0})=0$ and $\nabla^2 f(\mathbf{0})$ semi-definite positive, but $\mathbf{0}$ is a saddle point not a minimum.



Proof

The condition $\nabla f(x_\star)^T = \mathbf{0}$ comes from first order necessary conditions. Consider now a generic direction d, and the finite difference:

$$\frac{f(x_{\star} + \lambda d) - 2f(x_{\star}) + f(x_{\star} - \lambda d)}{\lambda^{2}} \ge 0$$

by using Taylor expansion for f(x)

$$f(\boldsymbol{x}_{\star} \pm \lambda \boldsymbol{d}) = f(\boldsymbol{x}_{\star}) \pm \nabla f(\boldsymbol{x}_{\star}) \lambda \boldsymbol{d} + \lambda^{2} \boldsymbol{d}^{T} \nabla^{2} f(\boldsymbol{x}_{\star}) \boldsymbol{d} + o(\lambda^{2})$$

and from the previous inequality

$$d^T \nabla^2 f(x_+) d + o(\lambda^2)/\lambda^2 > 0$$

taking the limit $\lambda \to 0$ and form the arbitrariness of d we have that $\nabla^2 f(x_*)$ must be semi-definite positive.

Proof

Consider now a generic direction d, and the Taylor expansion for $f(\boldsymbol{x})$

$$\begin{aligned} \mathsf{f}(\boldsymbol{x}_{\star} + \boldsymbol{d}) &= \mathsf{f}(\boldsymbol{x}_{\star}) + \nabla \mathsf{f}(\boldsymbol{x}_{\star}) \boldsymbol{d} + \boldsymbol{d}^{T} \nabla^{2} \mathsf{f}(\boldsymbol{x}_{\star}) \boldsymbol{d} + o(\|\boldsymbol{d}\|^{2}) \\ &\geq \mathsf{f}(\boldsymbol{x}_{\star}) + \lambda_{min} \|\boldsymbol{d}\|^{2} + o(\|\boldsymbol{d}\|^{2}) \end{aligned}$$

$$\geq f(x_{\star}) + \lambda_{min} \|d\| + o(\|d\|)$$

$$\geq f(x_{\star}) + \lambda_{min} \|d\|^{2} \left(1 + o(\|d\|^{2}) / \|d\|^{2}\right)$$

choosing d small enough we can write

$$f(x_{\star} + d) \ge f(x_{\star}) + \frac{\lambda_{min}}{2} \|d\|^2 > f(x_{\star}), \qquad d \ne 0, \|d\| \le \delta.$$

i.e. x_{+} is a strict minimum.

Second order sufficient condition

Lemma (Second order sufficient condition for local minimum)

Given $f \in C^2(\mathbb{R}^n)$ if a point $x_+ \in \mathbb{R}^n$ satisfy:

 $\nabla f(x_+)^T = 0;$

$$d^T \nabla^2 f(x_*) d > 0, \quad \forall d \in \mathbb{R}^n \setminus \{x_*\}$$

then
$$x_* \in \mathbb{R}^n$$
 is a strict local minimum.

Remark

Because $abla^2 f(x_\star)$ is symmetric we can write

$$\lambda_{\mathsf{min}} \boldsymbol{d}^T \boldsymbol{d} \leq \boldsymbol{d}^T \nabla^2 \mathsf{f}(\boldsymbol{x}_\star) \boldsymbol{d} \leq \lambda_{\mathsf{max}} \boldsymbol{d}^T \boldsymbol{d}$$

If $abla^2 f(x_\star)$ is positive definite we have $\lambda_{\min} > 0$.

General iterative scheme Outline

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- Backtracking Armijo line-search
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- Algorithms for line-search
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How to find a minimum

General iterative scheme

Given
$$f : \mathbb{R}^n \mapsto \mathbb{R}$$
: minimize _{$x \in \mathbb{R}^n$} $f(x)$.

We can solve the problem by solving the necessary condition. i.e by solving the nonlinear systems

$$\nabla f(x)^T = 0.$$

- Using such an approach we looses the information about f(x).
- Moreover such an approach can find solution corresponding to a maximum or saddle points.
- A better approach is to use all the information and try to build minimizing procedure, i.e. procedures that, starting from a point x_0 build a sequence $\{x_k\}$ such that $f(x_{k+1}) \le f(x_k)$. In this way, at least, we avoid to converge to a strict maximum.



- in practice very rare to be able to provide explicit minimizer.
- iterative method: given starting guess xn, generate the sequence.

$$\{x_k\}, \quad k = 1, 2, ...$$

- AIM: ensure that (a subsequence) has some favorable limiting properties:
 - satisfies first-order necessary conditions
 - satisfies second-order necessary conditions

General iterative scheme

Line-search Methods

A generic iterative minimization procedure can be sketched as follows:

- \mathbf{e} calculate a search direction \mathbf{p}_{l} from \mathbf{x}_{l}
- a ensure that this direction is a descent direction, i.e.

$$\nabla f(x_k)p_k < 0$$
, whenever $\nabla f(x_k)^T \neq 0$

so that, at least for small steps along pi, the objective function f(x) will be reduced

ullet use line-search to calculate a suitable step-length $\alpha_k>0$ so that

$$f(\mathbf{x}_k + \alpha_k \mathbf{p}_k) < f(\mathbf{x}_k)$$
.

Update the point:

$$x_{k+1} = x_k + \alpha_k p_k$$







Written with a pseudo-code the minimization procedure is the following algorithm:

Generic minimization algorithm

Given an initial guess x_0 , let k = 0;

while not converged do

Find a descent direction p_k at x_k ;

Compute a step size α_k using a line-search along p_k . Set $x_{k+1} = x_k + \alpha_k p_k$ and increase k by 1.

end while

The crucial points which differentiate the algorithms are:

- The computation of the direction p_k;
- The computation of the step size α_k.











Practical Line-search methods

• The first developed minimization algorithms try to solve

$$\alpha_k = \arg \min_{n \ge 0} f(x_k + \alpha p_k)$$

- performing exact line-search by univariate minimization;
 rather expensive and certainly not cost effective.
- rather expensive and certainly not cost effective.
 Modern methods implements inexact line-search:
 - u ensure steps are neither too long nor too short
 - u try to pick useful initial step size for fast convergence
 - best methods are based on:
 - backtracking–Armijo search;
 Armijo–Goldstein search;
 - Franke-Wolfe search:

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Existence of a descent step

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Lemma (Descent Lemma)

Suppose that f(x) satisfy the standard assumptions and that p_k is a descent direction at x_k , i.e. $\nabla f(x_k)p_k < 0$. Then we have

$$f(x_k + \alpha p_k) \le f(x_k) + \alpha \nabla f(x_k) p_k + \frac{\gamma}{2} \alpha^2 \|p_k\|^2$$

for all
$$\alpha \in [0, \alpha_k^{\star}]$$
 where $\alpha_k^{\star} = \frac{-2\nabla f(x_k)p_k}{\gamma \|p_k\|^2} > 0$

Assumption (Regularity assumption)

We assume $f \in C^1(\mathbb{R}^n)$ with Lipschitz continuous gradient, i.e. there exists $\gamma > 0$ such that

$$\|\nabla f(x) - \nabla f(y)\| \le \gamma \|x - y\|, \quad \forall x, y \in \mathbb{R}^n$$

backtracking line-search

To obtain a monotone decreasing sequence we can use the following algorithm:

Backtracking line-search

Given α_{init} (e.g., $\alpha_{init} = 1$); Given $\tau \in (0, 1)$ typically $\tau = 0.5$:

Let $\alpha^{(0)} = \alpha_{\text{init}}$; while not $f(x_k + \alpha^{(\ell)}p_k) < f(x_k)$ do set $\alpha^{(\ell+1)} = \tau \alpha^{(\ell)}$.

increase ℓ by 1;

Set $\alpha_k = \alpha^{(\ell)}$.

To be effective the previous algorithm should terminate in a finite number of steps. The next lemma assure that if p_k is a descent direction then the algorithm terminate.

Existence of a descent step

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Proof.

Let be $q(\alpha) = f(x_k + \alpha p_k)$ then we can write:

$$\begin{split} g(\alpha) - g(0) &= \int_0^\alpha g'(\xi)d\xi = \alpha g'(0) + \int_0^\alpha \left(g'(\xi) - g'(0)\right)d\xi \\ &= \alpha \nabla \mathbf{f}(\mathbf{x}_k)\mathbf{p}_k + \int_0^\alpha \left(\nabla \mathbf{f}(\mathbf{x}_k + \xi \mathbf{p}_k) - \nabla \mathbf{f}(\mathbf{x}_k)\right)\mathbf{p}_k d\xi \end{split}$$

$$\leq \alpha \nabla f(\boldsymbol{x}_k) \boldsymbol{p}_k + \int_0^{\alpha} \|\nabla f(\boldsymbol{x}_k + \xi \boldsymbol{p}_k) - \nabla f(\boldsymbol{x}_k)\| \|\boldsymbol{p}_k\| d\xi$$

$$\leq \alpha \nabla f(\boldsymbol{x}_k) \boldsymbol{p}_k + \|\boldsymbol{p}_k\|^2 \int_0^{\alpha} \gamma \xi d\xi$$

$$\leq \alpha \nabla f(x_k)p_k + \frac{\gamma \alpha^2}{2} \|p_k\|^2 = \alpha \left[\nabla f(x_k)p_k + \frac{\gamma \alpha}{2} \|p_k\|^2 \right].$$

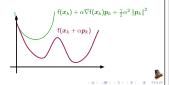
now the lemma follows trivially.



Existence of a descent step

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- The descent lemma means that there is a parabola that is entirely over the function f(x) in the direction p_i if this is a descent direction
- The second part of the lemma permits to ensure a minimal reduction if the step length is chosen to be $\alpha_k = \alpha_k^*/2$.



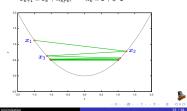
General iterative scheme

Descent direction failure

Steps may be too long

The objective function is $f(x) = x^2$ and the iterates are generated by the descent directions $p_k = (-1)^{k+1}$ from $x_0 = 2$ with:

$$x_{k+1} = x_k + \alpha_k p_k$$
, $\alpha_k = 2 + 3 \cdot 2^{-(k+1)}$



Descent direction failure

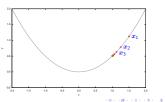
- . The simple request to have a descent direction may be not enough.
- In fact, step length may be asymptotically too short
- · Or step length may be asymptotically too long



Steps may be too short

The objective function is $f(x) = x^2$ and the iterates are generated by the descent directions $p_k = -1$ from $x_0 = 2$ with:

$$x_{k+1} = x_k + \alpha_k p_k, \qquad \alpha_k = 2^{-(k+1)}$$





Outline

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- The Armijo-Goldstein conditions
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 - Wolfe linesearch

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Backtracking Armijo line-search

Given α_{init} (e.g., $\alpha_{\mathrm{init}} = 1$); Given $\tau \in (0,1)$ typically $\tau = 0.5$; Let $\alpha^{(0)} = \alpha_{\mathrm{init}}$; while $\mathrm{not} \ f(x_k + \alpha^{(\ell)}p_k) \leq f(x_k) + \alpha^{(\ell)}\beta\nabla f(x_k)p_k$ do set $\alpha^{(\ell+1)} = \alpha^{(\ell)}$; increase ℓ by 1; end while Set $\alpha_k = \alpha^{(\ell)}$.

- Backtracking Armijo line-search prevents the step from getting too large.
- Now the question is: will the backtracking Armijo line-search terminate in a finite number of steps?

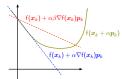
Armijo Conditio

Armijo condition

To prevent large steps relative to the decreasing of f(x) we require that

$$f(x_k + \alpha_k p_k) \le f(x_k) + \alpha_k \beta \nabla f(x_k) p_k$$

for some $\beta \in (0,1)$. Typical values of β ranges form 10^{-4} to 0.1.



Backtracking Armijo line-search

Finite termination of Armijo line-search

Theorem (Finite termination of Armijo linesearch)

Suppose that f(x) satisfy the standard assumptions and $\beta \in (0,1)$ and that p_k is a descent direction at x_k . Then the Armijo condition

$$f(x_k + \alpha_k p_k) \le f(x_k) + \alpha_k \beta \nabla f(x_k) p_k$$

is satisfied for all $\alpha_k \in [0, \omega_k]$ where $\omega_k = \frac{2(\beta-1)\nabla \mathsf{f}(\boldsymbol{x}_k)\boldsymbol{p}_k}{\gamma \, \|\boldsymbol{p}_k\|^2}$

Assumption (Regularity assumption)

We assume $f \in C^1(\mathbb{R}^n)$ with Lipschitz continuous gradient, i.e. there exists $\gamma > 0$ such that

$$\|\nabla f(x) - \nabla f(y)\| \le \gamma \|x - y\|, \quad \forall x, y \in \mathbb{R}^n$$







$$f(x + \alpha p) = f(x) + \alpha \nabla f(x)p + E$$
 where $|E| \le \frac{\gamma}{2}\alpha^2 ||p||^2$

Proof.

If $\alpha \le \omega_k$ we have $\alpha \gamma ||\mathbf{p}_k||^2 \le 2(\beta - 1)\nabla f(\mathbf{x}_k)\mathbf{p}_k$ and by using Taylor expansion

$$\begin{aligned} \mathsf{f}(\boldsymbol{x}_k + \alpha \boldsymbol{p}_k) &\leq \mathsf{f}(\boldsymbol{x}_k) + \alpha \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k + \frac{\gamma}{2} \alpha^2 \|\boldsymbol{p}_k\|^2 \\ &\leq \mathsf{f}(\boldsymbol{x}_k) + \alpha \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k + \alpha(\beta - 1) \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k \\ &\leq \mathsf{f}(\boldsymbol{x}_k) + \alpha \beta \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k \end{aligned}$$

Backtracking Armijo line-search

Backtracking-Armijo line-search

- The previous analysis permit to say that Backtracking-Armijo line-search ends in a finite number of steps.
- The line-search produce a step length not too long due to the condition

$$f(x_k + \alpha_k p_k) \le f(x_k) + \alpha_k \beta \nabla f(x_k) p_k$$

- The line-search produce a step length not too short due to the finite termination theorem
- Armijo line-search can be improved by adding some further requirements on the step length acceptance criteria.

Finite termination of Armijo line-search

Corollary (Finite termination of Armijo linesearch)

Suppose that f(x) satisfy the standard assumptions and $\beta \in (0,1)$ and that p_k is a descent direction at x_k . Then the step-size generated by then backtracking-Armijo line-search terminates with

$$\alpha_k \ge \min \{\alpha_{init}, \tau \omega_k\}, \quad \omega_k = 2(\beta - 1)\nabla f(x_k)p_k/(\gamma ||p_k||^2)$$

Proof

Line-search will terminate as soon as $\alpha^{(\ell)} \le \omega_k$:

- May be that α_{init} satisfies the Armijo condition ⇒ α_k = α_{init}.
 - Otherwise in the last line-search iteration we have

$$\alpha^{(\ell-1)} > \omega_k$$
, $\alpha_k = \alpha^{(\ell)} = \tau \alpha^{(\ell-1)} > \tau \omega_k$.

Combining these 2 cases gives the required result.



Backtracking Armijo line-search

Global convergence

Theorem (Global convergence)

Suppose that f(x) satisfy the standard assumptions, then, for the iterates generated by the Generic minimization algorithm with backtracking Armijo line-search either:

$$\circ$$
 or $\lim_{k\to\infty} f(x_k) = -\infty$;

$$\circ$$
 or $\lim_{k\to\infty} |\nabla f(x_k)p_k| \min \{1, ||p_k||^{-1}\} = 0$.

Remark

If the theorem, point 1 means that we found a stationary point in a finite number of steps. Point 2 means that function f(x) is unbounded below, so that a minimum does not exists. Point 3 alone do not imply convergence, but if $\nabla f(x_k)$ and p_k do not become orthogonal and $||p_k|| \neq 0$ then $||\nabla f(x_k)|| \rightarrow 0$.









Assume points 1 and 2 are not satisfied, then we prove point 3.

Consider

Consider
$$\mathsf{f}(x_{k+1}) \leq \mathsf{f}(x_k) + \alpha_k \beta \nabla \mathsf{f}(x_k) p_k \leq \mathsf{f}(x_0) + \sum_{k}^k \alpha_j \beta \nabla \mathsf{f}(x_j) p_j$$

by the fact that p_k is a descent direction we have that the series:

$$\sum_{i=0}^{\infty} \alpha_j \left| \nabla \mathsf{f}(\boldsymbol{x}_j) \boldsymbol{p}_j \right| \leq \beta^{-1} \lim_{k \to \infty} \left[\mathsf{f}(\boldsymbol{x}_0) - \mathsf{f}(\boldsymbol{x}_{k+1}) \right] < \infty$$

and then

$$\lim_{j\to\infty} \alpha_j |\nabla f(x_j)p_j| = 0$$

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Global convergence of backtracking Armijo line-search

Proof.

For $k \in \mathcal{K}_1$ we have $\alpha_k = \alpha_{\mathsf{init}}$ and

 $lpha_k \left|
abla {\sf f}(m{x}_k) m{p}_k
ight| = lpha_{\sf init} \left|
abla {\sf f}(m{x}_k) m{p}_k
ight|$ and from (A) we have

$$\lim_{k \in \mathcal{K}_1 o \infty} |
abla \mathsf{f}(x_k) p_k| = 0$$
 (*

For $k \in \mathcal{K}_2$ we have $\tau \omega_k \leq \alpha_k \leq \omega_k$ so

$$\alpha_k |\nabla f(x_k)p_k| \ge \tau \omega_k |\nabla f(x_k)p_k| \ge 2\tau (1 - \beta) \frac{|\nabla f(x_k)p_k|^2}{\gamma ||p_k||^2}$$

and from (B) we have

$$\lim_{k \in K_1 \to \infty} \frac{|\nabla f(x_k)p_k|}{\|p_k\|} = 0 \qquad (\star\star)$$

Combining (*) and (**) gives the required result.

wergence of backtracking Armijo line-search Backtracking Armijo line

Proof.

$$\alpha_k \geq \min\left\{\alpha_{\mathsf{init}}, \tau \omega_k\right\}, \qquad \omega_k = 2(\beta - 1) \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k / \big(\gamma \left\|\boldsymbol{p}_k\right\|^2\big)$$

and consider the two index set:

$$K_1 = \{k \mid \alpha_k = \alpha_{init}\}, \quad K_2 = \{k \mid \alpha_k < \alpha_{init}\},$$

Obviously $N=\mathcal{K}_1\cup\mathcal{K}_2$ and from $\lim_{k\to\infty}\alpha_k\,|\nabla f(x_k)p_k|=0$ we have

$$\lim_{k \in K_1 \to \infty} \alpha_k |\nabla f(x_k)p_k| = 0, \quad (A)$$

$$\lim_{k \in K_0 \to \infty} \alpha_k |\nabla f(x_k)p_k| = 0,$$
 (B)

Steepest descent algorithm

Steepest descent algorithm

Given an initial guess x_0 , let k = 0:

while not converged do

Compute a step-size α_k using a line-search along $-\nabla f(x_k)^T$. Set $x_{k+1} = x_k - \alpha_k \nabla f(x_k)^T$ and increase k by 1.

end while

- The steepest descent algorithm is simply the generic minimization algorithm with search direction the opposite of the gradient in x_L.
- The search direction $-\nabla f(x_k)^T$ is always a descent direction unless the point x_k is a stationary point.

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Corollary (Global convergence of steepest descent)

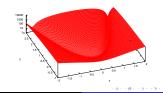
Suppose that f(x) satisfy the standard assumptions, then, for the iterates generated by the steepest descent algorithm with backtracking Armijo line-search either:

- ∇f(x_k)^T = 0 for some k ≥ 0;
- \circ or $\lim_{k\to\infty} f(x_k) = -\infty$;
- \bigcirc or $\lim_{k\to\infty} \nabla f(x_k)^T = 0$.

Backtracking Armijo line-search

- Although the steepest descent scheme is globally convergent it can be very slow!
- · A classical example is the Rosenbrock function:

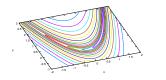
$$f(x,y) = 100(y-x^2)^2 + (x-1)^2$$



Sacktracking Armijo line-search

The Rosenbrock example

 This function has a unique minimum at (1,1)^T inside a banana shaped valley.

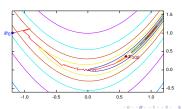


Backtracking Armijo line-search The Rosenbrock example

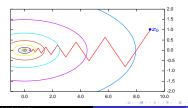
(3/3)

After 100 iteration starting from (−1.2.1)^T the approximate

minimum is far from the solution



• Given the function $f(x,y) = \frac{1}{2}x^2 + \frac{9}{2}y^2$ starting from $x_0 = (9,1)^T$ we have the zig-zag pattern toward $(0,0)^T$.



Wolfe-Zoutendijk global convergence

The Wolfe and Armijo Goldstein conditions

- The simple condition of descent step is in general not enough. for the convergence of a iterative minimization scheme.
- The condition of sufficient decrease of backtracking Armijo line-search may be insufficient on general inexact line-search algorithm.
- Adding another condition to the sufficient decrease condition such that we avoid too short step length we obtain globally convergent numerical procedure.
- O Depending on which additional condition is added we obtain the:
 - Wolfe conditions;
 - Armijo Goldstein conditions.

Outline



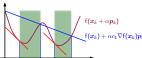
- Wolfe-Zoutendiik global convergence
- The Wolfe conditions
 - The Armijo-Goldstein conditions



The Wolfe conditions

Let c_1 and c_2 two constant such that $0 < c_1 < c_2 < 1$. We say that the step length α_{ν} satisfy the Wolfe conditions if α_{ν} satisfy:

- Sufficient decrease: f(x_k + α_kp_k) ≤ f(x_k) + c₁ α_k∇f(x_k)p_k;
- ② curvature condition: $\nabla f(x_k + \alpha_k p_k)p_k \ge c_2 \nabla f(x_k)p_k$.







The strong Wolfe conditions

Let c_1 and c_2 two constant such that $0 < c_1 < c_2 < 1$. We say that the step length α_k satisfy the strong Wolfe conditions if α_k satisfy:

- sufficient decrease: $f(x_k + \alpha_k p_k) \le f(x_k) + c_1 \alpha_k \nabla f(x_k) p_k$; • curvature condition: $|\nabla f(x_k + \alpha_k p_k) p_k| \le c_2 |\nabla f(x_k) p_k|$.
 - $\begin{array}{c} \mathsf{f}(x_k + \alpha p_k) \\ \mathsf{f}(x_k) + \alpha c_1 \nabla \mathsf{f}(x_k) p_t \end{array}$

fe conditions

Proof.

Define $\ell(\alpha)=\mathrm{f}(x_k)+\alpha c_1 \nabla \mathrm{f}(x_k)p_k$ and $g(\alpha)=\mathrm{f}(x_k+\alpha p_k)$. From $\lim_{\alpha\to\infty}\ell(\alpha)=-\infty$ and from condition 1 it follows that there exists $\alpha_{\star}>0$ such that

$$\ell(\alpha_{\star}) = g(\alpha_{\star})$$
 and $\ell(\alpha) > g(\alpha)$, $\forall \alpha \in (0, \alpha_{\star})$

so that all step length $\alpha \in (0, \alpha_\star)$ satisfy strong Wolfe condition 1. Because $\ell(0) = g(0)$ form Cauchy-Rolle theorem there exists $\alpha_{\star\star} \in (0, \alpha_\star)$ such that

$$g'(\alpha_{\star\star}) = \ell'(\alpha_{\star\star}) \Rightarrow$$

$$0 > \nabla f(x_k + \alpha_{++}p_k)p_k = c_1\nabla f(x_k)p_k > c_2\nabla f(x_k)p_k$$

by continuity we find an interval around $\alpha_{\star\star}$ with step lengths satisfying strong Wolfe conditions.

Existence of "Wolfe" step length

- The Wolfe condition seems guite restrictive.
- The next lemma answer to the question if a step length satisfying Wolfe conditions does exists.

Lemma (strong Wolfe step length)

Let $f: \mathbb{R}^n \to \mathbb{R}$ satisfying the regularity assumption. If the following condition are satisfied:

- p_k is a descent direction for the point x_k , i.e. $\nabla f(x_k)p_k < 0$;
- $f(x_k + \alpha p_k)$ is bounded from below, i.e. $\lim_{n\to\infty} f(x_k + \alpha p_k) > -\infty$.

then for any $0 < c_1 < c_2 < 1$ there exists an interval [a, b] such that all $\alpha_k \in [a, b]$ satisfy the strong Wolfe conditions.

The Wolfe condition

The Zoutendijk condition

Theorem (Zoutendijk)

Wolfe-Zoutendijk global convergence

Let $f:\mathbb{R}^n\mapsto\mathbb{R}$ satisfying the regularity assumption and bounded from below, i.e.

$$\inf_{x \in \mathbb{R}^n} f(x) > -\infty$$

Let $\{x_k\}$, $k=0,1,\ldots,\infty$ generated by a generic minimization algorithm where line-search satisfy Wolfe conditions, then

$$\sum_{k=1}^{\infty} (\cos \theta_k)^2 \left\| \nabla \mathsf{f}(\boldsymbol{x}_k)^T \right\|^2 < +\infty$$

where

$$\cos \theta_k = \frac{-\nabla f(x_k)p_k}{\|\nabla f(x_k)^T\| \|p_k\|}$$



Proof.

Using the second condition of Wolfe

$$\nabla f(x_k + \alpha_k p_k)p_k \ge c_2 \nabla f(x_k)p_k$$

 $(\nabla f(x_k + \alpha_k p_k) - \nabla f(x_k))p_k \ge (c_2 - 1)\nabla f(x_k)p_k$

 $= \alpha_{l} \gamma ||\mathbf{p}_{l}||^{2}$

by using Lipschitz regularity
$$\|\nabla f(x_k + \alpha_k p_k) - \nabla f(x_k)) p_k\| < \gamma \|x_{k+1} - x_k\| \|p_k\|$$

and using both inequality we obtain the estimate for α_k :

$$\alpha_k \ge \frac{c_2 - 1}{2 \|\mathbf{p}_k\|^2} \nabla f(\mathbf{x}_k) \mathbf{p}_k$$

Proof

Using the first condition of Wolfe and estimate of α_k

$$\begin{aligned} \mathsf{f}(\boldsymbol{x}_k + \alpha_k \boldsymbol{p}_k) &\leq \mathsf{f}(\boldsymbol{x}_k) + \alpha_k c_1 \nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k \\ &\leq \mathsf{f}(\boldsymbol{x}_k) - \frac{c_1 (1 - c_2)}{c_1 \|\boldsymbol{x}_k\|^2} \big(\nabla \mathsf{f}(\boldsymbol{x}_k) \boldsymbol{p}_k \big)^2 \end{aligned}$$

setting $A = c_1(1 - c_2)/\gamma$ and using the definition of $\cos \theta_k$

$$\mathsf{f}(\boldsymbol{x}_{k+1}) = \mathsf{f}(\boldsymbol{x}_k + \alpha_k \boldsymbol{p}_k) \leq \mathsf{f}(\boldsymbol{x}_k) - A(\cos\theta_k)^2 \left\| \nabla \mathsf{f}(\boldsymbol{x}_k)^T \right\|^2$$
 and by induction

$$f(x_{k+1})$$

$$\mathsf{f}(\boldsymbol{x}_{k+1}) \leq \mathsf{f}(\boldsymbol{x}_1) - A \sum_{j=1}^k (\cos \theta_j)^2 \left\| \nabla \mathsf{f}(\boldsymbol{x}_j)^T \right\|^2$$

Proof

The function f(x) is bounded from below, i.e

$$\inf_{x \in \mathbb{R}^n} f(x) > -\infty$$

so that

$$A\sum_{j=1}^k (\cos heta_j)^2 \left\|
abla \mathsf{f}(oldsymbol{x}_j)^T
ight\|^2 \leq \mathsf{f}(oldsymbol{x}_1) - \mathsf{f}(oldsymbol{x}_{k+1})$$

and

$$A\sum_{i=1}^{\infty}(\cos\theta_{j})^{2}\left\|\nabla \mathsf{f}(\boldsymbol{x}_{j})^{T}\right\|^{2}\leq\mathsf{f}(\boldsymbol{x}_{1})-\lim_{k\to\infty}\mathsf{f}(\boldsymbol{x}_{k+1})<+\infty$$

Corollary (Zoutendijk condition)

Let $f : \mathbb{R}^n \mapsto \mathbb{R}$ satisfying the regularity assumption and bounded from below. Let $\{x_k\}$, $k = 0, 1, ..., \infty$ generated by a generic minimization algorithm where line-search satisfy Wolfe conditions. then

$$\cos \theta_k \| \nabla f(x_k)^T \| \to 0$$
 where $\cos \theta_k = \frac{-\nabla f(x_k) p_k}{\| \nabla f(x_k)^T \| \| p_k \|}$

Remark

Wolfe-Zoutendijk global convergeno

If $\cos \theta_k \ge \delta > 0$ for all k from the Zoutendijk condition we have:

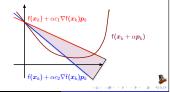
$$\|\nabla f(x_k)^T\| \rightarrow 0$$

i.e. the generic minimization algorithm where line-search satisfy Wolfe conditions converge to a stationary point.

The Armijo-Goldstein conditions

Let c_1 and c_2 two constant such that $0 < c_1 < c_2 < 1$. We say that the step length α_k satisfy the Wolfe conditions if α_k satisfy:

- $f(x_k + \alpha_k p_k) \le f(x_k) + c_1 \alpha_k \nabla f(x_k) p_k$;
- $f(x_k + \alpha_k p_k) \ge f(x_k) + c_2 \alpha_k \nabla f(x_k) p_k$;



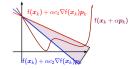
Algorithms for line-search

Outline

- General iterative scheme
 - Descent direction failure
- Racktracking Armijo line-search
 - · Global convergence of backtracking Armijo line-searc
 - Global convergence of steepest descent
- Wolfe–Zoutendijk global convergence
 - The Wolfe conditions
 - The Armijo-Goldstein conditions
- Algorithms for line-search
 - · Armijo Parabolic-Cubic search
 - Wolfe linesearch

The Armijo-Goldstein conditions

- Armijo-Goldstein conditions has very similar theoretical properties like the Wolfe conditions.
- Global convergence theorems can be established.
- The weakness of Armijo-Goldstein conditions respect to Wolfe conditions is that the former can exclude local minima's from the step length as you can see in the figure below.



Armijo Parabolic-Cubic searc

Armijo Parabolic-Cubic search

- Backtracking-Armijo line-search can be slow if a large number of reduction must be performed to satisfy Armijo condition.
- A better performance is obtained if instead of reducing by a fixed factor we use polynomial interpolation to estimate the location of the minimum.
- Assuming that that $f(x_k)$ and $\nabla f(x_k)p_k$ are known at the first step we know also $f(x_k + \lambda p_k)$ if λ is the first trial step.
- In this case a parabolic interpolation can be used to estimate the minimum.
- If we store the last trial step length, in the successive iteration we can use cubic interpolation to estimate the minima's.
- The resulting algorithm is in the following slides.



Algorithm (Armijo Parabolic-Cubic search

 $armijo_linesearch(f, x, p, \tau)$ $f_0 \leftarrow f(x)$; $\nabla f_0 \leftarrow \nabla f(x)p$; $\lambda \leftarrow 1$;

while $\lambda > \lambda_{\min}$ do

 $f_{\lambda} \leftarrow f(x + \lambda p)$;

if $f_{\lambda} \leq f_0 + \lambda \tau \nabla f_0$ then

return \(\lambda\): successful search else

if $\lambda = 1$ then $\lambda_{tmn} \leftarrow \nabla f_0 / [2(f_0 + \nabla f_0 - f_\lambda)];$

 $\lambda_{tmn} \leftarrow cubic(f_0, \nabla f_0, f_\lambda, \lambda, f_n, \lambda_n);$

 $\lambda_p \leftarrow \lambda$; $f_p \leftarrow f_{\lambda}$; $\lambda \leftarrow range(\lambda_{tmp}, \lambda/10, \lambda/2)$;

end if end while

return λ_{min} ; failed search

Armijo Parabolic-Cubic search

Algorithm (Armijo Parabolic-Cubic search

 $cubic(f_0, \nabla f_0, f_\lambda, \lambda, f_n, \lambda_n)$ Evaluate:

Algorithms for line-search

$$\begin{pmatrix} a \\ b \end{pmatrix} = \frac{1}{\lambda^2 \lambda^2 (\lambda - \lambda_p)} \begin{pmatrix} \lambda_p^2 & -\lambda^2 \\ -\lambda_n^3 & \lambda^3 \end{pmatrix} \begin{pmatrix} f_{\lambda} - f_0 - \lambda \nabla f_0 \\ f_p - f_0 - \lambda_p \nabla f_0 \end{pmatrix}$$

if a = 0 then

return $-\nabla f_0/(2b)$;

cubic is a quadratic

else $d \leftarrow b^2 - 3a \nabla f_0$ return $(-b + \sqrt{d})/(3a)$;

discriminant

end if

legitimate cubic

Algorithm (Armijo Parabolic-Cubic search

 $range(\lambda, a, b)$ if $\lambda < a$ then

return a:

else if $\lambda > b$ then return b:

else return λ ;

end if

Algorithms for line-search

Wolfe linesearch

 Wolfe linesearch is identical to the Armijo Parabolic-Cubic search, until a point satisfying the first condition is found.

At this point the Armijo algorithm stop while Wolfe search try to refine the search until the second condition is satisfied.

If the step estimated is too short then is is enlarged until it contains a minimum.

If the step estimated is too long it is reduced until the second condition is satisfied



hms for line-search W

```
Algorithm (Wolfe linesearch (1/3) wolfe_linesearch (f, x, p, c_1, c_2) f_0 \leftarrow f(x); \nabla f_0 \leftarrow \nabla f(x)p; \lambda \leftarrow 1; while \lambda \geq \lambda_{min} do f_0 \leftarrow f(x) + \lambda_{n} \nabla f_0 then g_0 to ZOOM, found a \lambda satisfying condition 1 else if \lambda = 1 then \lambda_{mp} \leftarrow \nabla f_0/[2(f_0 + \nabla f_0 - f_\lambda)]; else \lambda_{mp} \leftarrow \nabla f_0/[2(f_0 + \nabla f_0, f_\lambda, \lambda, f_p, \lambda_p)); end if \lambda_p \leftarrow \lambda; f_p \leftarrow f_\lambda; \lambda \leftarrow range(\lambda_{tmp}, \lambda/10, \lambda/2); end if end while
```

minimization

return λ_{min} ; failed search

```
Wolfe linesearch
Algorithm (Wolfe linesearch
       REFINE:
       \{\lambda_{in}, f_{in}, \nabla f_{in}\} \leftarrow \{\lambda, f_{\lambda}, \nabla f_{\lambda}\}; \Delta \leftarrow \lambda_n - \lambda_{in};
       while \Delta > \epsilon do
           \delta \lambda \leftarrow \Delta^2 \nabla f_{lo} / [2(f_{lo} + \nabla f_{lo} \Delta - f_n)]:
           \delta \lambda \leftarrow range(\delta \lambda, 0.2\Delta, 0.8\Delta):
           \lambda \leftarrow \lambda_{lo} + \delta \lambda: f_{\lambda} \leftarrow f(x + \lambda p):
           if f_{\lambda} \leq f_0 + \lambda c_1 \nabla f_0 then
               \nabla f_{\lambda} \leftarrow \nabla f(x + \lambda p)p:
                if \nabla f_{\lambda} \ge c_2 \nabla f_0 then return \lambda:
                                                                                                     found Wolfe point!
                \{\lambda_{lo}, f_{lo}, \nabla f_{lo}\} \leftarrow \{\lambda, f_{\lambda}, \nabla f_{\lambda}\}; \Delta \leftarrow \Delta - \delta \lambda;
            else
                \{\lambda_p, f_p\} \leftarrow \{\lambda, f_{\lambda}\}; \Delta \leftarrow \delta \lambda;
            end if
       end while
       return \(\lambda\): failed search
```

Algorithm (Wolfe linesearch ZOOM: $\nabla f_{\lambda} \leftarrow \nabla f(x + \lambda p)p$: if $\nabla f_{\lambda} \ge c_2 \nabla f_0$ then return λ ; found Wolfe point! if $\lambda = 1$ then forward search of an interval bracketing a minimum while $\lambda \le \lambda_{max}$ do $\{\lambda_n, f_n\} \leftarrow \{\lambda, f_\lambda\};$ save values $\lambda \leftarrow 2\lambda$; $f_{\lambda} \leftarrow f(x + \lambda p)$; if not $f_{\lambda} \leq f_0 + \lambda c_1 \nabla f_0$ then $\{\lambda_n, f_n\} \rightleftharpoons \{\lambda, f_\lambda\}$; go to REFINE; swap values end if $\nabla f_{\lambda} \leftarrow \nabla f(x + \lambda p)p$; if $\nabla f_{\lambda} \ge c_2 \nabla f_0$ then return λ ; found Wolfe point! end while return λ_{max} ; failed search end if

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