# Conjugate Direction minimization

Lectures for PHD course on Numerical optimization

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### Outline

- Introduction
- 2 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- 4 Conjugate Gradient method
- 5 Conjugate Gradient convergence rate
- 6 Preconditioning the Conjugate Gradient method
- Nonlinear Conjugate Gradient extension



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# Generic minimization algorithm

In the following we study the convergence rate of the Generic minimization algorithm applied to a quadratic function  $\mathbf{q}(x)$  with exact line search. The function

$$\mathsf{q}(\boldsymbol{x}) = \frac{1}{2}\boldsymbol{x}^T\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}^T\boldsymbol{x} + c$$

can be viewed as a n-dimensional generalization of the 1-dimensional parabolic model.

### Generic minimization algorithm

Given an initial guess  $x_0$ , let k=0;

while not converged do

Find a descent direction  $p_k$  at  $x_k$ ;

Compute a step size  $\alpha_k$  using a line-search along  $p_k$ .

Set  $x_{k+1} = x_k + \alpha_k p_k$  and increase k by 1.

end while





### Assumption (Symmetry)

The matrix A is assumed to be symmetric, in fact,

$$\boldsymbol{A} = \boldsymbol{A}^{Symm} + \boldsymbol{A}^{Skew}$$

where

$$oldsymbol{A}^{Symm} = rac{1}{2}ig[oldsymbol{A} + oldsymbol{A}^Tig], \qquad oldsymbol{A}^{Symm} = (oldsymbol{A}^{Symm})^T \ oldsymbol{A}^{Skew} = rac{1}{2}ig[oldsymbol{A} - oldsymbol{A}^Tig], \qquad oldsymbol{A}^{Skew} = -(oldsymbol{A}^{Skew})^T$$

moreover

$$\boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x} = \boldsymbol{x}^T \boldsymbol{A}^{Symm} \boldsymbol{x} + \boldsymbol{x}^T \boldsymbol{A}^{Skew} \boldsymbol{x} = \boldsymbol{x}^T \boldsymbol{A}^{Symm} \boldsymbol{x}$$

so that only the symmetric part of A contribute to q(x).





## Assumption (SPD)

The matrix  $\boldsymbol{A}$  is assumed to be symmetric and positive definite, in fact,

$$\nabla \mathsf{q}(\boldsymbol{x})^T = \frac{1}{2} \big( \boldsymbol{A} + \boldsymbol{A}^T \big) \boldsymbol{x} - \boldsymbol{b} = \boldsymbol{A} \boldsymbol{x} - \boldsymbol{b}$$

and

$$\nabla^2 \mathbf{q}(\boldsymbol{x}) = \frac{1}{2} \big( \boldsymbol{A} + \boldsymbol{A}^T \big) = \boldsymbol{A}$$

From the sufficient condition for a minimum we have that  $abla \mathsf{q}(x_\star)^T = \mathbf{0}$ , i.e.

$$Ax_{\star} = b$$

and  $abla^2 \mathsf{q}(oldsymbol{x}_\star) = oldsymbol{A}$  is SPD.





# The toy problem

 In the following we study the convergence rate of the Steepest Descent and Conjugate Gradient methods applied to

$$q(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x} - \boldsymbol{b}^T \boldsymbol{x} + c$$

where A is an SPD matrix.

• This assumption simplify the analysis but it is also useful in the non linear case. In fact, by expanding a generic function f(x) near its minimum  $x_{\star}$  we have

$$\begin{split} \mathsf{f}(\boldsymbol{x}) &= \mathsf{f}(\boldsymbol{x}_{\star}) + \nabla \mathsf{f}(\boldsymbol{x}_{\star})(\boldsymbol{x} - \boldsymbol{x}_{\star}) \\ &+ \frac{1}{2}(\boldsymbol{x} - \boldsymbol{x}_{\star})^{T} \nabla^{2} \mathsf{f}(\boldsymbol{x}_{\star})(\boldsymbol{x} - \boldsymbol{x}_{\star}) + \mathcal{O}(\|\boldsymbol{x} - \boldsymbol{x}_{\star}\|^{3}) \end{split}$$



# The toy problem

By setting

$$egin{aligned} oldsymbol{A} &= 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}), \ oldsymbol{b} &= 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} - 
abla \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} + rac{1}{2} oldsymbol{x}_{\star}^T 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} \ c &= \mathsf{f}(oldsymbol{x}_{\star}) - 
abla \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} + rac{1}{2} oldsymbol{x}_{\star}^T 
abla^2 \mathsf{f}(oldsymbol{x}_{\star}) oldsymbol{x}_{\star} \end{aligned}$$

we have

$$f(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x} - \boldsymbol{b}^T \boldsymbol{x} + c + \mathcal{O}(\|\boldsymbol{x} - \boldsymbol{x}_{\star}\|^3)$$

• So that we expect that when an iterate  $x_k$  is near  $x_\star$  then we can neglect  $\mathcal{O}(\|x-x_\star\|^3)$  and the asymptotic behavior is the same of the quadratic problem.



# The toy problem

 we can rewrite the quadratic problem in many different way as follows

$$q(\boldsymbol{x}) = \frac{1}{2}(\boldsymbol{x} - \boldsymbol{x}_{\star})^{T} \boldsymbol{A}(\boldsymbol{x} - \boldsymbol{x}_{\star}) + c'$$
$$= \frac{1}{2}(\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b})^{T} \boldsymbol{A}^{-1} (\boldsymbol{A}\boldsymbol{x} - \boldsymbol{b}) + c'$$

where

$$c' = c + \frac{1}{2} \boldsymbol{x}_{\star}^{T} \boldsymbol{A} \boldsymbol{x}_{\star}$$

 This last forms are useful in the study of the steepest descent method.





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# The steepest descent for quadratic functions

### The steepest descent minimization algorithm

Given an initial guess  $x_0$ , let k = 0;

while not converged do

Choose as descent direction  $p_k = -\nabla q(x_k)^T = b - Ax_k$ ;

Compute a step size  $\alpha_k$  using a line-search along  $p_k$ .

Set  $x_{k+1} = x_k + \alpha_k p_k$  and increase k by 1.

end while

### Definition (Residual)

The expressions

$$r(x) = b - Ax$$
,  $r_k = b - Ax_k$ 

are called the residual. We obviously have  ${m r}({m x}) = - 
abla {f q}({m x})^T$  and  ${m r}({m x}_\star) = {f 0}.$ 



(1/3)

# The steepest descent for quadratic functions

We can solve exactly the problem

$$\alpha_k = \underset{\alpha>0}{\operatorname{arg\,min}} \ \mathsf{q}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)$$

because  $p(\alpha) = \mathsf{q}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)$  is a parabola. In fact

$$\frac{\mathrm{d}p(\alpha)}{\mathrm{d}\alpha} = \frac{\mathrm{d}q(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)}{\mathrm{d}\alpha} = -\nabla q(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)\boldsymbol{r}_k = 0$$

but

$$0 = -\nabla q(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k) \boldsymbol{r}_k = \boldsymbol{r}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k)^T \boldsymbol{r}_k = (\boldsymbol{b} - \boldsymbol{A}(\boldsymbol{x}_k - \alpha \boldsymbol{r}_k))^T \boldsymbol{r}_k$$
$$= (\boldsymbol{r}_k - \alpha \boldsymbol{A} \boldsymbol{r}_k)^T \boldsymbol{r}_k$$

and the minimum is at lpha set to  $\dfrac{m{r}_k^Tm{r}_k}{m{r}_k^Tm{A}m{r}_k}.$ 



(2/3)

# The steepest descent for quadratic functions

(3/3)

### The steepest descent minimization algorithm

Given an initial guess  $x_0$ , let k=0;

while not converged do

Compute  $r_k = b - Ax_k$ ;

Compute the step size  $\alpha_k = \frac{\boldsymbol{r}_k^T \boldsymbol{r}_k}{\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k}$ ;

Set  $x_{k+1} = x_k + \alpha_k r_k$  and increase k by 1.

end while

Or more compactly

$$oldsymbol{x}_{k+1} = oldsymbol{x}_k + rac{oldsymbol{r}_k^Toldsymbol{r}_k}{oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k}oldsymbol{r}_k$$



# The steepest descent reduction step

We want bound  $q(x_{k+1})$  by  $q(x_k)$ :

$$\begin{split} \mathbf{q}(\boldsymbol{x}_{k+1}) &= \mathbf{q} \left(\boldsymbol{x}_k + \alpha_k \boldsymbol{r}_k\right) \\ &= \frac{1}{2} \left(\boldsymbol{A} \boldsymbol{x}_k + \alpha_k \boldsymbol{A} \boldsymbol{r}_k - \boldsymbol{b}\right)^T \boldsymbol{A}^{-1} \left(\boldsymbol{A} \boldsymbol{x}_k + \alpha_k \boldsymbol{A} \boldsymbol{r}_k - \boldsymbol{b}\right) + c' \\ &= \frac{1}{2} \left(\alpha_k \boldsymbol{A} \boldsymbol{r}_k - \boldsymbol{r}_k\right)^T \boldsymbol{A}^{-1} \left(\alpha_k \boldsymbol{A} \boldsymbol{r}_k - \boldsymbol{r}_k\right) + c' \\ &= \frac{1}{2} \boldsymbol{r}_k^T \boldsymbol{A}^{-1} \boldsymbol{r}_k + \frac{1}{2} \alpha_k^2 \boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k - \alpha_k \boldsymbol{r}_k^T \boldsymbol{r}_k + c' \\ &= \mathbf{q}(\boldsymbol{x}_k) + \frac{1}{2} \alpha_k \left(\alpha_k \boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k - 2 \boldsymbol{r}_k^T \boldsymbol{r}_k\right) \end{split}$$



# The steepest descent reduction step

Substituting  $\alpha_k = rac{m{r}_k^Tm{r}_k}{m{r}_k^Tm{A}m{r}_k}$  we obtain

$$\mathsf{q}(oldsymbol{x}_{k+1}) = \mathsf{q}(oldsymbol{x}_k) - rac{1}{2} rac{(oldsymbol{r}_k^T oldsymbol{r}_k)^2}{oldsymbol{r}_k^T oldsymbol{A} oldsymbol{r}_k}$$

this shows that the steepest descent method reduce at each step the objective function  $\mathbf{q}(\boldsymbol{x})$ .

Using the expression  $\mathbf{q}(\boldsymbol{x}) = \frac{1}{2} \boldsymbol{r}(\boldsymbol{x})^T \boldsymbol{A}^{-1} \boldsymbol{r}(\boldsymbol{x}) + c'$  we can write:

$$\frac{1}{2} \boldsymbol{r}_{k+1}^T \boldsymbol{A}^{-1} \boldsymbol{r}_{k+1} = \frac{1}{2} \boldsymbol{r}_k^T \boldsymbol{A}^{-1} \boldsymbol{r}_k - \frac{1}{2} \frac{(\boldsymbol{r}_k^T \boldsymbol{r}_k)^2}{\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k}$$



(2/3)



# The steepest descent reduction step

(3/3)

or better

$$m{r}_{k+1}^T m{A}^{-1} m{r}_{k+1} = m{r}_k^T m{A}^{-1} m{r}_k \left( 1 - \frac{(m{r}_k^T m{r}_k)^2}{(m{r}_k^T m{A}^{-1} m{r}_k)(m{r}_k^T m{A} m{r}_k)} \right)$$

noticing that  $m{r}_k = m{b} - m{A}m{x}_k = m{A}m{x}_\star - m{A}m{x}_k = m{A}(m{x}_\star - m{x}_k)$  we have

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k+1}\|_{\boldsymbol{A}}^2 = \|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}}^2 \left(1 - \frac{(\boldsymbol{r}_{k}^T \boldsymbol{r}_{k})^2}{(\boldsymbol{r}_{k}^T \boldsymbol{A}^{-1} \boldsymbol{r}_{k})(\boldsymbol{r}_{k}^T \boldsymbol{A} \boldsymbol{r}_{k})}\right)$$

where

$$\|\boldsymbol{x}\|_{\boldsymbol{A}} = \sqrt{\boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x}}$$

is the energy norm induced by the SPD matrix A.



The estimate of the convergence rate for the steepest descent method is linked to the estimate of the term

$$\frac{(\boldsymbol{r}_k^T\boldsymbol{r}_k)^2}{(\boldsymbol{r}_k^T\boldsymbol{A}^{-1}\boldsymbol{r}_k)(\boldsymbol{r}_k^T\boldsymbol{A}\boldsymbol{r}_k)}$$

in particular we can prove

### Lemma (Kantorovic)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the following inequality is valid

$$1 \le \frac{(\boldsymbol{x}^T \boldsymbol{A} \boldsymbol{x})(\boldsymbol{x}^T \boldsymbol{A}^{-1} \boldsymbol{x})}{(\boldsymbol{x}^T \boldsymbol{x})^2} \le \frac{(M+m)^2}{4 M m}$$

for all  $x \neq 0$ . Where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.



Proof. (1/5).

STEP 1: problem reformulation. First of all notice that

$$\frac{(\boldsymbol{x}^T\boldsymbol{A}\boldsymbol{x})(\boldsymbol{x}^T\boldsymbol{A}^{-1}\boldsymbol{x})}{(\boldsymbol{x}^T\boldsymbol{x})^2} = \frac{(\boldsymbol{y}^T\boldsymbol{A}\boldsymbol{y})(\boldsymbol{y}^T\boldsymbol{A}^{-1}\boldsymbol{y})}{(\boldsymbol{y}^T\boldsymbol{y})^2}$$

for all  $y = \alpha x$  with  $\alpha \neq 0$ . Choosing  $\alpha = ||x||^{-1}$  have:





Proof. (2/5).

STEP 2: eigenvector expansions. Matrix  $A \in \mathbb{R}^{n \times n}$  is an SPD matrix so that there exists  $u_1, u_2, \ldots, u_n$  a complete orthonormal eigenvectors set with  $0 < \lambda_1 \leq \lambda_2 \leq \cdots \leq \lambda_n$  corresponding eigenvalues. Let be  $x \in \mathbb{R}^n$  then

$$oldsymbol{x} = \sum_{k=1}^n lpha_k oldsymbol{u}_k, \qquad oldsymbol{x}^T oldsymbol{x} = \sum_{k=1}^n lpha_k^2$$

so that  $(\boldsymbol{x}^T\boldsymbol{A}\boldsymbol{x})(\boldsymbol{x}^T\boldsymbol{A}^{-1}\boldsymbol{x}) = h(\alpha_1,\ldots,\alpha_n)$  where

$$h(\alpha_1, \dots, \alpha_n) = \left(\sum_{k=1}^n \alpha_k^2 \lambda_k\right) \left(\sum_{k=1}^n \alpha_k^2 \lambda_k^{-1}\right)$$

then the lemma can be reformulated:

- Find maxima and minima of  $h(\alpha_1, \ldots, \alpha_n)$
- subject to  $\sum_{k=1}^{n} \alpha_k^2 = 1$ .



Proof. (3/5).

STEP 3: problem reduction. By using Lagrange multiplier maxima and minima are the stationary points of:

$$g(\alpha_1, \dots, \alpha_n, \mu) = h(\alpha_1, \dots, \alpha_n) + \mu \left( \sum_{k=1}^n \alpha_k^2 - 1 \right)$$

setting  $A=\sum_{k=1}^n \alpha_k^2 \lambda_k$  and  $B=\sum_{k=1}^n \alpha_k^2 \lambda_k^{-1}$  we have

$$\frac{\partial g(\alpha_1, \dots, \alpha_n, \mu)}{\partial \alpha_k} = 2\alpha_k (\lambda_k B + \lambda_k^{-1} A + \mu) = 0$$

so that

- ② Or  $\lambda_k$  is a root of the quadratic polynomial  $\lambda^2 B + \lambda \mu + A$ . in any case there are at most 2 coefficients  $\alpha$ 's not zero. <sup>a</sup>



<sup>&</sup>lt;sup>a</sup>the argument should be improved in the case of multiple eigenvalues

Proof. (4/5).

STEP 4: problem reformulation. say  $\alpha_i$  and  $\alpha_j$  are the only non zero coefficients, then  $\alpha_i^2 + \alpha_j^2 = 1$  and we can write

$$h(\alpha_1, \dots, \alpha_n) = \left(\alpha_i^2 \lambda_i + \alpha_j^2 \lambda_j\right) \left(\alpha_i^2 \lambda_i^{-1} + \alpha_j^2 \lambda_j^{-1}\right)$$

$$= \alpha_i^4 + \alpha_j^4 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i}\right)$$

$$= \alpha_i^2 (1 - \alpha_j^2) + \alpha_j^2 (1 - \alpha_i^2) + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i}\right)$$

$$= 1 + \alpha_i^2 \alpha_j^2 \left(\frac{\lambda_i}{\lambda_j} + \frac{\lambda_j}{\lambda_i} - 2\right)$$

$$= 1 + \alpha_i^2 (1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_i}$$





### Proof.

(5/5).

STEP 5: bounding maxima and minima. notice that

$$0 \le \beta(1-\beta) \le \frac{1}{4}, \quad \forall \beta \in [0,1]$$

$$1 \le 1 + \alpha_i^2 (1 - \alpha_i^2) \frac{(\lambda_i - \lambda_j)^2}{\lambda_i \lambda_j} \le 1 + \frac{(\lambda_i - \lambda_j)^2}{4\lambda_i \lambda_j} = \frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j}$$

to bound  $(\lambda_i + \lambda_j)^2/(4\lambda_i\lambda_j)$  consider the function  $f(x) = (1+x)^2/x$  which is increasing for  $x \ge 1$  so that we have

$$\frac{(\lambda_i + \lambda_j)^2}{4\lambda_i \lambda_j} \le \frac{(M+m)^2}{4 M m}$$

and finally

$$1 \le h(\alpha_1, \dots, \alpha_n) \le \frac{(M+m)^2}{4 M m}$$



# Convergence rate of Steepest Descent

The Kantorovich inequality permits to prove:

### Theorem (Convergence rate of Steepest Descent)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the steepest descent method:

$$oldsymbol{x}_{k+1} = oldsymbol{x}_k + rac{oldsymbol{r}_k^Toldsymbol{r}_k}{oldsymbol{r}_k^Toldsymbol{A}oldsymbol{r}_k}oldsymbol{r}_k$$

converge to the solution  $x_\star = A^{-1}b$  with at least linear q-rate in the norm  $\|\cdot\|_A$ . Moreover we have the error estimate

$$\left\| \boldsymbol{x}_{k+1} - \boldsymbol{x}_{\star} \right\|_{\boldsymbol{A}} \leq \frac{\kappa - 1}{\kappa + 1} \left\| \boldsymbol{x}_{k} - \boldsymbol{x}_{\star} \right\|_{\boldsymbol{A}}$$

 $\kappa = M/m$  is the condition number where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.



#### Proof.

Remember from slide  $N^{\circ}16$ 

$$\left\| oldsymbol{x}_{\star} - oldsymbol{x}_{k+1} 
ight\|_{oldsymbol{A}}^2 = \left\| oldsymbol{x}_{\star} - oldsymbol{x}_{k} 
ight\|_{oldsymbol{A}}^2 \left( 1 - rac{(oldsymbol{r}_{k}^T oldsymbol{r}_{k})^2}{(oldsymbol{r}_{k}^T oldsymbol{A}^{-1} oldsymbol{r}_{k})(oldsymbol{r}_{k}^T oldsymbol{A} oldsymbol{r}_{k})} 
ight)$$

from Kantorovich inequality

$$1 - \frac{(\boldsymbol{r}_k^T \boldsymbol{r}_k)^2}{(\boldsymbol{r}_k^T \boldsymbol{A}^{-1} \boldsymbol{r}_k)(\boldsymbol{r}_k^T \boldsymbol{A} \boldsymbol{r}_k)} \le 1 - \frac{4 \, M \, m}{(M+m)^2} = \frac{(M-m)^2}{(M+m)^2}$$

so that

$$\left\|oldsymbol{x}_{\star} - oldsymbol{x}_{k+1}
ight\|_{oldsymbol{A}} \leq rac{M-m}{M+m} \left\|oldsymbol{x}_{\star} - oldsymbol{x}_{k}
ight\|_{oldsymbol{A}}$$





### Remark (One step convergence)

The steepest descent method can converge in one iteration if  $\kappa = 1$  or when  $\mathbf{r}_0 = \mathbf{u}_k$  where  $\mathbf{u}_k$  is an eigenvector of  $\mathbf{A}$ .

- In the first case  $(\kappa = 1)$  we have  $\mathbf{A} = \beta \mathbf{I}$  for some  $\beta > 0$  so it is not interesting.
- 2 In the second case we have

$$\frac{(\boldsymbol{u}_k^T\boldsymbol{u}_k)^2}{(\boldsymbol{u}_k^T\boldsymbol{A}^{-1}\boldsymbol{u}_k)(\boldsymbol{u}_k^T\boldsymbol{A}\boldsymbol{u}_k)} = \frac{(\boldsymbol{u}_k^T\boldsymbol{u}_k)^2}{\lambda_k^{-1}(\boldsymbol{u}_k^T\boldsymbol{u}_k)\lambda_k(\boldsymbol{u}_k^T\boldsymbol{u}_k)} = 1$$

in both cases we have  $r_1 = 0$  i.e. we have found the solution.



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# Conjugate direction method

### Definition (Conjugate vector)

Given two vectors p and q in  $\mathbb{R}^n$  are conjugate respect to A if they are orthogonal respect the scalar product induced by A; i.e.,

$$\boldsymbol{p}^T \boldsymbol{A} \boldsymbol{q} = \sum_{i,j=1}^n A_{ij} p_i q_j = 0.$$

Clearly, n vectors  $p_1, p_2, \dots p_n \in \mathbb{R}^n$  that are pair wise conjugated respect to A form a base of  $\mathbb{R}^n$ .



### Problem (Linear system)

Find the minimum of  $q(x) = \frac{1}{2}x^TAx - b^Tx + c$  is equivalent to solve the first order necessary condition, i.e.

Find  $\mathbf{x}_{\star} \in \mathbb{R}^n$  such that:  $A\mathbf{x}_{\star} = \mathbf{b}$ .

#### Observation

Consider  $x_0 \in \mathbb{R}^n$  and decompose the error  $e_0 = x_\star - x_0$  by the conjugate vectors  $p_1$ ,  $p_2, \ldots, p_n \in \mathbb{R}^n$ :

$$e_0 = x_\star - x_0 = \sigma_1 p_1 + \sigma_2 p_2 + \cdots + \sigma_n p_n.$$

Evaluating the coefficients  $\sigma_1$ ,  $\sigma_2$ , ...,  $\sigma_n \in \mathbb{R}$  is equivalent to solve the problem  $Ax_* = b$ , because knowing  $e_0$  we have

$$x_{\star} = x_0 + e_0.$$





#### Observation

Using conjugacy the coefficients  $\sigma_1$ ,  $\sigma_2$ , ...,  $\sigma_n \in \mathbb{R}$  can be computed as

$$\sigma_i = rac{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{e}_0}{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i}, \qquad for \ i = 1, 2, \dots, n.$$

In fact, for all  $1 \le i \le n$ , we have

$$egin{aligned} oldsymbol{p}_i^T oldsymbol{A} oldsymbol{e}_0 &= oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_1 + \sigma_2 oldsymbol{p}_2 + \ldots + \sigma_n oldsymbol{p}_n oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_n, \ &= \sigma_1 oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i + \sigma_2 oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_2 + \ldots + \sigma_n oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_n, \ &= \sigma_i oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i, \end{aligned}$$

because  $\mathbf{p}_i^T \mathbf{A} \mathbf{p}_j = 0$  for  $i \neq j$ .





The conjugate direction method evaluate the coefficients  $\sigma_1$ ,  $\sigma_2, \ldots, \sigma_n \in \mathbb{R}$  recursively in n steps, solving for  $k \geq 0$  the minimization problem:

### Conjugate direction method

Given  $x_0$ ;  $k \leftarrow 0$ ; repeat  $k \leftarrow k+1$ ; Find  $x_k \in x_0 + \mathcal{V}_k$  such that:

$$oldsymbol{x}_k = rg \min_{oldsymbol{x} \in oldsymbol{x}_0 + \mathcal{V}_k} \|oldsymbol{x}_\star - oldsymbol{x}\|_{oldsymbol{A}}$$

until k=n

where  $\mathcal{V}_k$  is the subspace of  $\mathbb{R}^n$  generated by the first k conjugate direction; i.e.,

$$\mathcal{V}_k = \text{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \dots, \boldsymbol{p}_k\}.$$





# Step: $\boldsymbol{x}_0 o \boldsymbol{x}_1$

At the first step we consider the subspace  $x_0 + \mathrm{SPAN}\{p_1\}$  which consists in vectors of the form

$$\boldsymbol{x}(\alpha) = \boldsymbol{x}_0 + \alpha \boldsymbol{p}_1 \qquad \alpha \in \mathbb{R}$$

The minimization problem becomes:

### Minimization step $oldsymbol{x}_0 ightarrow oldsymbol{x}_1$

Find  $x_1 = x_0 + \alpha_1 p_1$  (i.e., find  $\alpha_1$ !) such that:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{1}\|_{\boldsymbol{A}} = \min_{\alpha \in \mathbb{R}} \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})\|_{\boldsymbol{A}},$$



# Solving first step method 1

The minimization problem is the minimum respect to  $\alpha$  of the quadratic:

$$\Phi(\alpha) = \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})\|_{\boldsymbol{A}}^{2},$$

$$= (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1}))^{T} \boldsymbol{A} (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + \alpha \boldsymbol{p}_{1})),$$

$$= (\boldsymbol{e}_{0} - \alpha \boldsymbol{p}_{1})^{T} \boldsymbol{A} (\boldsymbol{e}_{0} - \alpha \boldsymbol{p}_{1}),$$

$$= \boldsymbol{e}_{0}^{T} \boldsymbol{A} \boldsymbol{e}_{0} - 2\alpha \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{e}_{0} + \alpha^{2} \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{p}_{1}.$$

minimum is found by imposing:

$$\frac{\mathrm{d}\Phi(\alpha)}{\mathrm{d}\alpha} = -2\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{e}_0 + 2\alpha\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{p}_1 = 0 \quad \Rightarrow \quad \alpha_1 = \frac{\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{e}_0}{\boldsymbol{p}_1^T\boldsymbol{A}\boldsymbol{p}_1}$$





Remember the error expansion:

$$\boldsymbol{x}_{\star} - \boldsymbol{x}_0 = \sigma_1 \boldsymbol{p}_1 + \sigma_2 \boldsymbol{p}_2 + \cdots + \sigma_n \boldsymbol{p}_n.$$

Let  $x(\alpha) = x_0 + \alpha p_1$ , the difference  $x_{\star} - x(\alpha)$  becomes:

$$\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha) = (\sigma_1 - \alpha)\boldsymbol{p}_1 + \sigma_2\boldsymbol{p}_2 + \ldots + \sigma_n\boldsymbol{p}_n$$

due to conjugacy the error  $\| \boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha) \|_{\boldsymbol{A}}$  becomes

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2}$$

$$= \left( (\sigma_{1} - \alpha)\boldsymbol{p}_{1} + \sum_{i=2}^{n} \sigma_{i}\boldsymbol{p}_{i} \right)^{T} \boldsymbol{A} \left( (\sigma_{1} - \alpha)\boldsymbol{p}_{1} + \sum_{j=2}^{n} \sigma_{j}\boldsymbol{p}_{i} \right)$$

$$= (\sigma_{1} - \alpha)^{2} \boldsymbol{p}_{1}^{T} \boldsymbol{A} \boldsymbol{p}_{1} + \sum_{j=2}^{n} \sigma_{j}^{2} \boldsymbol{p}_{j}^{T} \boldsymbol{A} \boldsymbol{p}_{j}$$





# Solving first step method 2

(2/2)

Because

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2} = (\sigma_{1} - \alpha)^{2} \|\boldsymbol{p}_{1}\|_{\boldsymbol{A}}^{2} + \sum_{i=2}^{n} \sigma_{2}^{2} \|\boldsymbol{p}_{i}\|_{\boldsymbol{A}}^{2},$$

we have that

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha_1)\|_{\boldsymbol{A}}^2 = \sum_{i=2}^n \sigma_i^2 \|\boldsymbol{p}_i\|_{\boldsymbol{A}}^2 \le \|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^2 \qquad \text{for all } \alpha \neq \sigma_1$$

so that minimum is found by imposing  $\alpha_1 = \sigma_1$ :

$$\alpha_1 = \frac{\boldsymbol{p}_1^T \boldsymbol{A} \boldsymbol{e}_0}{\boldsymbol{p}_1^T \boldsymbol{A} \boldsymbol{p}_1}$$

This argument can be generalized for all k > 1 (see next slides).



# Step, $oldsymbol{x}_{k-1} ightarrow oldsymbol{x}_k$

For the step from k-1 to k we consider the subspace of  $\mathbb{R}^n$ 

$$\mathcal{V}_k = \operatorname{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \dots, \boldsymbol{p}_k\}$$

which contains vectors of the form:

$$\boldsymbol{x}(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = \boldsymbol{x}_0 + \alpha^{(1)} \boldsymbol{p}_1 + \alpha^{(2)} \boldsymbol{p}_2 + \dots + \alpha^{(k)} \boldsymbol{p}_k$$

The minimization problem becomes:

### Minimization step $oldsymbol{x}_{k-1} ightarrow oldsymbol{x}_k$

Find  $x_k = x_0 + \alpha_1 p_1 + \alpha_2 p_2 + \ldots + \alpha_k p_k$  (i.e.  $\alpha_1, \alpha_2, \ldots, \alpha_k$ ) such that:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} = \min_{\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)} \in \mathbb{R}} \|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})\|_{\boldsymbol{A}}$$





# Solving kth Step: $\boldsymbol{x}_{k-1} \rightarrow \boldsymbol{x}_k$

(1/2)

Remember the error expansion:

$$\boldsymbol{x}_{\star} - \boldsymbol{x}_0 = \sigma_1 \boldsymbol{p}_1 + \sigma_2 \boldsymbol{p}_2 + \cdots + \sigma_n \boldsymbol{p}_n.$$

Consider a vector of the form

$$x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = x_0 + \alpha^{(1)}p_1 + \alpha^{(2)}p_2 + \dots + \alpha^{(k)}p_k$$

the error  $x_{\star} - x(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)})$  can be written as

$$\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) = \boldsymbol{x}_{\star} - \boldsymbol{x}_{0} - \sum_{i=1}^{k} \alpha^{(i)} \boldsymbol{p}_{i},$$

$$= \sum_{i=1}^{k} (\sigma_i - \alpha^{(i)}) \mathbf{p}_i + \sum_{i=k+1}^{n} \sigma_i \mathbf{p}_i.$$





### Solving kth Step: $\boldsymbol{x}_{k-1} \rightarrow \boldsymbol{x}_k$

using conjugacy of  $p_i$  we obtain the norm of the error:

$$\left\| \boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha^{(1)}, \alpha^{(2)}, \dots, \alpha^{(k)}) \right\|_{\boldsymbol{A}}^{2}$$

$$= \sum_{i=1}^{k} \left( \sigma_{i} - \alpha^{(i)} \right)^{2} \left\| \boldsymbol{p}_{i} \right\|_{\boldsymbol{A}}^{2} + \sum_{i=k+1}^{n} \sigma_{i}^{2} \left\| \boldsymbol{p}_{i} \right\|_{\boldsymbol{A}}^{2}.$$

So that minimum is found by imposing  $\alpha_i = \sigma_i$ : for i = 1, 2, ..., k.

$$\alpha_i = \frac{\boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{e}_0}{\boldsymbol{p}_i^T \boldsymbol{A} \boldsymbol{p}_i} \qquad i = 1, 2, \dots, k$$



(1/3)

### Successive one dimensional minimization

• notice that  $\alpha_i = \sigma_i$  and that

$$x_k = x_0 + \alpha_1 p_1 + \dots + \alpha_k p_k$$
  
=  $x_{k-1} + \alpha_k p_k$ 

- so that  $x_{k-1}$  contains k-1 coefficients  $\alpha_i$  for the minimization.
- if we consider the one dimensional minimization on the subspace  $x_{k-1} + \operatorname{SPAN}\{p_k\}$  we find again  $x_k$ !





### Successive one dimensional minimization

Consider a vector of the form

$$\boldsymbol{x}(\alpha) = \boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_k$$

remember that  $x_{k-1} = x_0 + \alpha_1 p_1 + \cdots + \alpha_{k-1} p_{k-1}$  so that the error  $x_{\star} - x(\alpha)$  can be written as

$$\mathbf{x}_{\star} - \mathbf{x}(\alpha) = \mathbf{x}_{\star} - \mathbf{x}_{0} - \sum_{i=1}^{k-1} \alpha_{i} \mathbf{p}_{i} + \alpha \mathbf{p}_{k}$$

$$= \sum_{i=1}^{k-1} (\sigma_{i} - \alpha_{i}) \mathbf{p}_{i} + (\sigma_{k} - \alpha) \mathbf{p}_{k} + \sum_{i=k+1}^{n} \sigma_{i} \mathbf{p}_{i}.$$

due to the equality  $\sigma_i = \alpha_i$  the blue part of the expression is 0.



Using conjugacy of  $p_i$  we obtain the norm of the error:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}(\alpha)\|_{\boldsymbol{A}}^{2} = \left(\sigma_{k} - \alpha\right)^{2} \|\boldsymbol{p}_{k}\|_{\boldsymbol{A}}^{2} + \sum_{i=k+1}^{n} \sigma_{i}^{2} \|\boldsymbol{p}_{i}\|_{\boldsymbol{A}}^{2}.$$

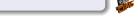
So that minimum is found by imposing  $\alpha = \sigma_k$ :

$$lpha_k = rac{oldsymbol{p}_k^T oldsymbol{A} oldsymbol{e}_0}{oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_k}$$

#### Remark

This observation permit to perform the minimization on the k-dimensional space  $x_0 + \mathcal{V}_k$  as successive one dimensional minimizations along the conjugate directions  $p_k$ !.





#### Problem (one dimensional successive minimization)

Find  $x_k = x_{k-1} + \alpha_k p_k$  such that:

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} = \min_{\alpha \in \mathbb{R}} \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_{k})\|_{\boldsymbol{A}},$$

The solution is the minimum respect to  $\alpha$  of the quadratic:

$$\Phi(\alpha) = (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_{k}))^{T} \boldsymbol{A} (\boldsymbol{x}_{\star} - (\boldsymbol{x}_{k-1} + \alpha \boldsymbol{p}_{k})),$$

$$= (\boldsymbol{e}_{k-1} - \alpha \boldsymbol{p}_{k})^{T} \boldsymbol{A} (\boldsymbol{e}_{k-1} - \alpha \boldsymbol{p}_{k}),$$

$$= \boldsymbol{e}_{k-1}^{T} \boldsymbol{A} \boldsymbol{e}_{k-1} - 2\alpha \boldsymbol{p}_{k}^{T} \boldsymbol{A} \boldsymbol{e}_{k-1} + \alpha^{2} \boldsymbol{p}_{k}^{T} \boldsymbol{A} \boldsymbol{p}_{k}.$$

minimum is found by imposing:

$$\frac{\mathrm{d}\Phi(\alpha)}{\mathrm{d}\alpha} = -2\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} + 2\alpha \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k = 0 \quad \Rightarrow \quad \boxed{\alpha_k = \frac{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1}}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k}}$$





• In the case of minimization on the subspace  $x_0 + \mathcal{V}_k$  we have:

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_0 / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k$$

• In the case of one dimensional minimization on the subspace  $x_{k-1} + \operatorname{SPAN}\{p_k\}$  we have:

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k$$

• Apparently they are different results, however by using the conjugacy of the vectors  $p_i$  we have

$$\begin{aligned} \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} &= \boldsymbol{p}_k^T \boldsymbol{A} (\boldsymbol{x}_{\star} - \boldsymbol{x}_{k-1}) \\ &= \boldsymbol{p}_k^T \boldsymbol{A} (\boldsymbol{x}_{\star} - (\boldsymbol{x}_0 + \alpha_1 \boldsymbol{p}_1 + \dots + \alpha_{k-1} \boldsymbol{p}_{k-1})) \\ &= \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_0 - \alpha_1 \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_1 - \dots - \alpha_{k-1} \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_{k-1} \\ &= \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_0 \end{aligned}$$



- The one step minimization in the space  $x_0 + \mathcal{V}_n$  and the successive minimization in the space  $x_{k-1} + \operatorname{SPAN}\{p_k\}$ ,  $k = 1, 2, \ldots, n$  are equivalent if  $p_i$ s are conjugate.
- The successive minimization is useful when  $p_i$ s are not known in advance but must be computed as the minimization process proceeds.
- The evaluation of  $\alpha_k$  is apparently not computable because  $e_i$  is not known. However noticing

$$Ae_k = A(x_\star - x_k) = b - Ax_k = r_k$$

we can write

$$\alpha_k = \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{e}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k = \boldsymbol{p}_k^T \boldsymbol{r}_{k-1} / \boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k =$$

• Finally for the residual is valid the recurrence

$$r_k = b - Ax_k = b - A(x_{k-1} + \alpha_k p_k) = r_{k-1} - \alpha_k Ap_k.$$



## Conjugate direction minimization

### Algorithm (Conjugate direction minimization)

$$\begin{aligned} k &\leftarrow 0; \ \boldsymbol{x}_0 \ \textit{assigned}; \\ \boldsymbol{r}_0 &\leftarrow \boldsymbol{b} - \boldsymbol{A}\boldsymbol{x}_0; \\ \textbf{while not converged do} \\ k &\leftarrow k+1; \\ \alpha_k &\leftarrow \frac{\boldsymbol{r}_{k-1}^T \boldsymbol{p}_k^T}{\boldsymbol{p}_k \boldsymbol{A} \boldsymbol{p}_k}; \\ \boldsymbol{x}_k &\leftarrow \boldsymbol{x}_{k-1} + \alpha_k \boldsymbol{p}_k; \\ \boldsymbol{r}_k &\leftarrow \boldsymbol{r}_{k-1} - \alpha_k \boldsymbol{A} \boldsymbol{p}_k; \\ \textbf{end while} \end{aligned}$$

### Observation (Computazional cost)

The conjugate direction minimization requires at each step one matrix–vector product for the evaluation of  $\alpha_k$  and two update AXPY for  $x_k$  and  $r_k$ .



### Monotonic behavior of the error

#### Remark (Monotonic behavior of the error)

The energy norm of the error  $\|e_k\|_A$  is monotonically decreasing in k. In fact:

$$e_k = x_{\star} - x_k = \alpha_{k+1} p_{k+1} + \ldots + \alpha_n p_n$$

and by conjugacy

$$\left\|m{e}_{k}
ight\|_{m{A}}^{2} = \left\|m{x}_{\star} - m{x}_{k}
ight\|_{m{A}}^{2} = \sigma_{k+1}^{2} \left\|m{p}_{k+1}
ight\|_{m{A}}^{2} + \ldots + \sigma_{n}^{2} \left\|m{p}_{n}
ight\|_{m{A}}^{2}.$$

Finally from this relation we have  $e_n = 0$ .



### Outline<sup>1</sup>

- Introduction
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- Conjugate direction method
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## Conjugate Gradient method

The Conjugate Gradient method combine the Conjugate Direction method with an orthogonalization process (like Gram-Schmidt) applied to the residual to construct the conjugate directions. In fact, because  $\boldsymbol{A}$  define a scalar product in the next slide we prove:

- each residue is orthogonal to the previous conjugate directions, and consequently linearly independent from the previous conjugate directions.
- if the residual is not null is can be used to construct a new conjugate direction.



## Orthogonality of the residue $oldsymbol{r}_k$ respect $\mathcal{V}_k$

• The residue  $r_k$  is orthogonal to  $p_1$ ,  $p_2$ , ...,  $p_k$ . In fact, from the error expansion

$$\boldsymbol{e}_k = \alpha_{k+1} \boldsymbol{p}_{k+1} + \alpha_{k+2} \boldsymbol{p}_{k+2} + \dots + \alpha_n \boldsymbol{p}_n$$

because  $\boldsymbol{r}_k = \boldsymbol{A}\boldsymbol{e}_k$ , for  $i=1,2,\ldots,k$  we have

$$\mathbf{p}_{i}^{T} \mathbf{r}_{k} = \mathbf{p}_{i}^{T} \mathbf{A} \mathbf{e}_{k}$$

$$= \mathbf{p}_{i}^{T} \mathbf{A} \sum_{j=k+1}^{n} \alpha_{j} \mathbf{p}_{j} = \sum_{j=k+1}^{n} \alpha_{j} \mathbf{p}_{i}^{T} \mathbf{A} \mathbf{p}_{j}$$

$$= 0$$





## Building new conjugate direction

- The conjugate direction method build one new direction at each step.
- ullet If  $oldsymbol{r}_k 
  eq oldsymbol{0}$  it can be used to build the new direction  $oldsymbol{p}_{k+1}$  by a Gram-Schmidt orthogonalization process

$$p_{k+1} = r_k + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \ldots + \beta_k^{(k+1)} p_k,$$

where the k coefficients  $\beta_1^{(k+1)}$ ,  $\beta_2^{(k+1)}$ ,  $\ldots$ ,  $\beta_k^{(k+1)}$  must satisfy:

$$p_i^T A p_{k+1} = 0,$$
 for  $i = 1, 2, ..., k$ .





### Building new conjugate direction

(repeating from previous slide)

$$p_{k+1} = r_k + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \dots + \beta_k^{(k+1)} p_k,$$

expanding the expression:

$$0 = \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{k+1},$$

$$= \boldsymbol{p}_{i}^{T} \boldsymbol{A} (\boldsymbol{r}_{k} + \beta_{1}^{(k+1)} \boldsymbol{p}_{1} + \beta_{2}^{(k+1)} \boldsymbol{p}_{2} + \dots + \beta_{k}^{(k+1)} \boldsymbol{p}_{k}),$$

$$= \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{r}_{k} + \beta_{i}^{(k+1)} \boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{i},$$

$$\Rightarrow \beta_{i}^{(k+1)} = -\frac{\boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{r}_{k}}{\boldsymbol{p}_{i}^{T} \boldsymbol{A} \boldsymbol{p}_{i}} \qquad i = 1, 2, \dots, k$$





The choice of the residual  $r_k \neq 0$  for the construction of the new conjugate direction  $p_{k+1}$  has three important consequences:

- **1** simplification of the expression for  $\alpha_k$ ;
- ② Orthogonality of the residual  $r_k$  from the previous residue  $r_0$ ,  $r_1, \ldots, r_{k-1}$ ;
- § three point formula and simplification of the coefficients  $\beta_i^{(k+1)}$ .

this facts will be examined in the next slides.



## Simplification of the expression for $\alpha_k$

Writing the expression for  $p_k$  from the orthogonalization process

$$p_k = r_{k-1} + \beta_1^{(k+1)} p_1 + \beta_2^{(k+1)} p_2 + \ldots + \beta_{k-1}^{(k+1)} p_{k-1},$$

using orthogonality of  $r_{k-1}$  and the vectors  $p_1$ ,  $p_2$ ,  $\ldots$ ,  $p_{k-1}$ , (see slide N.48) we have

$$egin{aligned} m{r}_{k-1}^T m{p}_k &= m{r}_{k-1}^T ig( m{r}_{k-1} + eta_1^{(k+1)} m{p}_1 + eta_3^{(k+1)} m{p}_2 + \ldots + eta_{k-1}^{(k+1)} m{p}_{k-1} ig), \ &= m{r}_{k-1}^T m{r}_{k-1}. \end{aligned}$$

recalling the definition of  $\alpha_k$  it follows:

$$\alpha_k = \frac{\boldsymbol{e}_{k-1}^T \boldsymbol{A} \boldsymbol{p}_k}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k} = \frac{\boldsymbol{r}_{k-1}^T \boldsymbol{p}_k}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k} = \boxed{\frac{\boldsymbol{r}_{k-1}^T \boldsymbol{r}_{k-1}}{\boldsymbol{p}_k^T \boldsymbol{A} \boldsymbol{p}_k}}$$





### Orthogonally of the residue $r_k$ from $r_0$ , $r_1$ , ..., $r_{k-1}$

From the definition of  $p_{i+1}$  it follows:

$$\begin{aligned} \boldsymbol{p}_{i+1} &= \boldsymbol{r}_i + \beta_1^{(i+1)} \boldsymbol{p}_1 + \beta_2^{(i+1)} \boldsymbol{p}_2 + \ldots + \beta_i^{(i+1)} \boldsymbol{p}_i, \\ &\Rightarrow \quad \boldsymbol{r}_i \in \text{SPAN}\{\boldsymbol{p}_1, \boldsymbol{p}_2, \ldots, \boldsymbol{p}_i, \boldsymbol{p}_{i+1}\} = \mathcal{V}_{i+1} \end{aligned} \quad \text{(obvious)}$$

using orthogonality of  $r_k$  and the vectors  $p_1$ ,  $p_2$ , ...,  $p_k$ , (see slide N.48) for i < k we have

$$egin{aligned} oldsymbol{r}_k^T oldsymbol{r}_i &= oldsymbol{r}_k^T oldsymbol{p}_{i+1} - \sum_{j=1}^i eta_j^{(i+1)} oldsymbol{p}_j igg), \ &= oldsymbol{r}_k^T oldsymbol{p}_{i+1} - \sum_{j=1}^i eta_j^{(i+1)} oldsymbol{r}_k^T oldsymbol{p}_j = 0. \end{aligned}$$





# Three point formula and simplification of $eta_i^{(k+1)}$

From the relation  $m{r}_k^Tm{r}_i = m{r}_k^T(m{r}_{i-1} - lpha_im{A}m{p}_i)$  we deduce

$$m{r}_k^T m{A} m{p}_i = rac{m{r}_k^T m{r}_{i-1} - m{r}_k^T m{r}_i}{lpha_i} = egin{cases} -m{r}_k^T m{r}_k / lpha_k & ext{if } i = k; \\ 0 & ext{if } i < k; \end{cases}$$

remembering that  $lpha_k = m{r}_{k-1}^T m{r}_{k-1} \ / \ m{p}_k^T m{A} m{p}_k$  we obtain

$$eta_i^{(k+1)} = -rac{oldsymbol{r}_k^T oldsymbol{A} oldsymbol{p}_i}{oldsymbol{p}_i^T oldsymbol{A} oldsymbol{p}_i} = \left\{ egin{array}{c} rac{oldsymbol{r}_k^T oldsymbol{r}_k}{oldsymbol{r}_{k-1}^T oldsymbol{r}_{k-1}} & i = k; \ 0 & i < k; \end{array} 
ight.$$

i.e. there is only one non zero coefficient  $\beta_k^{(k+1)}$ , so we write  $\beta_k = \beta_k^{(k+1)}$  and obtain the three point formula:

$$\boldsymbol{p}_{k+1} = \boldsymbol{r}_k + \beta_k \boldsymbol{p}_k$$





## Conjugate gradient algorithm

#### initial step:

$$k \leftarrow 0; \ oldsymbol{x}_0 ext{ assigned}; \ oldsymbol{r}_0 \leftarrow oldsymbol{b} - oldsymbol{A} oldsymbol{x}_0; \ oldsymbol{p}_1 \leftarrow oldsymbol{r}_0; \ ext{while} \ \|oldsymbol{r}_k\| > \epsilon \ ext{do} \ k \leftarrow k+1;$$

#### Conjugate direction method

$$lpha_k \leftarrow rac{oldsymbol{r}_{k-1}^T oldsymbol{r}_{k-1}}{oldsymbol{p}_k^T oldsymbol{A} oldsymbol{p}_k}; \ oldsymbol{x}_k \leftarrow oldsymbol{x}_{k-1} + lpha_k oldsymbol{p}_k; \ oldsymbol{r}_k \leftarrow oldsymbol{r}_{k-1} - lpha_k oldsymbol{A} oldsymbol{p}_k;$$

#### Residual orthogonalization

$$eta_k \leftarrow rac{oldsymbol{r}_k^T oldsymbol{r}_k}{oldsymbol{r}_{k-1}^T oldsymbol{r}_{k-1}}; \ oldsymbol{p}_{k+1} \leftarrow oldsymbol{r}_k + eta_k oldsymbol{p}_k;$$

end while



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## Polynomial residual expansions

(1/5)

From the Conjugate Gradient iterative scheme on slide 55 we have

### Lemma

There exists k-degree polynomial  $P_k(x)$  and  $Q_k(x)$  such that

$$\boldsymbol{r}_k = P_k(\boldsymbol{A})\boldsymbol{r}_0$$

$$k=0,1,\ldots,n$$

$$\boldsymbol{p}_k = Q_{k-1}(\boldsymbol{A})\boldsymbol{r}_0 \qquad k = 1, 2, \dots, n$$

Moreover  $P_k(0) = 1$  for all k.

### Proof.

(1/2).

The proof is by induction.

Base k=0

$$p_1 = r_0$$

so that  $P_0(x) = 1$  and  $Q_0(x) = 1$ .



### Polynomial residual expansions

(2/5)

Proof. (2/2).

let the expansion valid for k-1 Consider the recursion for the residual:

$$r_k = r_{k-1} - \alpha_k \mathbf{A} \mathbf{p}_k$$

$$= P_{k-1}(\mathbf{A}) \mathbf{r}_0 + \alpha_k \mathbf{A} Q_{k-1}(\mathbf{A}) \mathbf{r}_0$$

$$= (P_{k-1}(\mathbf{A}) + \alpha_k \mathbf{A} Q_{k-1}(\mathbf{A})) \mathbf{r}_0$$

then  $P_k(x) = P_{k-1}(x) + \alpha_k x Q_{k-1}(x)$  and  $P_k(0) = P_{k-1}(0) = 1$ . Consider the recursion for the conjugate direction

$$\mathbf{p}_{k+1} = P_k(\mathbf{A})\mathbf{r}_0 + \beta_k Q_{k-1}(\mathbf{A})\mathbf{r}_0$$
$$= (P_k(\mathbf{A}) + \beta_k Q_{k-1}(\mathbf{A}))\mathbf{r}_0$$

then  $Q_k(x) = P_k(x) + \beta_k Q_{k-1}(x)$ .



We have the following trivial equality

$$\mathcal{V}_k = \operatorname{SPAN} \{ \boldsymbol{p}_1, \boldsymbol{p}_2, \dots \boldsymbol{p}_k \}$$

$$= \operatorname{SPAN} \{ \boldsymbol{r}_0, \boldsymbol{r}_1, \dots \boldsymbol{r}_{k-1} \}$$

$$= \{ q(\boldsymbol{A}) \boldsymbol{r}_0 \mid q \in \mathbb{P}^{k-1}, \}$$

$$= \{ p(\boldsymbol{A}) \boldsymbol{e}_0 \mid p \in \mathbb{P}^k, p(0) = 0 \}$$

In this way the optimality of CG step can be written as

$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} \leq \|\boldsymbol{x}_{\star} - \boldsymbol{x}\|_{\boldsymbol{A}}, \qquad \forall \boldsymbol{x} \in \boldsymbol{x}_{0} + \mathcal{V}_{k}$$
$$\|\boldsymbol{x}_{\star} - \boldsymbol{x}_{k}\|_{\boldsymbol{A}} \leq \|\boldsymbol{x}_{\star} - (\boldsymbol{x}_{0} + p(\boldsymbol{A})\boldsymbol{e}_{0})\|_{\boldsymbol{A}}, \qquad \forall p \in \mathbb{P}^{k}, \ p(0) = 0$$

 $\forall P \in \mathbb{P}^k, P(0) = 1$  $\|x_{\star} - x_{k}\|_{\Lambda} < \|P(A)e_{0}\|_{\Lambda}$ 





### Polynomial residual expansions

(4/5)

Recalling that

$$A^{-1}r_k = A^{-1}(b - Ax_k) = x_\star - x_k = e_k$$

we can write

$$e_k = \mathbf{x}_{\star} - \mathbf{x}_k = \mathbf{A}^{-1} \mathbf{r}_k$$

$$= \mathbf{A}^{-1} P_k(\mathbf{A}) \mathbf{r}_0$$

$$= P_k(\mathbf{A}) \mathbf{A}^{-1} \mathbf{r}_0$$

$$= P_k(\mathbf{A}) (\mathbf{x}_{\star} - \mathbf{x}_0)$$

$$= P_k(\mathbf{A}) e_0.$$

due to the optimality of the conjugate gradient we have:



Using the results of slide 59 and 60 we can write

$$\boldsymbol{e}_k = P_k(\boldsymbol{A})\boldsymbol{e}_0,$$

$$\|e_k\|_{A} = \|P_k(A)e_0\|_{A} \le \|P(A)e_0\|_{A} \qquad \forall P \in \mathbb{P}^k, P(0) = 1$$

and from this equation we have the estimate

$$\|e_k\|_{A} \le \inf_{P \in \mathbb{P}^k, P(0)=1} \|P(A)e_0\|_{A}$$

So an estimate of the form

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \left\| P(\boldsymbol{A}) \boldsymbol{e}_0 \right\|_{\boldsymbol{A}} \le C_k \left\| \boldsymbol{e}_0 \right\|_{\boldsymbol{A}}$$

can be used to proof a convergence rate theorem, as for the steepest descent algorithm.



### Convergence rate calculation

#### Lemma

Let  $m{A} \in \mathbb{R}^{n \times n}$  an SPD matrix, and  $p \in \mathbb{P}^k$  a polynomial, then

$$\|p(\boldsymbol{A})\boldsymbol{x}\|_{\boldsymbol{A}} \leq \|p(\boldsymbol{A})\|_2 \|\boldsymbol{x}\|_{\boldsymbol{A}}$$

Proof. (1/2).

The matrix A is SPD so that we can write

$$\mathbf{A} = \mathbf{U}^T \mathbf{\Lambda} \mathbf{U}, \qquad \mathbf{\Lambda} = \text{DIAG}\{\lambda_1, \lambda_2, \dots, \lambda_n\}$$

where  $m{U}$  is an orthogonal matrix (i.e.  $m{U}^Tm{U}=m{I}$ ) and  $m{\Lambda} \geq m{0}$  is diagonal. We can define the SPD matrix  $m{A}^{1/2}$  as follows

$$A^{1/2} = U^T \Lambda^{1/2} U, \qquad \Lambda^{1/2} = \text{DIAG}\{\lambda_1^{1/2}, \lambda_2^{1/2}, \dots, \lambda_n^{1/2}\}$$

and obviously  $oldsymbol{A}^{1/2}oldsymbol{A}^{1/2}=oldsymbol{A}.$ 



#### Proof.

(2/2).

Notice that

$$\|m{x}\|_{m{A}}^2 = m{x}^Tm{A}m{x} = m{x}^Tm{A}^{1/2}m{A}^{1/2}m{x} = \left\|m{A}^{1/2}m{x}
ight\|_2^2$$

so that

$$\begin{aligned} \|p(\mathbf{A})\mathbf{x}\|_{\mathbf{A}} &= \left\|\mathbf{A}^{1/2}p(\mathbf{A})\mathbf{x}\right\|_{2} \\ &= \left\|p(\mathbf{A})\mathbf{A}^{1/2}\mathbf{x}\right\|_{2} \\ &\leq \|p(\mathbf{A})\|_{2} \left\|\mathbf{A}^{1/2}\mathbf{x}\right\|_{2} \\ &= \|p(\mathbf{A})\|_{2} \|\mathbf{x}\|_{\mathbf{A}} \end{aligned}$$



#### Lemma

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix, and  $p \in \mathbb{P}^k$  a polynomial, then

$$\|p(\boldsymbol{A})\|_2 = \max_{\lambda \in \sigma(\boldsymbol{A})} |p(\lambda)|$$

#### Proof.

The matrix  $p(\boldsymbol{A})$  is symmetric, and for a generic symmetric matrix  $\boldsymbol{B}$  we have

$$\|\boldsymbol{B}\|_2 = \max_{\lambda \in \sigma(\boldsymbol{B})} |\lambda|$$

observing that if  $\lambda$  is an eigenvalue of  $\boldsymbol{A}$  then  $p(\lambda)$  is an eigenvalue of  $p(\boldsymbol{A})$  the thesis easily follows.



Starting the error estimate

$$\|\boldsymbol{e}_k\|_{\boldsymbol{A}} \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \|P(\boldsymbol{A})\boldsymbol{e}_0\|_{\boldsymbol{A}}$$

Combining the last two lemma we easily obtain the estimate

$$\|\boldsymbol{e}_{k}\|_{\boldsymbol{A}} \leq \inf_{P \in \mathbb{P}^{k}, P(0)=1} \left[ \max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \right] \|\boldsymbol{e}_{0}\|_{\boldsymbol{A}}$$

• The convergence rate is estimated by bounding the constant

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \left[ \max_{\lambda \in \sigma(\mathbf{A})} |P(\lambda)| \right]$$





### Finite termination of Conjugate Gradient

### Theorem (Finite termination of Conjugate Gradient)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix, the the Conjugate Gradient applied to the linear system Ax = b terminate finding the exact solution in at most n-step.

#### Proof.

From the estimate

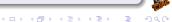
$$\|e_k\|_{oldsymbol{A}} \leq \inf_{P \in \mathbb{P}^k, P(0)=1} \left[ \max_{\lambda \in \sigma(oldsymbol{A})} |P(\lambda)| \right] \|e_0\|_{oldsymbol{A}}$$

choosing

$$P(x) = \prod_{\lambda \in \sigma(\mathbf{A})} (x - \lambda) / \prod_{\lambda \in \sigma(\mathbf{A})} (0 - \lambda)$$

we have  $\max_{\lambda \in \sigma(A)} |P(\lambda)| = 0$  and  $||e_n||_{\mathbf{A}} = 0$ .





## Convergence rate of Conjugate Gradient

The constant

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \left[ \max_{\lambda \in \sigma(\mathbf{A})} |P(\lambda)| \right]$$

is not easy to evaluate,

2 The following bound, is useful

$$\max_{\lambda \in \sigma(\boldsymbol{A})} |P(\lambda)| \le \max_{\lambda \in [\lambda_1, \lambda_n]} |P(\lambda)|$$

in particular the final estimate will be obtained by

$$\inf_{P \in \mathbb{P}^k, P(0)=1} \left[ \max_{\lambda \in \sigma(\mathbf{A})} |P(\lambda)| \right] \le \max_{\lambda \in [\lambda_1, \lambda_n]} |\bar{P}_k(\lambda)|$$

where  $\bar{P}_k(x)$  is an opportune k-degree polynomial for which  $\bar{P}_k(0)=1$  and it is easy to evaluate  $\max_{\lambda\in[\lambda_1,\lambda_n]}\left|\bar{P}_k(\lambda)\right|$ .





**1** The Chebyshev Polynomials of the First Kind are the right polynomial for this estimate. This polynomial have the following definition in the interval [-1,1]:

$$T_k(x) = \cos(k \arccos(x))$$

2 Another equivalent definition valid in the interval  $(-\infty,\infty)$  is the following

$$T_k(x) = \frac{1}{2} \left[ \left( x + \sqrt{x^2 - 1} \right)^k + \left( x - \sqrt{x^2 - 1} \right)^k \right]$$

lacksquare In spite of these definition,  $T_k(x)$  is effectively a polynomial.

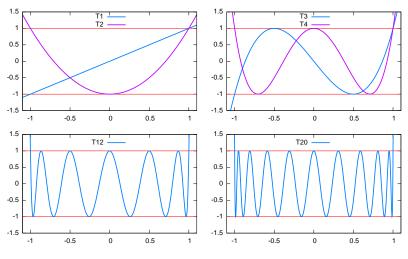




## Chebyshev Polynomials

(2/4)

Some example of Chebyshev Polynomials.





## Chebyshev Polynomials

lacksquare It is easy to show that  $T_k(x)$  is a polynomial by the use of

$$\cos(\alpha + \beta) = \cos \alpha \cos \beta - \sin \alpha \sin \beta$$
$$\cos(\alpha + \beta) + \cos(\alpha - \beta) = 2\cos \alpha \cos \beta$$

let  $\theta = \arccos(x)$ :

- **1**  $T_0(x) = \cos(0\,\theta) = 1$ ;
- **2**  $T_1(x) = \cos(1 \theta) = x$ ;
- **3**  $T_2(x) = \cos(2\theta) = \cos(\theta)^2 \sin(\theta)^2 = 2\cos(\theta)^2 1 = 2x^2 1;$
- $T_{k+1}(x) + T_{k-1}(x) = \cos((k+1)\theta) + \cos((k-1)\theta)$   $= 2\cos(k\theta)\cos(\theta) = 2xT_k(x)$
- 2 In general we have the following recurrence:
  - **1**  $T_0(x) = 1$ ;
  - **2**  $T_1(x) = x$ ;
  - $T_{k+1}(x) = 2 x T_k(x) T_{k-1}(x).$





(4/4)

- Solving the recurrence:
  - **1**  $T_0(x) = 1$ ;
  - **2**  $T_1(x) = x$ ;
  - $T_{k+1}(x) = 2xT_k(x) T_{k-1}(x).$
- We obtain the explicit form of the Chebyshev Polynomials

$$T_k(x) = \frac{1}{2} \left[ \left( x + \sqrt{x^2 - 1} \right)^k + \left( x - \sqrt{x^2 - 1} \right)^k \right]$$

 The translated and scaled polynomial is useful in the study of the conjugate gradient method:

$$T_k(x; a, b) = T_k\left(\frac{a+b-2x}{b-a}\right)$$

where we have  $|T_k(x; a, b)| \le 1$  for all  $x \in [a, b]$ .





## Convergence rate of Conjugate Gradient method

### Theorem (Convergence rate of Conjugate Gradient method)

Let  $A \in \mathbb{R}^{n \times n}$  an SPD matrix then the Conjugate Gradient method converge to the solution  $x_\star = A^{-1}b$  with at least linear r-rate in the norm  $\|\cdot\|_A$ . Moreover we have the error estimate

$$\|\boldsymbol{e}_k\|_{\boldsymbol{A}} \lesssim 2\left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k \|\boldsymbol{e}_0\|_{\boldsymbol{A}}$$

 $\kappa = M/m$  is the condition number where  $m = \lambda_1$  is the smallest eigenvalue of A and  $M = \lambda_n$  is the biggest eigenvalue of A.

The expression  $a_k \lesssim b_k$  means that for all  $\epsilon > 0$  there exists  $k_0 > 0$  such that:

$$a_k \le (1 - \epsilon)b_k, \quad \forall k > k_0$$





### Proof.

From the estimate

$$\|\boldsymbol{e}_{k}\|_{\boldsymbol{A}} \leq \max_{\lambda \in [m,M]} |P(\lambda)| \|\boldsymbol{e}_{0}\|_{\boldsymbol{A}}, \qquad P \in \mathbb{P}^{k}, P(0) = 1$$

choosing  $P(x)=T_k(x;m,M)/T_k(0;m,M)$  from the fact that  $|T_k(x;m,M)|\leq 1$  for  $x\in [m,M]$  we have

$$\|e_k\|_{\mathbf{A}} \le T_k(0; m, M)^{-1} \|e_0\|_{\mathbf{A}} = T_k \left(\frac{M+m}{M-m}\right)^{-1} \|e_0\|_{\mathbf{A}}$$

observe that  $\frac{M+m}{M-m} = \frac{\kappa+1}{\kappa-1}$  and

$$T_k \left(\frac{\kappa+1}{\kappa-1}\right)^{-1} = 2\left[\left(\frac{\sqrt{\kappa}+1}{\sqrt{\kappa}-1}\right)^k + \left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k\right]^{-1}$$

finally notice that  $\left(\frac{\sqrt{\kappa}-1}{\sqrt{\kappa}+1}\right)^k \to 0$  as  $k \to \infty$ .



## Outline

- Introduction
- 2 Convergence rate of Steepest Descent iterative scheme
- Conjugate direction method
- 4 Conjugate Gradient method
- 5 Conjugate Gradient convergence rate
- 6 Preconditioning the Conjugate Gradient method
- Nonlinear Conjugate Gradient extension



# Preconditioning

## Problem (Preconditioned linear system)

Given  $A, P \in \mathbb{R}^{n \times n}$ , with A an SPD matrix and P non singular matrix and  $b \in \mathbb{R}^n$ .

Find 
$$x_{\star} \in \mathbb{R}^n$$
 such that:  $oldsymbol{P}^{-T} oldsymbol{A} x_{\star} = oldsymbol{P}^{-T} oldsymbol{b}$ .

A good choice for P should be such that  $M = P^T P \approx A$ , where  $\approx$  denotes that M is an approximation of A in some sense to precise later.

Notice that:

• P non singular imply:

$$P^{-T}(b-Ax)=0 \iff b-Ax=0;$$

 $oldsymbol{A}$  SPD imply  $\widetilde{oldsymbol{A}} = oldsymbol{P}^{-T} oldsymbol{A} oldsymbol{P}^{-1}$  is also SPD (obvious proof).



Now we reformulate the preconditioned system:

## Problem (Preconditioned linear system)

Given  $A, P \in \mathbb{R}^{n \times n}$ , with A an SPD matrix and P non singular matrix and  $b \in \mathbb{R}^n$  the preconditioned problem is the following:

Find 
$$\widetilde{m{x}_{\star}} \in \mathbb{R}^n$$
 such that:  $\widetilde{m{A}}\widetilde{m{x}_{\star}} = \widetilde{m{b}}$ 

where

$$\widetilde{A} = P^{-T}AP^{-1}$$
  $\widetilde{b} = P^{-T}b$ 

notice that if  $x_\star$  is the solution of the linear system Ax = b then  $\widetilde{x_\star} = Px_\star$  is the solution of the linear system  $\widetilde{A}x = \widetilde{b}$ .





# PCG: preliminary version

### initial step:

$$\begin{array}{l} k \leftarrow 0; \ \boldsymbol{x}_0 \ \text{assigned}; \\ \widetilde{\boldsymbol{x}}_0 \leftarrow \boldsymbol{P} \boldsymbol{x}_0; \ \widetilde{\boldsymbol{r}}_0 \leftarrow \widetilde{\boldsymbol{b}} - \widetilde{\boldsymbol{A}} \widetilde{\boldsymbol{x}}_0; \ \widetilde{\boldsymbol{p}}_1 \leftarrow \widetilde{\boldsymbol{r}}_0; \\ \text{while } \|\widetilde{\boldsymbol{r}}_k\| > \epsilon \ \text{do} \\ k \leftarrow k + 1: \end{array}$$

## Conjugate direction method

$$\widetilde{\alpha}_{k} \leftarrow \frac{\widetilde{r}_{k-1}^{T} \widetilde{r}_{k-1}}{\widetilde{p}_{k}^{T} \widetilde{A} \widetilde{p}_{k}};$$

$$\widetilde{x}_{k} \leftarrow \widetilde{x}_{k-1} + \widetilde{\alpha}_{k} \widetilde{p}_{k};$$

$$\widetilde{r}_{k} \leftarrow \widetilde{r}_{k-1} - \widetilde{\alpha}_{k} \widetilde{A} \widetilde{p}_{k};$$

## Residual orthogonalization

$$\widetilde{\beta}_{k} \leftarrow \frac{\widetilde{r}_{k}^{T} \widetilde{r}_{k}}{\widetilde{r}_{k-1}^{T} \widetilde{r}_{k-1}};$$

$$\widetilde{p}_{k+1} \leftarrow \widetilde{r}_{k} + \widetilde{\beta}_{k} \widetilde{p}_{k};$$

#### end while

## final step

$$oldsymbol{P}^{-1}\widetilde{oldsymbol{x}}_{k}$$
;



Conjugate gradient algorithm applied to  $\widetilde{A}\widetilde{x}=\widetilde{b}$  require the evaluation of thing like:

$$\widetilde{A}\widetilde{p}_k = P^{-T}AP^{-1}\widetilde{p}_k.$$

this can be done without evaluate directly the matrix  $\widetilde{A}$ , by the following operations:

- lacksquare solve  $m{P}m{s}_k' = \widetilde{m{p}}_k$  for  $m{s}_k' = m{P}^{-1}\widetilde{m{p}}_k$ ;
- $oldsymbol{\circ}$  evaluate  $s_k'' = A s_k';$
- $\bullet$  solve  $P^Ts_k''' = s_k''$  for  $s_k''' = P^{-T}s''$ .

Step 1 and 3 require the solution of two auxiliary linear system. This is not a big problem if P and  $P^T$  are triangular matrices (see e.g. incomplete Cholesky).



However... we can reformulate the algorithm using only the matrices A and P!

#### Definition

For all  $k \geq 1$ , we introduce the vector  $q_k = P^{-1}\widetilde{p}$ .

#### Observation

If the vectors  $\widetilde{p}_1$ ,  $\widetilde{p}_2$ , ...  $\widetilde{p}_k$  for all  $1 \leq k \leq n$  are  $\widetilde{A}$ -conjugate, then the corresponding vectors  $q_1$ ,  $q_2$ , ...  $q_k$  are A-conjugate. In fact:

$$\boldsymbol{q}_{j}^{T}\boldsymbol{A}\boldsymbol{q}_{i} = \underbrace{\widetilde{\boldsymbol{p}}_{j}^{T}\boldsymbol{P}^{-T}}_{=\boldsymbol{q}_{i}^{T}}\boldsymbol{A}\underbrace{\boldsymbol{P}^{-1}\widetilde{\boldsymbol{p}}_{i}}_{=\boldsymbol{q}_{j}^{T}} = \widetilde{\boldsymbol{p}}_{j}^{T}\underbrace{\widetilde{\boldsymbol{A}}}_{=\boldsymbol{P}^{-T}\boldsymbol{A}\boldsymbol{P}^{-1}} = 0, \quad \text{if } i \neq j,$$

that is a consequence of  $\widetilde{A}$ -conjugation of vectors  $\widetilde{p}_i$ .



#### **Definition**

For all  $k \geq 1$ , we introduce the vectors

$$\boldsymbol{x}_k = \boldsymbol{x}_{k-1} + \widetilde{\alpha}_k \boldsymbol{q}_k.$$

### Observation

If we assume, by construction,  $\widetilde{x}_0 = Px_0$ , then we have

$$\widetilde{\boldsymbol{x}}_k = \boldsymbol{P} \boldsymbol{x}_k,$$
 for all  $k$  with  $1 \le k \le n$ .

In fact, if  $\widetilde{x}_{k-1} = Px_{k-1}$  (inductive hypothesis), then

$$egin{aligned} \widetilde{m{x}}_k &= \widetilde{m{x}}_{k-1} + \widetilde{lpha}_k \widetilde{m{p}}_k & [ extit{preconditioned CG}] \ &= m{P}m{x}_{k-1} + \widetilde{lpha}_k m{P}m{q}_k & [ extit{inductive Hyp. defs of }m{q}_k] \ &= m{P}m{(x}_{k-1} + \widetilde{lpha}_k m{q}_k) & [ extit{obvious}] \ &= m{P}m{x}_k & [ extit{defs. of }m{x}_k] \end{aligned}$$



#### Observation

Because  $\widetilde{x}_k = Px_k$  for all  $k \geq 0$ , we have the recurrence between the corresponding residue  $\widetilde{r}_k = \widetilde{b} - \widetilde{A}\widetilde{x}$  and  $r_k = b - Ax_k$ :

$$\widetilde{\boldsymbol{r}}_k = \boldsymbol{P}^{-T} \boldsymbol{r}_k.$$

In fact,

$$egin{aligned} \widetilde{m{r}}_k &= \widetilde{m{b}} - \widetilde{m{A}}\widetilde{m{x}}_k, & [ ext{defs. of } \widetilde{m{r}}_k] \ &= m{P}^{-T}m{b} - m{P}^{-T}m{A}m{P}^{-1}m{P}m{x}_k, & [ ext{defs. of } \widetilde{m{b}}, \ \widetilde{m{A}}, \ \widetilde{m{x}}_k] \ &= m{P}^{-T}\left(m{b} - m{A}m{x}_k
ight), & [ ext{obvious}] \ &= m{P}^{-T}m{r}_k. & [ ext{defs. of } m{r}_k] \end{aligned}$$





#### **Definition**

For all k, with  $1 \le k \le n$ , the vector  $\mathbf{z}_k$  is the solution of the linear system

$$Mz_k = r_k$$
.

where  $M = P^T P$ . Formally,

$$z_k = M^{-1}r_k = P^{-1}P^{-T}r_k.$$

Using the vectors  $\{z_k\}$ ,

- we can express  $\widetilde{\alpha}_k$  and  $\widetilde{\beta}_k$  in terms of A, the residual  $r_k$ , and conjugate direction  $q_k$ ;
- we can build a recurrence relation for the A-conjugate directions q<sub>k</sub>.



#### Observation

$$egin{aligned} \widetilde{lpha}_k &= rac{\widetilde{m{r}}_{k-1}^T \widetilde{m{r}}_{k-1}}{\widetilde{m{p}}_k^T \widetilde{m{A}} \widetilde{m{p}}_k} = rac{m{r}_{k-1} m{P}^{-1} m{P}^{-1} m{r}_{k-1}}{m{q}_k^T m{P}^T m{P}^{-T} m{A} m{P}^{-1} m{P} m{q}_k} = rac{m{r}_{k-1} m{M}^{-1} m{r}_{k-1}}{m{q}_k m{A} m{q}_k}, \ &= \boxed{rac{m{r}_{k-1} m{z}_{k-1}}{m{q}_k m{A} m{q}_k}.} \end{aligned}$$

### Observation

$$egin{aligned} \widetilde{eta}_k &= rac{\widetilde{m{r}}_k^T \widetilde{m{r}}_k}{\widetilde{m{r}}_{k-1}^T \widetilde{m{r}}_{k-1}} = rac{m{r}_k^T m{P}^{-1} m{P}^{-1} m{r}_k}{m{r}_{k-1}^T m{P}^{-1} m{P}^{-T} m{r}_{k-1}} = rac{m{r}_k^T m{M}^{-1} m{r}_k}{m{r}_{k-1}^T m{M}^{-1} m{r}_{k-1}}, \ &= \boxed{ rac{m{r}_k^T m{z}_k}{m{r}_{k-1}^T m{z}_{k-1}}. \end{aligned} }$$



#### Observation

Using the vector  $z_k = M^{-1}r_k$ , the following recurrence is true

$$\boldsymbol{q}_{k+1} = \boldsymbol{z}_k + \widetilde{\beta}_k \boldsymbol{q}_k$$

In fact:

$$egin{aligned} \widetilde{m{p}}_{k+1} &= \widetilde{m{r}}_k + \widetilde{eta}_k \widetilde{m{p}}_k & [ ext{preconditioned CG}] \ m{P}^{-1} \widetilde{m{p}}_{k+1} &= m{P}^{-1} \widetilde{m{r}}_k + \widetilde{eta}_k m{P}^{-1} \widetilde{m{p}}_k & [ ext{left mult } m{P}^{-1}] \ m{P}^{-1} \widetilde{m{p}}_{k+1} &= m{P}^{-1} m{P}^{-T} m{r}_k + \widetilde{m{\beta}}_k m{P}^{-1} \widetilde{m{p}}_k & [m{r}_{k+1} &= m{P}^{-T} m{r}_{k+1}] \ m{P}^{-1} \widetilde{m{p}}_{k+1} &= m{M}^{-1} m{r}_k + \widetilde{m{\beta}}_k m{P}^{-1} \widetilde{m{p}}_k & [m{M}^{-1} &= m{P}^{-1} m{P}^{-T}] \ m{q}_{k+1} &= m{z}_k + \widetilde{m{\beta}}_k m{q}_k & [m{q}_k &= m{P}^{-1} \widetilde{m{p}}_k] \end{aligned}$$





## PCG: final version

### initial step:

$$k \leftarrow 0$$
;  $x_0$  assigned;  $r_0 \leftarrow b - Ax_0$ ;  $q_1 \leftarrow r_0$ ; while  $||z_k|| > \epsilon$  do  $k \leftarrow k + 1$ :

### Conjugate direction method

$$\widetilde{lpha}_k \leftarrow rac{oldsymbol{r}_{k-1}^T oldsymbol{z}_{k-1}}{oldsymbol{q}_k^T oldsymbol{A} oldsymbol{q}_k}; \ oldsymbol{x}_k \leftarrow oldsymbol{x}_{k-1} + \widetilde{lpha}_k oldsymbol{q}_k;$$

$$x_k \leftarrow x_{k-1} + \alpha_k q_k$$
;

$$r_k \leftarrow r_{k-1} - \widetilde{\alpha}_k A q_k;$$

## Preconditioning

$$oldsymbol{z}_k = oldsymbol{M}^{-1} oldsymbol{r}_k;$$

#### Residual orthogonalization

$$\widetilde{\beta}_k \leftarrow \frac{\boldsymbol{r}_k^T \boldsymbol{z}_k}{\boldsymbol{r}_{k-1}^T \boldsymbol{z}_{k-1}};$$

$$oldsymbol{q}_{k+1} \leftarrow oldsymbol{z}_k + \widetilde{eta}_k oldsymbol{q}_k;$$





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# Nonlinear Conjugate Gradient extension

- The conjugate gradient algorithm can be extended for nonlinear minimization.
- ② Fletcher and Reeves extend CG for the minimization of a general non linear function f(x) as follows:
  - **1** Substitute the evaluation of  $\alpha_k$  by an line search
  - **2** Substitute the residual  $r_k$  with the gradient  $\nabla \mathsf{f}(x_k)$
- $oldsymbol{\circ}$  We also translate the index for the search direction  $oldsymbol{p}_k$  to be more consistent with the gradients. The resulting algorithm is in the next slide



# Fletcher and Reeves Nonlinear Conjugate Gradient

#### initial step:

$$\begin{aligned} k &\leftarrow 0; \ \boldsymbol{x}_0 \ \text{assigned}; \\ f_0 &\leftarrow \mathbf{f}(\boldsymbol{x}_0); \ \boldsymbol{g}_0 \leftarrow \nabla \mathbf{f}(\boldsymbol{x}_0)^T; \\ \boldsymbol{p}_0 &\leftarrow -\boldsymbol{g}_0; \\ \mathbf{while} \ \|\boldsymbol{g}_k\| &> \epsilon \ \mathbf{do} \\ k &\leftarrow k+1; \\ \mathbf{Conjugate \ direction \ method} \\ \mathbf{Compute} \ \alpha_k \ \text{by line-search}; \\ \boldsymbol{x}_k &\leftarrow \boldsymbol{x}_{k-1} + \alpha_k \boldsymbol{p}_{k-1}; \\ \boldsymbol{g}_k &\leftarrow \nabla \mathbf{f}(\boldsymbol{x}_k)^T; \\ \mathbf{Residual \ orthogonalization} \\ \boldsymbol{\beta}_k^{FR} &\leftarrow \frac{\boldsymbol{g}_k^T \boldsymbol{g}_k}{\boldsymbol{g}_{k-1}^T \boldsymbol{g}_{k-1}}; \\ \boldsymbol{p}_k &\leftarrow -\boldsymbol{g}_k + \boldsymbol{\beta}_k^{FR} \boldsymbol{p}_{k-1}; \\ \mathbf{end \ while} \end{aligned}$$



- lacktriangledown To ensure convergence and apply Zoutendijk global convergence theorem we need to ensure that  $m{p}_k$  is a descent direction.
- $oldsymbol{2}$   $oldsymbol{p}_0$  is a descent direction by construction, for  $oldsymbol{p}_k$  we have

$$\left\|oldsymbol{g}_k^Toldsymbol{p}_k = -\left\|oldsymbol{g}_k
ight\|^2 + eta_k^{FR}oldsymbol{g}_k^Toldsymbol{p}_{k-1}$$

if the line-search is exact than  $g_k^T p_{k-1} = 0$  because  $p_{k-1}$  is the direction of the line-search. So by induction  $p_k$  is a descent direction.

- Exact line-search is expensive, however if we use inexact line-search with strong Wolfe conditions
  - sufficient decrease:  $f(x_k + \alpha_k p_k) \le f(x_k) + c_1 \alpha_k \nabla f(x_k) p_k$ ;
  - **2** curvature condition:  $|\nabla f(x_k + \alpha_k p_k)p_k| \leq c_2 |\nabla f(x_k)p_k|$ .

with  $0 < c_1 < c_2 < 1/2$  then we can prove that  $p_k$  is a descent direction.



The previous consideration permits to say that Fletcher and Reeves nonlinear conjugate gradient method with strong Wolfe line-search is globally convergent<sup>1</sup>

To prove globally convergence we need the following lemma:

## Lemma (descent direction bound)

Suppose we apply Fletcher and Reeves nonlinear conjugate gradient method to  $f(\boldsymbol{x})$  with strong Wolfe line-search with  $0 < c_2 < 1/2$ . The the method generates descent direction  $\boldsymbol{p}_k$  that satisfy the following inequality

$$-\frac{1}{1-c_2} \le \frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{q}_k\|^2} \le -\frac{1-2c_2}{1-c_2}, \qquad k = 0, 1, 2, \dots$$



<sup>&</sup>lt;sup>1</sup>globally here means that Zoutendijk like theorem apply (3) + (3) + (3) + (3)

Proof. (1/3).

The proof is by induction. First notice that the function

$$t(\xi) = \frac{2\xi - 1}{1 - \xi}$$

is monotonically increasing on the interval [0,1/2] and that t(0)=-1 and t(1/2)=0. Hence, because of  $c_2\in(0,1/2)$  we have:

$$-1 < \frac{2c_2 - 1}{1 - c_2} < 0. \tag{*}$$

base of induction k = 0: For k = 0 we have  $p_0 = -g_0$  so that  $g_0^T p_0 / ||g_0||^2 = -1$ . From  $(\star)$  the lemma inequality is trivially satisfied.





Proof. (2/3).

Using update direction formula's of the algorithm:

$$eta_k^{FR} = rac{oldsymbol{g}_k^T oldsymbol{g}_k}{oldsymbol{g}_{k-1}^T oldsymbol{g}_{k-1}} \qquad oldsymbol{p}_k = -oldsymbol{g}_k + eta_k^{FR} oldsymbol{p}_{k-1}$$

we can write

$$\frac{{\boldsymbol{g}_k^T \boldsymbol{p}_k}}{{{{\left\| {\boldsymbol{g}_k} \right\|}^2}}} = - 1 + \beta _k^{FR} \frac{{\boldsymbol{g}_k^T \boldsymbol{p}_{k - 1}}}{{{{\left\| {\boldsymbol{g}_k} \right\|}^2}}} = - 1 + \frac{{\boldsymbol{g}_k^T \boldsymbol{p}_{k - 1}}}{{{{\left\| {\boldsymbol{g}_{k - 1}} \right\|}^2}}}$$

and by using second strong Wolfe condition:

$$-1 + c_2 \frac{\boldsymbol{g}_{k-1}^T \boldsymbol{p}_{k-1}}{\|\boldsymbol{g}_{k-1}\|^2} \le \frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{g}_k\|^2} \le -1 - c_2 \frac{\boldsymbol{g}_{k-1}^T \boldsymbol{p}_{k-1}}{\|\boldsymbol{g}_{k-1}\|^2}$$





Proof. (3/3).

by induction we have

$$\frac{1}{1 - c_2} \ge -\frac{\boldsymbol{g}_{k-1}^T \boldsymbol{p}_{k-1}}{\|\boldsymbol{g}_{k-1}\|^2} > 0$$

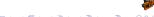
so that

$$\frac{\boldsymbol{g}_{k}^{T}\boldsymbol{p}_{k}}{\left\|\boldsymbol{g}_{k}\right\|^{2}} \leq -1 - c_{2} \frac{\boldsymbol{g}_{k-1}^{T}\boldsymbol{p}_{k-1}}{\left\|\boldsymbol{g}_{k-1}\right\|^{2}} \leq -1 + c_{2} \frac{1}{1 - c_{2}} = \frac{2c_{2} - 1}{1 - c_{2}}$$

and

$$\frac{\boldsymbol{g}_{k}^{T}\boldsymbol{p}_{k}}{\|\boldsymbol{q}_{k}\|^{2}} \ge -1 + c_{2}\frac{\boldsymbol{g}_{k-1}^{T}\boldsymbol{p}_{k-1}}{\|\boldsymbol{q}_{k-1}\|^{2}} \ge -1 - c_{2}\frac{1}{1 - c_{2}} = -\frac{1}{1 - c_{2}}$$





The inequality of the the previous lemma can be written as:

$$\frac{1}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} \ge -\frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{g}_k\| \|\boldsymbol{p}_k\|} \ge \frac{1 - 2c_2}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} > 0$$

Remembering the Zoutendijk theorem we have

$$\sum_{k=1}^{\infty} (\cos \theta_k)^2 \|\boldsymbol{g}_k\|^2 < \infty, \quad \text{where} \quad \cos \theta_k = -\frac{\boldsymbol{g}_k^T \boldsymbol{p}_k}{\|\boldsymbol{g}_k\| \|\boldsymbol{p}_k\|}$$

- **3** so that if  $\|g_k\|/\|p_k\|$  is bounded from below we have that  $\cos \theta_k \geq \delta$  for all k and then from Zoutendijk theorem the scheme converge.
- ① Unfortunately this bound cant be proved so that Zoutendijk theorem cant be applied directly. However it is possible to prove a weaker results, i.e. that  $\liminf_{k\to\infty}\|g_k\|=0!$





# Convergence of Fletcher and Reeves method

## Assumption (Regularity assumption)

We assume  $f \in C^1(\mathbb{R}^n)$  with Lipschitz continuous gradient, i.e. there exists  $\gamma > 0$  such that

$$\|\nabla f(\boldsymbol{x})^T - \nabla f(\boldsymbol{y})^T\| \le \gamma \|\boldsymbol{x} - \boldsymbol{y}\|, \quad \forall \boldsymbol{x}, \boldsymbol{y} \in \mathbb{R}^n$$



## Theorem (Convergence of Fletcher and Reeves method)

Suppose the method of Fletcher and Reeves is implemented with strong Wolfe line-search with  $0 < c_1 < c_2 < 1/2$ . If  $f(\boldsymbol{x})$  and  $\boldsymbol{x}_0$  satisfy the previous regularity assumptions, then

$$\liminf_{k \to \infty} \|\boldsymbol{g}_k\| = 0$$

Proof. (1/4).

From previous Lemma we have

$$\cos \theta_k \ge \frac{1}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} \qquad k = 1, 2, \dots$$

substituting in Zoutendijk condition we have  $\sum_{k=1}^{\infty} \frac{\|m{g}_k\|^4}{\|m{p}_k\|^2} < \infty.$ 

The proof is by contradiction. in fact if theorem is not true than the series diverge. Next we want to bound  $\|p_k\|$ .



## Proof. (bounding $\|p_k\|$ )

(2/4).

Using second Wolfe condition and previous Lemma

$$\left| \boldsymbol{g}_{k}^{T} \boldsymbol{p}_{k-1} \right| \leq -c_{2} \boldsymbol{g}_{k}^{T} \boldsymbol{p}_{k-1} \leq \frac{c_{2}}{1-c_{2}} \left\| \boldsymbol{g}_{k-1} \right\|^{2}$$

using  $oldsymbol{p}_k \leftarrow -oldsymbol{g}_k + eta_k^{FR} oldsymbol{p}_{k-1}$  we have

$$\|\boldsymbol{p}_{k}\|^{2} \leq \|\boldsymbol{g}_{k}\|^{2} + 2\beta_{k}^{FR} |\boldsymbol{g}_{k}^{T} \boldsymbol{p}_{k-1}| + (\beta_{k}^{FR})^{2} \|\boldsymbol{p}_{k-1}\|^{2}$$

$$\leq \|\boldsymbol{g}_{k}\|^{2} + \frac{2c_{2}}{1 - c_{2}} \beta_{k}^{FR} \|\boldsymbol{g}_{k-1}\|^{2} + (\beta_{k}^{FR})^{2} \|\boldsymbol{p}_{k-1}\|^{2}$$

recall that  $eta_k^{FR} \leftarrow \|oldsymbol{g}_k\|^2 \, / \, \|oldsymbol{g}_{k-1}\|^2$  then

$$\left\| \boldsymbol{p}_{k} \right\|^{2} \leq \frac{1 + c_{2}}{1 - c_{2}} \left\| \boldsymbol{g}_{k} \right\|^{2} + (\beta_{k}^{FR})^{2} \left\| \boldsymbol{p}_{k-1} \right\|^{2}$$





## Proof. (bounding $\|\boldsymbol{p}_k\|$ )

(3/4).

setting  $c_3 = \frac{1+c_2}{1-c_2}$  and using repeatedly the last inequality we obtain:

$$\begin{split} \left\| \boldsymbol{p}_{k} \right\|^{2} & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{2} + (\beta_{k}^{FR})^{2} \left( c_{3} \left\| \boldsymbol{g}_{k-1} \right\|^{2} + (\beta_{k-1}^{FR})^{2} \left\| \boldsymbol{p}_{k-2} \right\|^{2} \right) \\ & = c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \left( \left\| \boldsymbol{g}_{k} \right\|^{-2} + \left\| \boldsymbol{g}_{k-1} \right\|^{-2} \right) + \frac{\left\| \boldsymbol{g}_{k} \right\|^{4}}{\left\| \boldsymbol{g}_{k-2} \right\|^{4}} \left\| \boldsymbol{p}_{k-2} \right\|^{2} \\ & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \left( \left\| \boldsymbol{g}_{k} \right\|^{-2} + \left\| \boldsymbol{g}_{k-1} \right\|^{-2} + \left\| \boldsymbol{g}_{k-2} \right\|^{-2} \right) \\ & + \frac{\left\| \boldsymbol{g}_{k} \right\|^{4}}{\left\| \boldsymbol{g}_{k-3} \right\|^{4}} \left\| \boldsymbol{p}_{k-3} \right\|^{2} \\ & \leq c_{3} \left\| \boldsymbol{g}_{k} \right\|^{4} \sum_{i=1}^{k} \left\| \boldsymbol{g}_{i} \right\|^{-2} \end{split}$$



### Proof.

(4/4).

Suppose now by contradiction there exists  $\delta>0$  such that  $\|g_k\|\geq \delta$  a by using the regularity assumptions we have

$$\|\boldsymbol{p}_k\|^2 \le c_3 \|\boldsymbol{g}_k\|^4 \sum_{j=1}^k \|\boldsymbol{g}_j\|^{-2} \le c_3 \|\boldsymbol{g}_k\|^4 \delta^{-2} k$$

Substituting in Zoutendijk condition we have

$$\infty > \sum_{k=1}^{\infty} \frac{\|g_k\|^4}{\|p_k\|^2} \ge \frac{\delta^2}{c_4} \sum_{k=1}^{\infty} \frac{1}{k} = \infty$$

this contradict assumption.





<sup>&</sup>lt;sup>a</sup>the correct assumption is that there exists  $k_0$  such that  $\|g_k\| \ge \delta$  for  $k \ge k_0$  but this complicate a little bit the following inequality without introducing new idea.

## Weakness of Fletcher and Reeves method

- Suppose that  $p_k$  is a bad search direction, i.e.  $\cos \theta_k \approx 0$ .
- From the descent direction bound Lemma (see slide 90) we have

$$\frac{1}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} \ge \cos \theta_k \ge \frac{1 - 2c_2}{1 - c_2} \frac{\|\boldsymbol{g}_k\|}{\|\boldsymbol{p}_k\|} > 0$$

- so that to have  $\cos \theta_k \approx 0$  we needs  $\|\boldsymbol{p}_k\| \gg \|\boldsymbol{g}_k\|$ .
- since  $p_k$  is a bad direction near orthogonal to  $g_k$  it is likely that the step is small and  $x_{k+1} \approx x_k$ . If so we have also  $g_{k+1} \approx g_k$  and  $\beta_{k+1}^{FR} \approx 1$ .
- but remember that  $m{p}_{k+1} \leftarrow -m{g}_{k+1} + eta_{k+1}^{FR} m{p}_k$ , so that  $m{p}_{k+1} pprox m{p}_k$ .
- This means that a long sequence of unproductive iterates will follows.



# Polack and Ribiére Nonlinear Conjugate Gradient

- The previous problem can be elided if we restart anew when the iterate stagnate.
- 2 Restarting is obtained by simply set  $\beta_k^{FR} = 0$ .
- **3** A more elegant solution can be obtained with a new definition of  $\beta_k$  due to Polack and Ribiére is the following:

$$\beta_k^{PR} = \frac{\boldsymbol{g}_k^T (\boldsymbol{g}_k - \boldsymbol{g}_{k-1})}{\boldsymbol{g}_{k-1}^T \boldsymbol{g}_{k-1}}$$

• This definition of  $\beta_k^{PR}$  is identical of  $\beta_k^{FR}$  in the case of quadratic function because  ${\boldsymbol g}_k^T{\boldsymbol g}_{k-1}=0$ . The definition differs in non linear case and in particular when there is stagnation i.e.  ${\boldsymbol g}_k \approx {\boldsymbol g}_{k-1}$  we have  $\beta_k^{PR} \approx 0$ , i.e. we have an automatic restart.



# Polack and Ribiére Nonlinear Conjugate Gradient

### initial step:

$$\begin{aligned} k &\leftarrow 0; \ \boldsymbol{x}_0 \ \text{assigned}; \\ f_0 &\leftarrow \mathsf{f}(\boldsymbol{x}_0); \ \boldsymbol{g}_0 \leftarrow \nabla \mathsf{f}(\boldsymbol{x}_0)^T; \\ \boldsymbol{p}_0 &\leftarrow -\boldsymbol{g}_0; \\ \textbf{while} \ \|\boldsymbol{g}_k\| > \epsilon \ \textbf{do} \\ k &\leftarrow k+1; \\ \textbf{Conjugate direction method} \\ \textbf{Compute} \ \alpha_k \ \text{by line-search}; \\ \boldsymbol{x}_k &\leftarrow \boldsymbol{x}_{k-1} + \alpha_k \boldsymbol{p}_{k-1}; \\ \boldsymbol{g}_k &\leftarrow \nabla \mathsf{f}(\boldsymbol{x}_k)^T; \\ \textbf{Residual orthogonalization} \\ \beta_k^{PR} &\leftarrow \frac{\boldsymbol{g}_k^T(\boldsymbol{g}_k - \boldsymbol{g}_{k-1})}{\boldsymbol{g}_{k-1}^T \boldsymbol{g}_{k-1}}; \\ \boldsymbol{p}_k &\leftarrow -\boldsymbol{g}_k + \beta_k^{PR} \boldsymbol{p}_{k-1}; \\ \textbf{end while} \end{aligned}$$



(1/2)

- Although the modification is minimal, for the Polack and Ribiére method with strong Wolfe line-search it can happen that  $p_k$  is not a descent direction.
- If  $p_k$  is not a descent direction we can restart i.e. set  $\beta_k^{PR}=0$  or modify  $\beta_k^{PR}$  as follows

$$\beta_k^{PR+} = \max\{\beta_k^{PR}, 0\}$$

this new coefficient with a modified Wolfe line-search ensure that  $p_k$  is a descent direction.





- Polack and Ribiére choice on the average perform better than Fletcher and Reeves but there is not convergence results!
- Although there is not convergence results there is a negative results due to Powell:

#### Theorem

Consider the Polack and Ribiére method with exact line-search. There exists a twice continuously differentiable function  $f: \mathbb{R}^3 \mapsto \mathbb{R}$  and a starting point  $x_0$  such that the sequence of gradients  $\{ \|g_k\| \}$  is bounded away from zero.

 However is spite of this results Polack and Ribiére is the first choice among conjugate direction methods.



## Other choices

• There are many other modification of the coefficient  $\beta_k$  that collapse to the same coefficient in the case o quadratic function. One important choice is the Hestenes and Stiefel choice

$$\beta_k^{HS} = \frac{\boldsymbol{g}_k^T (\boldsymbol{g}_k - \boldsymbol{g}_{k-1})}{(\boldsymbol{g}_k^T - \boldsymbol{g}_{k-1}^T) \boldsymbol{p}_{k-1}}$$

 For this choice there is similar convergence results of Fletcher and Reeves and similar performance.



## References



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